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dsPIC30F2010
Data Sheet

High-Performance,
16-bit Digital Signal Controllers

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dsPIC30F2010

High-Performance, 16-bit Digital Signal Controller

Note:	This data sheet summarizes features of this group of dsPIC30F devices and is not intended to be a complete reference source. For more information on the CPU, peripherals, register descriptions and general device functionality, refer to the "*dsPIC30F Family Reference Manual*" (DS70046). For more information on the device instruction set and programming, refer to the "*16-bit MCU and DSC Programmer's Reference Manual*" (DS70157).

High-Performance Modified RISC CPU:

- Modified Harvard architecture
- C compiler optimized instruction set architecture
- 83 base instructions with flexible addressing modes
- 24-bit wide instructions, 16-bit wide data path
- 12 Kbytes on-chip Flash program space
- 512 bytes on-chip data RAM
- 1 Kbyte nonvolatile data EEPROM
- 16 x 16-bit working register array
- Up to 30 MIPs operation:
  - DC to 40 MHz external clock input
  - 4 MHz-10 MHz oscillator input with PLL active (4x, 8x, 16x)
- 27 interrupt sources
- Three external interrupt sources
- Eight user-selectable priority levels for each interrupt
- Four processor exceptions and software traps

DSP Engine Features:

- Modulo and Bit-Reversed modes
- Two 40-bit wide accumulators with optional saturation logic
- 17-bit x 17-bit single-cycle hardware fractional/ integer multiplier
- Single-cycle Multiply-Accumulate (MAC) operation
- 40-stage Barrel Shifter
- Dual data fetch

Peripheral Features:

- High current sink/source I/O pins: 25 mA/25 mA
- Three 16-bit timers/counters; optionally pair up 16-bit timers into 32-bit timer modules
- Four 16-bit capture input functions
- Two 16-bit compare/PWM output functions
  - Dual Compare mode available
- 3-wire SPI modules (supports 4 Frame modes)
- I^{2}C™ module supports Multi-Master/Slave mode and 7-bit/10-bit addressing
- Addressable UART modules with FIFO buffers

Motor Control PWM Module Features:

- Six PWM output channels
  - Complementary or Independent Output modes
  - Edge and Center-Aligned modes
- Four duty cycle generators
- Dedicated time base with four modes
- Programmable output polarity
- Dead-time control for Complementary mode
- Manual output control
- Trigger for synchronized A/D conversions

Quadrature Encoder Interface Module Features:

- Phase A, Phase B and Index Pulse input
- 16-bit up/down position counter
- Count direction status
- Position Measurement (x2 and x4) mode
- Programmable digital noise filters on inputs
- Alternate 16-bit Timer/Counter mode
- Interrupt on position counter rollover/underflow

Analog Features:

- 10-bit Analog-to-Digital Converter (ADC) with:
  - 1 Msps (for 10-bit A/D) conversion rate
  - Six input channels
  - Conversion available during Sleep and Idle
- Programmable Brown-out Reset

Special Digital Signal Controller Features:

- Enhanced Flash program memory:
  - 10,000 erase/write cycle (min.) for industrial temperature range, 100K (typical)
- Data EEPROM memory:
  - 100,000 erase/write cycle (min.) for industrial temperature range, 1M (typical)
- Self-reprogrammable under software control
- Power-on Reset (POR), Power-up Timer (PWRT) and Oscillator Start-up Timer (OST)
- Flexible Watchdog Timer (WDT) with on-chip low-power RC oscillator for reliable operation
- Fail-Safe Clock Monitor (FSCM) operation
- Detects clock failure and switches to on-chip Low-Power RC (LPRC) oscillator
- Programmable code protection
- In-Circuit Serial Programming™ (ICSP™) programming capability
- Selectable Power Management modes
  - Sleep, Idle and Alternate Clock modes

CMOS Technology:

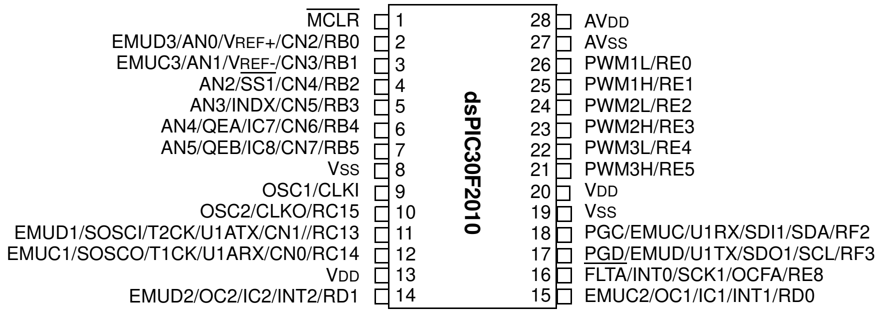
- Low-power, high-speed Flash technology
- Wide operating voltage range (2.5V to 5.5V)
- Industrial and Extended temperature ranges
- Low power consumption

dsPIC30F Motor Control and Power Conversion Family

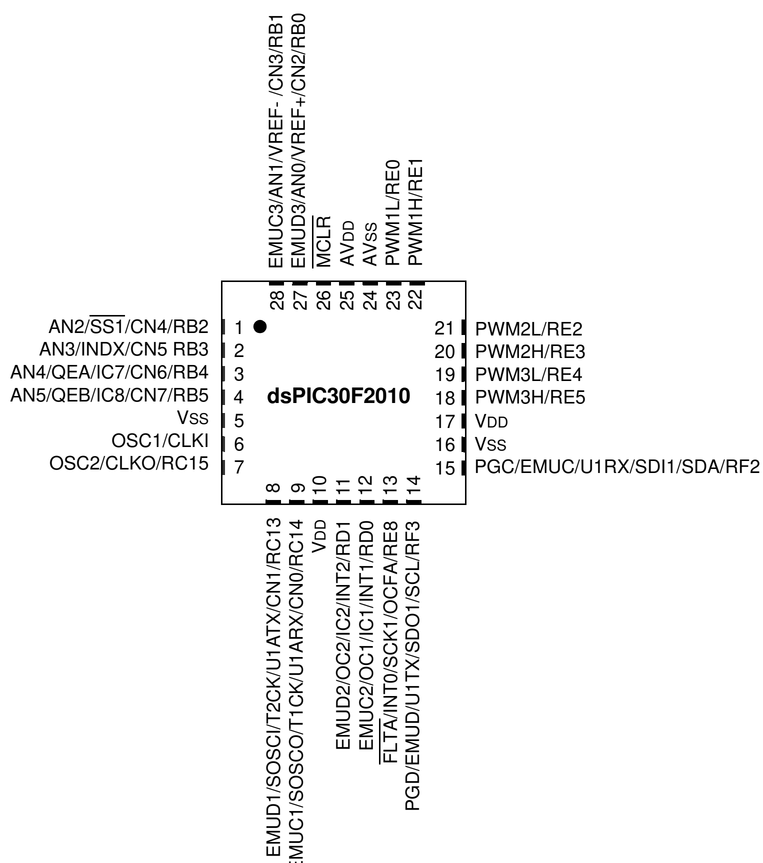
Device	Pins	Program Mem. Bytes/ Instructions	SRAM Bytes	EEPROM Bytes	Timer 16-bit	Input Cap	Output Comp/Std PWM	Motor Control PWM	A/D 10-bit 1 Msps	QEI	UART	SPI	I²C™
dsPIC30F2010	28	12K/4K	512	1024	3	4	2	6 ch	6 ch	Yes	1	1	1

Pin Diagrams

28-Pin SDIP and SOIC



28-Pin QFN-S(1)



Note 1: The metal plane at the bottom of the device is not connected to any pins and is recommended to be connected to VSS externally.

Table of Contents

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1.0 DEVICE OVERVIEW

Note:	This data sheet summarizes features of this group of dsPIC30F devices and is not intended to be a complete reference source. For more information on the CPU, peripherals, register descriptions and general device functionality, refer to the *"dsPIC30F Family Reference Manual"* (DS70046). For more information on the device instruction set and programming, refer to the *"16-bit MCU and DSC Programmer's Reference Manual"* (DS70157).

This document contains device specific information for the dsPIC30F2010 device. The dsPIC30F devices contain extensive Digital Signal Processor (DSP) functionality within a high-performance 16-bit microcontroller (MCU) architecture. Figure 1-1 shows a device block diagram for the dsPIC30F2010 device.

FIGURE 1-1: dsPIC30F2010 BLOCK DIAGRAM

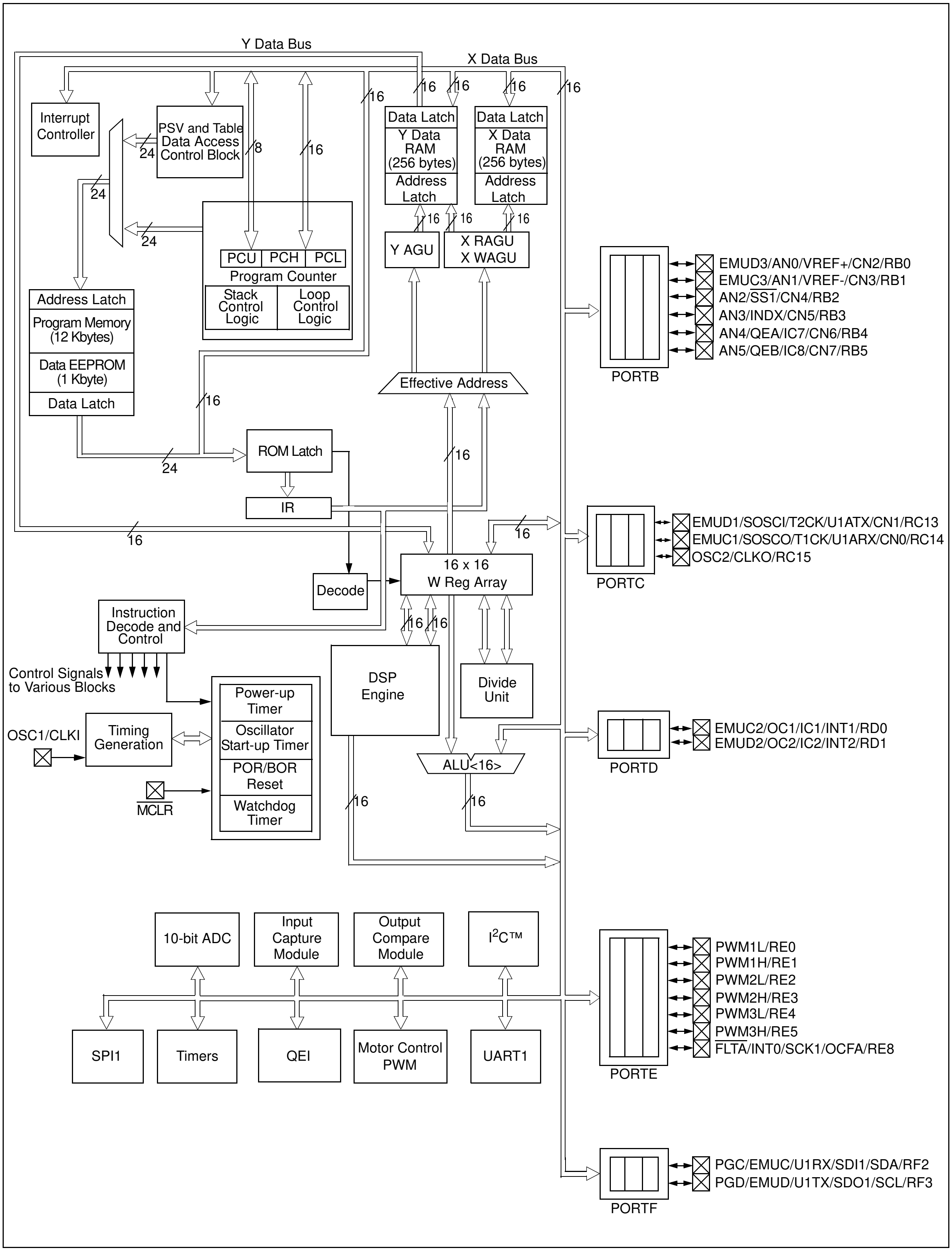


Table 1-1 provides a brief description of device I/O pinouts and the functions that may be multiplexed to a port pin. Multiple functions may exist on one port pin. When multiplexing occurs, the peripheral module's functional requirements may force an override of the data direction of the port pin.

TABLE 1-1: PINOUT I/O DESCRIPTIONS

Pin Name	Pin Type	Buffer Type	Description
AN0-AN5	I	Analog	Analog input channels.
AVDD	P	P	Positive supply for analog module. This pin must be connected at all times.
AVSS	P	P	Ground reference for analog module. This pin must be connected at all times.
CLKI CLKO	I O	ST/CMOS —	External clock source input. Always associated with OSC1 pin function. Oscillator crystal output. Connects to crystal or resonator in Crystal Oscillator mode. Optionally functions as CLKO in RC and EC modes. Always associated with OSC2 pin function.
CN0-CN7	I	ST	Input change notification inputs. Can be software programmed for internal weak pull-ups on all inputs.
EMUD EMUC EMUD1 EMUC1 EMUD2 EMUC2 EMUD3 EMUC3	I/O I/O I/O I/O I/O I/O I/O I/O	ST ST ST ST ST ST ST ST	ICD Primary Communication Channel data input/output pin. ICD Primary Communication Channel clock input/output pin. ICD Secondary Communication Channel data input/output pin. ICD Secondary Communication Channel clock input/output pin. ICD Tertiary Communication Channel data input/output pin. ICD Tertiary Communication Channel clock input/output pin. ICD Quaternary Communication Channel data input/output pin. ICD Quaternary Communication Channel clock input/output pin.
IC1, IC2, IC7, IC8	I	ST	Capture inputs. The dsPIC30F2010 has four capture inputs. The inputs are numbered for consistency with the inputs on larger device variants.
INDX QEA QEB	I I I	ST ST ST	Quadrature Encoder Index Pulse input. Quadrature Encoder Phase A input in QEI mode. Auxiliary Timer External Clock/Gate input in Timer mode. Quadrature Encoder Phase B input in QEI mode. Auxiliary Timer External Clock/Gate input in Timer mode.
INT0 INT1 INT2	I I I	ST ST ST	External interrupt 0 External interrupt 1 External interrupt 2
$\overline{FLTA}$ PWM1L PWM1H PWM2L PWM2H PWM3L PWM3H	I O O O O O O	ST — — — — — —	PWM Fault A input PWM 1 Low output PWM 1 High output PWM 2 Low output PWM 2 High output PWM 3 Low output PWM 3 High output
$\overline{MCLR}$	I/P	ST	Master Clear (Reset) input or programming voltage input. This pin is an active-low Reset to the device.
OCFA OC1-OC2	I O	ST —	Compare Fault A input (for Compare channels 1, 2, 3 and 4). Compare outputs.
OSC1 OSC2	I I/O	ST/CMOS —	Oscillator crystal input. ST buffer when configured in RC mode; CMOS otherwise. Oscillator crystal output. Connects to crystal or resonator in Crystal Oscillator mode. Optionally functions as CLKO in RC and EC modes.

Legend: CMOS = CMOS compatible input or output; Analog = Analog input
ST = Schmitt Trigger input with CMOS levels; O = Output
I = Input; P = Power

TABLE 1-1: PINOUT I/O DESCRIPTIONS (CONTINUED)

Pin Name	Pin Type	Buffer Type	Description
PGD PGC	I/O I	ST ST	In-Circuit Serial Programming™ (ICSP™) data input/output pin. In-Circuit Serial Programming clock input pin.
RB0-RB5	I/O	ST	PORTB is a bidirectional I/O port.
RC13-RC14	I/O	ST	PORTC is a bidirectional I/O port.
RD0-RD1	I/O	ST	PORTD is a bidirectional I/O port.
RE0-RE5, RE8	I/O	ST	PORTE is a bidirectional I/O port.
RF2, RF3	I/O	ST	PORTF is a bidirectional I/O port.
SCK1 SDI1 SDO1 $\overline{SS1}$	I/O I O I	ST ST — ST	Synchronous serial clock input/output for SPI1. SPI1 Data In. SPI1 Data Out. SPI1 Slave Synchronization.
SCL SDA	I/O I/O	ST ST	Synchronous serial clock input/output for I^{2}C™. Synchronous serial data input/output for I^{2}C.
SOSCO SOSCI	O I	— ST/CMOS	32 kHz low-power oscillator crystal output. 32 kHz low-power oscillator crystal input. ST buffer when configured in RC mode; CMOS otherwise.
T1CK T2CK	I I	ST ST	Timer1 external clock input. Timer2 external clock input.
U1RX U1TX U1ARX U1ATX	I O I O	ST — ST —	UART1 Receive. UART1 Transmit. UART1 Alternate Receive. UART1 Alternate Transmit.
VDD	P	—	Positive supply for logic and I/O pins.
VSS	P	—	Ground reference for logic and I/O pins.
VREF+	I	Analog	Analog Voltage Reference (High) input.
VREF-	I	Analog	Analog Voltage Reference (Low) input.

Legend: CMOS = CMOS compatible input or output; Analog = Analog input
ST = Schmitt Trigger input with CMOS levels; O = Output
I = Input; P = Power

2.0 CPU ARCHITECTURE OVERVIEW

Note:	This data sheet summarizes features of this group of dsPIC30F devices and is not intended to be a complete reference source. For more information on the CPU, peripherals, register descriptions and general device functionality, refer to the "*dsPIC30F Family Reference Manual*" (DS70046). For more information on the device instruction set and programming, refer to the "*16-bit MCU and DSC Programmer's Reference Manual*" (DS70157).

2.1 Core Overview

The core has a 24-bit instruction word. The Program Counter (PC) is 23 bits wide with the Least Significant bit (LSb) always clear (see **Section 3.1 "Program Address Space"**), and the Most Significant bit (MSb) is ignored during normal program execution, except for certain specialized instructions. Thus, the PC can address up to 4M instruction words of user program space. An instruction prefetch mechanism is used to help maintain throughput. Program loop constructs, free from loop count management overhead, are supported using the `DO` and `REPEAT` instructions, both of which are interruptible at any point.

The working register array consists of 16x16-bit registers, each of which can act as data, address or offset registers. One working register (W15) operates as a software Stack Pointer for interrupts and calls.

The data space is 64 Kbytes (32K words) and is split into two blocks, referred to as X and Y data memory. Each block has its own independent Address Generation Unit (AGU). Most instructions operate solely through the X memory AGU, which provides the appearance of a single unified data space. The Multiply-Accumulate (`MAC`) class of dual source DSP instructions operate through both the X and Y AGUs, splitting the data address space into two parts (see **Section 3.2 "Data Address Space"**). The X and Y data space boundary is device specific and cannot be altered by the user. Each data word consists of 2 bytes, and most instructions can address data either as words or bytes.

There are two methods of accessing data stored in program memory:

- The upper 32 Kbytes of data space memory can be mapped into the lower half (user space) of program space at any 16K program word boundary, defined by the 8-bit Program Space Visibility Page (PSVPAG) register. This lets any instruction access program space as if it were data space, with a limitation that the access requires an additional cycle. Moreover, only the lower 16 bits of each instruction word can be accessed using this method.
- Linear indirect access of 32K word pages within program space is also possible using any working register, via table read and write instructions. Table read and write instructions can be used to access all 24 bits of an instruction word.

Overhead-free circular buffers (Modulo Addressing) are supported in both X and Y address spaces. This is primarily intended to remove the loop overhead for DSP algorithms.

The X AGU also supports Bit-Reversed Addressing on destination effective addresses, to greatly simplify input or output data reordering for radix-2 FFT algorithms. Refer to **Section 4.0 "Address Generator Units"** for details on Modulo and Bit-Reversed Addressing.

The core supports Inherent (no operand), Relative, Literal, Memory Direct, Register Direct, Register Indirect, Register Offset and Literal Offset Addressing modes. Instructions are associated with predefined Addressing modes, depending upon their functional requirements.

For most instructions, the core is capable of executing a data (or program data) memory read, a working register (data) read, a data memory write and a program (instruction) memory read per instruction cycle. As a result, 3-operand instructions are supported, allowing C = A + B operations to be executed in a single cycle.

A DSP engine has been included to significantly enhance the core arithmetic capability and throughput. It features a high-speed 17-bit by 17-bit multiplier, a 40-bit ALU, two 40-bit saturating accumulators and a 40-bit bidirectional barrel shifter. Data in the accumulator or any working register can be shifted up to 15 bits right or 16 bits left in a single cycle. The DSP instructions operate seamlessly with all other instructions and have been designed for optimal real-time performance. The `MAC` class of instructions can concurrently fetch two data operands from memory, while multiplying two W registers. To enable this concurrent fetching of data operands, the data space has been split for these instructions and linear for all others. This has been achieved in a transparent and flexible manner, by dedicating certain working registers to each address space for the `MAC` class of instructions.

The core does not support a multi-stage instruction pipeline. However, a single stage instruction prefetch mechanism is used, which accesses and partially decodes instructions a cycle ahead of execution, in order to maximize available execution time. Most instructions execute in a single cycle, with certain exceptions.

The core features a vectored exception processing structure for traps and interrupts, with 62 independent vectors. The exceptions consist of up to 8 traps (of which 4 are reserved) and 54 interrupts. Each interrupt is prioritized based on a user-assigned priority between 1 and 7 (1 being the lowest priority and 7 being the highest) in conjunction with a predetermined 'natural order'. Traps have fixed priorities, ranging from 8 to 15.

2.2 Programmer's Model

The programmer's model is shown in Figure 2-1 and consists of 16 x 16-bit working registers (W0 through W15), 2 x 40-bit accumulators (ACCA and ACCB), STATUS Register (SR), Data Table Page register (TBLPAG), Program Space Visibility Page register (PSVPAG), DO and REPEAT registers (DOSTART, DOEND, DCOUNT and RCOUNT) and Program Counter (PC). The working registers can act as data, address or offset registers. All registers are memory mapped. W0 acts as the W register for file register addressing.

Some of these registers have a shadow register associated with each of them, as shown in Figure 2-1. The shadow register is used as a temporary holding register and can transfer its contents to or from its host register upon the occurrence of an event. None of the shadow registers are accessible directly. The following rules apply for transfer of registers into and out of shadows.

- `PUSH.S` and `POP.S`
  W0, W1, W2, W3, SR (DC, N, OV, Z and C bits only) are transferred.
- `DO` instruction
  DOSTART, DOEND, DCOUNT shadows are pushed on loop start, and popped on loop end.

When a byte operation is performed on a working register, only the Least Significant Byte (LSB) of the target register is affected. However, a benefit of memory mapped working registers is that both the Least and Most Significant Bytes can be manipulated through byte wide data memory space accesses.

2.2.1 SOFTWARE STACK POINTER/ FRAME POINTER

The dsPIC® DSC devices contain a software stack. W15 is the dedicated software Stack Pointer (SP), and will be automatically modified by exception processing and subroutine calls and returns. However, W15 can be referenced by any instruction in the same manner as all other W registers. This simplifies the reading, writing and manipulation of the Stack Pointer (e.g., creating stack frames).

Note:	In order to protect against misaligned stack accesses, W15<0> is always clear.

W15 is initialized to 0x0800 during a Reset. The user may reprogram the SP during initialization to any location within data space.

W14 has been dedicated as a Stack Frame Pointer as defined by the `LNK` and `ULNK` instructions. However, W14 can be referenced by any instruction in the same manner as all other W registers.

2.2.2 STATUS REGISTER

The dsPIC DSC core has a 16-bit STATUS Register (SR), the LSB of which is referred to as the SR Low Byte (SRL) and the MSB as the SR High Byte (SRH). See Figure 2-1 for SR layout.

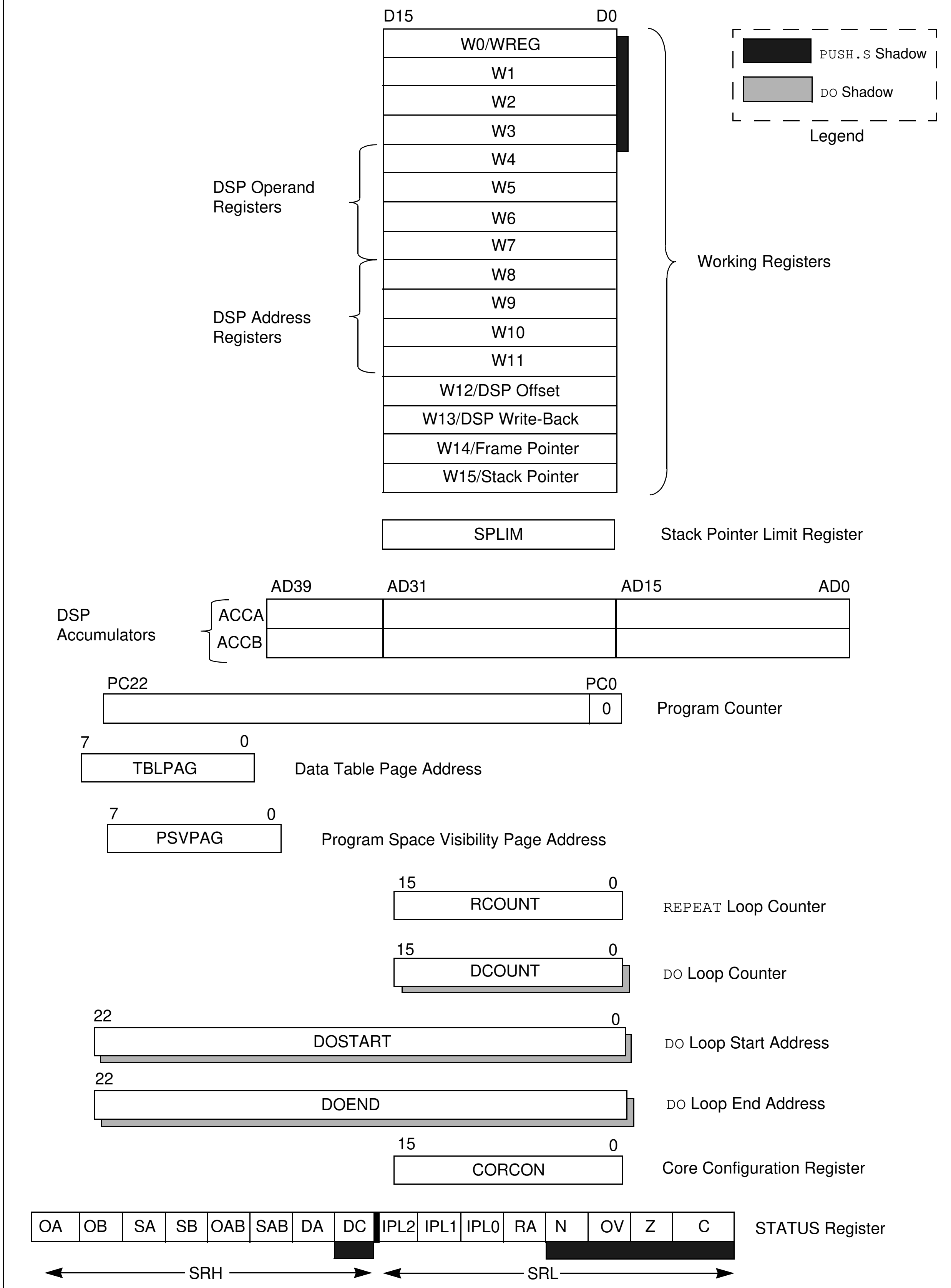
SRL contains all the MCU ALU operation status flags (including the Z bit), as well as the CPU Interrupt Priority Level status bits, IPL<2:0>, and the REPEAT active status bit, RA. During exception processing, SRL is concatenated with the MSB of the PC to form a complete word value which is then stacked.

The upper byte of the STATUS register contains the DSP adder/subtracter status bits, the DO Loop Active bit (DA) and the Digit Carry (DC) status bit.

2.2.3 PROGRAM COUNTER

The Program Counter is 23 bits wide. Bit 0 is always clear. Therefore, the PC can address up to 4M instruction words.

FIGURE 2-1: PROGRAMMER'S MODEL



2.3 Divide Support

The dsPIC DSC devices feature a 16/16-bit signed fractional divide operation, as well as 32/16-bit and 16/16-bit signed and unsigned integer divide operations, in the form of single instruction iterative divides. The following instructions and data sizes are supported:

- `DIVF` – 16/16 signed fractional divide
- `DIV.sd` – 32/16 signed divide
- `DIV.ud` – 32/16 unsigned divide
- `DIV.sw` – 16/16 signed divide
- `DIV.uw` – 16/16 unsigned divide

The 16/16 divides are similar to the 32/16 (same number of iterations), but the dividend is either zero-extended or sign-extended during the first iteration.

The divide instructions must be executed within a `REPEAT` loop. Any other form of execution (e.g., a series of discrete divide instructions) will not function correctly because the instruction flow depends on RCOUNT. The divide instruction does not automatically set up the RCOUNT value, and it must, therefore, be explicitly and correctly specified in the `REPEAT` instruction, as shown in Table 2-2 (`REPEAT` will execute the target instruction {operand value + 1} times). The `REPEAT` loop count must be set up for 18 iterations of the `DIV/DIVF` instruction. Thus, a complete divide operation requires 19 cycles.

Note: The Divide flow is interruptible; however, the user needs to save the context as appropriate.

2.4 DSP Engine

The DSP engine consists of a high-speed 17-bit x 17-bit multiplier, a barrel shifter, and a 40-bit adder/subtracter (with two target accumulators, round and saturation logic).

The DSP engine also has the capability to perform inherent accumulator-to-accumulator operations, which require no additional data. These instructions are `ADD`, `SUB`, and `NEG`.

The DSP engine has various options selected through various bits in the CPU Core Configuration Register (CORCON), as listed below:

- Fractional or integer DSP multiply (IF).
- Signed or unsigned DSP multiply (US).
- Conventional or convergent rounding (RND).
- Automatic saturation on/off for ACCA (SATA).
- Automatic saturation on/off for ACCB (SATB).
- Automatic saturation on/off for writes to data memory (SATDW).
- Accumulator Saturation mode selection (ACCSAT).

Note: For CORCON layout, see Table 3-3.

A block diagram of the DSP engine is shown in Figure 2-2.

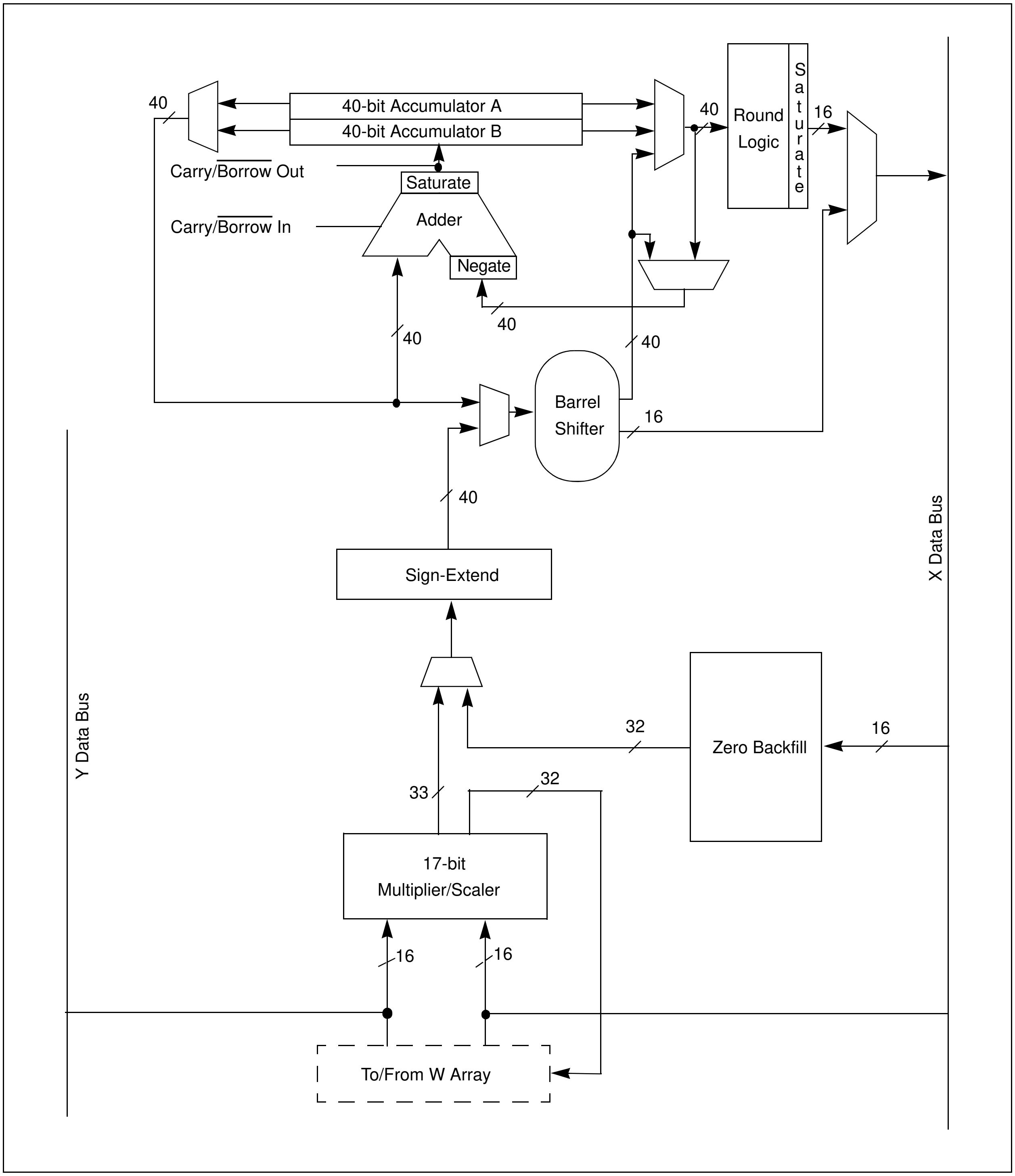
TABLE 2-1: DSP INSTRUCTION SUMMARY

Instruction	Algebraic Operation	ACC WB?
`CLR`	$A = 0$	Yes
`ED`	$A = (x - y)^2$	No
`EDAC`	$A = A + (x - y)^2$	No
`MAC`	$A = A + (x \bullet y)$	Yes
`MAC`	$A = A + x^2$	No
`MOVSAC`	No change in A	Yes
`MPY`	$A = x \bullet y$	No
`MPY.N`	$A = - x \bullet y$	No
`MSC`	$A = A - x \bullet y$	Yes

TABLE 2-2: DIVIDE INSTRUCTIONS

Instruction	Function
`DIVF`	Signed fractional divide: Wm/Wn →W0; Rem →W1
`DIV.sd`	Signed divide: (Wm + 1:Wm)/Wn →W0; Rem →W1
`DIV.ud`	Unsigned divide: (Wm + 1:Wm)/Wn →W0; Rem →W1
`DIV.sw` (or `DIV.s`)	Signed divide: Wm/Wn →W0; Rem →W1
`DIV.uw` (or `DIV.u`)	Unsigned divide: Wm/Wn →W0; Rem →W1

FIGURE 2-2: DSP ENGINE BLOCK DIAGRAM



2.4.1 MULTIPLIER

The 17 x 17-bit multiplier is capable of signed or unsigned operation and can multiplex its output using a scaler to support either 1.31 fractional (Q31) or 32-bit integer results. Unsigned operands are zero-extended into the 17th bit of the multiplier input value. Signed operands are sign-extended into the 17th bit of the multiplier input value. The output of the 17 x 17-bit multiplier/scaler is a 33-bit value, which is sign-extended to 40 bits. Integer data is inherently represented as a signed two's complement value, where the MSB is defined as a sign bit. Generally speaking, the range of an N-bit two's complement integer is -2^{N-1} to $2^{N-1} - 1$. For a 16-bit integer, the data range is -32768 (0x8000) to 32767 (0x7FFF), including '0'. For a 32-bit integer, the data range is -2,147,483,648 (0x8000 0000) to 2,147,483,645 (0x7FFF FFFF).

When the multiplier is configured for fractional multiplication, the data is represented as a two's complement fraction, where the MSB is defined as a sign bit and the radix point is implied to lie just after the sign bit (QX format). The range of an N-bit two's complement fraction with this implied radix point is -1.0 to $(1-2^{1-N})$. For a 16-bit fraction, the Q15 data range is -1.0 (0x8000) to 0.999969482 (0x7FFF), including '0' and has a precision of 3.01518×10^{-5}. In Fractional mode, a 16x16 multiply operation generates a 1.31 product, which has a precision of 4.65661×10^{-10}.

The same multiplier is used to support the MCU multiply instructions, which include integer 16-bit signed, unsigned and mixed sign multiplies.

The `MUL` instruction may be directed to use byte or word-sized operands. Byte operands will direct a 16-bit result, and word operands will direct a 32-bit result to the specified register(s) in the W array.

2.4.2 DATA ACCUMULATORS AND ADDER/SUBTRACTER

The data accumulator consists of a 40-bit adder/subtracter with automatic sign extension logic. It can select one of two accumulators (A or B) as its pre-accumulation source and post-accumulation destination. For the `ADD` and `LAC` instructions, the data to be accumulated or loaded can be optionally scaled via the barrel shifter, prior to accumulation.

2.4.2.1 Adder/Subtracter, Overflow and Saturation

The adder/subtracter is a 40-bit adder with an optional zero input into one side and either true or complement data into the other input. In the case of addition, the carry/borrow input is active high and the other input is true data (not complemented), whereas in the case of subtraction, the carry/borrow input is active low and the other input is complemented. The adder/subtracter generates overflow status bits SA/SB and OA/OB, which are latched and reflected in the STATUS Register.

- Overflow from bit 39: this is a catastrophic overflow in which the sign of the accumulator is destroyed.
- Overflow into guard bits 32 through 39: this is a recoverable overflow. This bit is set whenever all the guard bits are not identical to each other.

The adder has an additional saturation block which controls accumulator data saturation, if selected. It uses the result of the adder, the overflow status bits described above, and the SATA/B (CORCON<7:6>) and ACCSAT (CORCON<4>) mode control bits to determine when and to what value to saturate.

Six STATUS register bits have been provided to support saturation and overflow; they are:

1. OA:
   ACCA overflowed into guard bits
2. OB:
   ACCB overflowed into guard bits
3. SA:
   ACCA saturated (bit 31 overflow and saturation)
   or
   ACCA overflowed into guard bits and saturated (bit 39 overflow and saturation)
4. SB:
   ACCB saturated (bit 31 overflow and saturation)
   or
   ACCB overflowed into guard bits and saturated (bit 39 overflow and saturation)
5. OAB:
   Logical OR of OA and OB
6. SAB:
   Logical OR of SA and SB

The OA and OB bits are modified each time data passes through the adder/subtracter. When set, they indicate that the most recent operation has overflowed into the accumulator guard bits (bits 32 through 39). The OA and OB bits can also optionally generate an arithmetic warning trap when set and the corresponding overflow trap flag enable bit (OVATE, OVBTE) in the INTCON1 register (refer to **Section 5.0 "Interrupts"**) is set. This allows the user to take immediate action, for example, to correct system gain.

The SA and SB bits are modified each time data passes through the adder/subtracter, but can only be cleared by the user. When set, they indicate that the accumulator has overflowed its maximum range (bit 31 for 32-bit saturation, or bit 39 for 40-bit saturation) and will be saturated (if saturation is enabled). When saturation is not enabled, SA and SB default to bit 39 overflow and thus indicate that a catastrophic overflow has occurred. If the COVTE bit in the INTCON1 register is set, SA and SB bits will generate an arithmetic warning trap when saturation is disabled.

The overflow and saturation status bits can optionally be viewed in the Status Register (SR) as the logical OR of OA and OB (in bit OAB), and the logical OR of SA and SB (in bit SAB). This allows programmers to check one bit in the STATUS register to determine if either accumulator has overflowed, or one bit to determine if either accumulator has saturated. This would be useful for complex number arithmetic which typically uses both the accumulators.

The device supports three Saturation and Overflow modes.

1. Bit 39 Overflow and Saturation:
   When bit 39 overflow and saturation occurs, the saturation logic loads the maximally positive 9.31 (0x7FFFFFFFFF) or maximally negative 9.31 value (0x8000000000) into the target accumulator. The SA or SB bit is set and remains set until cleared by the user. This is referred to as 'super saturation' and provides protection against erroneous data or unexpected algorithm problems (e.g., gain calculations).
2. Bit 31 Overflow and Saturation:
   When bit 31 overflow and saturation occurs, the saturation logic then loads the maximally positive 1.31 value (0x007FFFFFFF) or maximally negative 1.31 value (0x0080000000) into the target accumulator. The SA or SB bit is set and remains set until cleared by the user. When this Saturation mode is in effect, the guard bits are not used (so the OA, OB or OAB bits are never set).
3. Bit 39 Catastrophic Overflow
   The bit 39 overflow status bit from the adder is used to set the SA or SB bit, which remain set until cleared by the user. No saturation operation is performed and the accumulator is allowed to overflow (destroying its sign). If the COVTE bit in the INTCON1 register is set, a catastrophic overflow can initiate a trap exception.

2.4.2.2 Accumulator 'Write-Back'

The `MAC` class of instructions (with the exception of `MPY`, `MPY.N`, `ED` and `EDAC`) can optionally write a rounded version of the high word (bits 31 through 16) of the accumulator that is not targeted by the instruction into data space memory. The write is performed across the X bus into combined X and Y address space. The following addressing modes are supported:

1. W13, Register Direct:
   The rounded contents of the non-target accumulator are written into W13 as a 1.15 fraction.
2. [W13]+=2, Register Indirect with Post-Increment:
   The rounded contents of the non-target accumulator are written into the address pointed to by W13 as a 1.15 fraction. W13 is then incremented by 2 (for a word write).

2.4.2.3 Round Logic

The round logic is a combinational block, which performs a conventional (biased) or convergent (unbiased) round function during an accumulator write (store). The Round mode is determined by the state of the RND bit in the CORCON register. It generates a 16-bit, 1.15 data value which is passed to the data space write saturation logic. If rounding is not indicated by the instruction, a truncated 1.15 data value is stored and the least significant word (lsw) is simply discarded.

Conventional rounding takes bit 15 of the accumulator, zero-extends it and adds it to the ACCxH word (bits 16 through 31 of the accumulator). If the ACCxL word (bits 0 through 15 of the accumulator) is between 0x8000 and 0xFFFF (0x8000 included), ACCxH is incremented. If ACCxL is between 0x0000 and 0x7FFF, ACCxH is left unchanged. A consequence of this algorithm is that over a succession of random rounding operations, the value will tend to be biased slightly positive.

Convergent (or unbiased) rounding operates in the same manner as conventional rounding, except when ACCxL equals 0x8000. If this is the case, the Least Significant bit (bit 16 of the accumulator) of ACCxH is examined. If it is '`1`', ACCxH is incremented. If it is '`0`', ACCxH is not modified. Assuming that bit 16 is effectively random in nature, this scheme will remove any rounding bias that may accumulate.

The `SAC` and `SAC.R` instructions store either a truncated (`SAC`) or rounded (`SAC.R`) version of the contents of the target accumulator to data memory, via the X bus (subject to data saturation, see **Section 2.4.2.4 "Data Space Write Saturation"**). Note that for the `MAC` class of instructions, the accumulator write-back operation will function in the same manner, addressing combined MCU (X and Y) data space though the X bus. For this class of instructions, the data is always subject to rounding.

2.4.2.4 Data Space Write Saturation

In addition to adder/subtracter saturation, writes to data space may also be saturated, but without affecting the contents of the source accumulator. The data space write saturation logic block accepts a 16-bit, 1.15 fractional value from the round logic block as its input, together with overflow status from the original source (accumulator) and the 16-bit round adder. These are combined and used to select the appropriate 1.15 fractional value as output to write to data space memory.

If the SATDW bit in the CORCON register is set, data (after rounding or truncation) is tested for overflow and adjusted accordingly. For input data greater than 0x007FFF, data written to memory is forced to the maximum positive 1.15 value, 0x7FFF. For input data less than 0xFF8000, data written to memory is forced to the maximum negative 1.15 value, 0x8000. The Most Significant bit of the source (bit 39) is used to determine the sign of the operand being tested.

If the SATDW bit in the CORCON register is not set, the input data is always passed through unmodified under all conditions.

2.4.3 BARREL SHIFTER

The barrel shifter is capable of performing up to 15-bit arithmetic or logic right shifts, or up to 16-bit left shifts in a single cycle. The source can be either of the two DSP accumulators or the X bus (to support multi-bit shifts of register or memory data).

The shifter requires a signed binary value to determine both the magnitude (number of bits) and direction of the shift operation. A positive value will shift the operand right. A negative value will shift the operand left. A value of 0 will not modify the operand.

The barrel shifter is 40 bits wide, thereby obtaining a 40-bit result for DSP shift operations and a 16-bit result for MCU shift operations. Data from the X bus is presented to the barrel shifter between bit positions 16 to 31 for right shifts, and bit positions 0 to 15 for left shifts.

3.0 MEMORY ORGANIZATION

Note:	This data sheet summarizes features of this group of dsPIC30F devices and is not intended to be a complete reference source. For more information on the CPU, peripherals, register descriptions and general device functionality, refer to the "*dsPIC30F Family Reference Manual*" (DS70046). For more information on the device instruction set and programming, refer to the *"16-bit MCU and DSC Programmer's Reference Manual"* (DS70157).

3.1 Program Address Space

The program address space is 4M instruction words. It is addressable by a 24-bit value from either the 23-bit PC, table instruction Effective Address (EA), or data space EA, when program space is mapped into data space, as defined by Table 3-1. Note that the program space address is incremented by two between successive program words, in order to provide compatibility with data space addressing.

User program space access is restricted to the lower 4M instruction word address range (0x000000 to 0x7FFFFE), for all accesses other than TBLRD/TBLWT, which use TBLPAG<7> to determine user or configuration space access. In Table 3-1, Read/Write instructions, bit 23 allows access to the Device ID, the User ID and the Configuration bits. Otherwise, bit 23 is always clear.

Note:	The address map shown in Figure 3-1 is conceptual, and the actual memory configuration may vary across individual devices depending on available memory.

FIGURE 3-1: PROGRAM SPACE MEMORY MAP FOR dsPIC30F2010

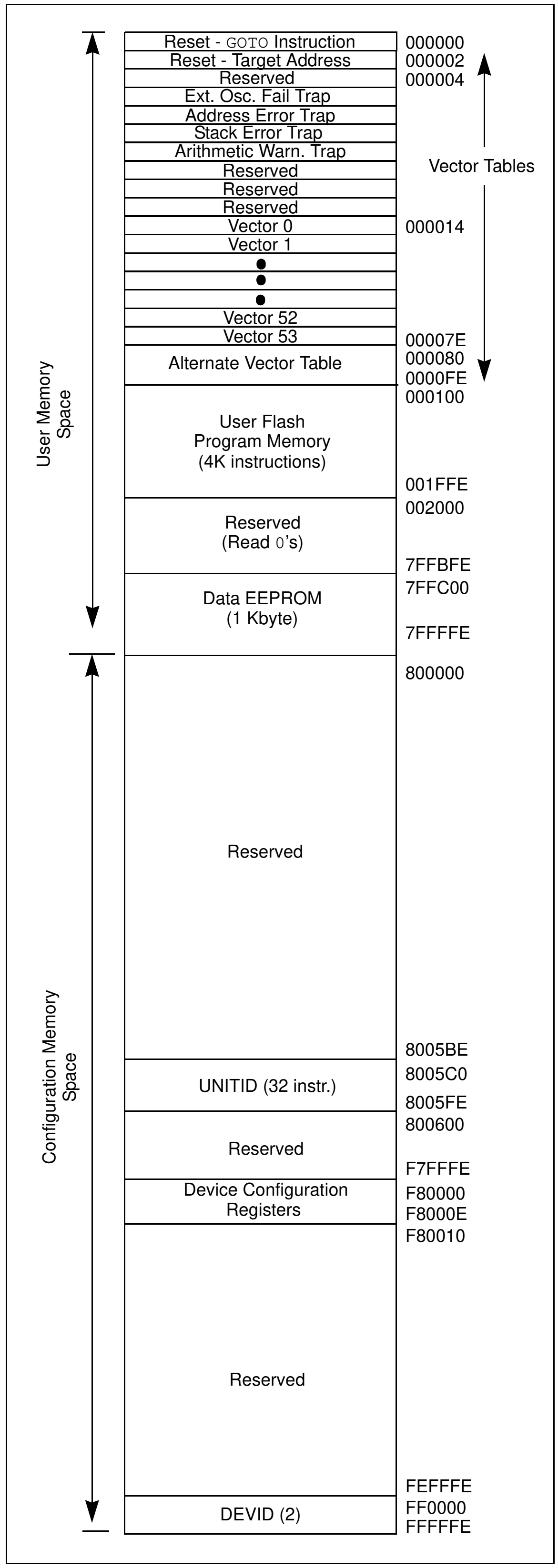
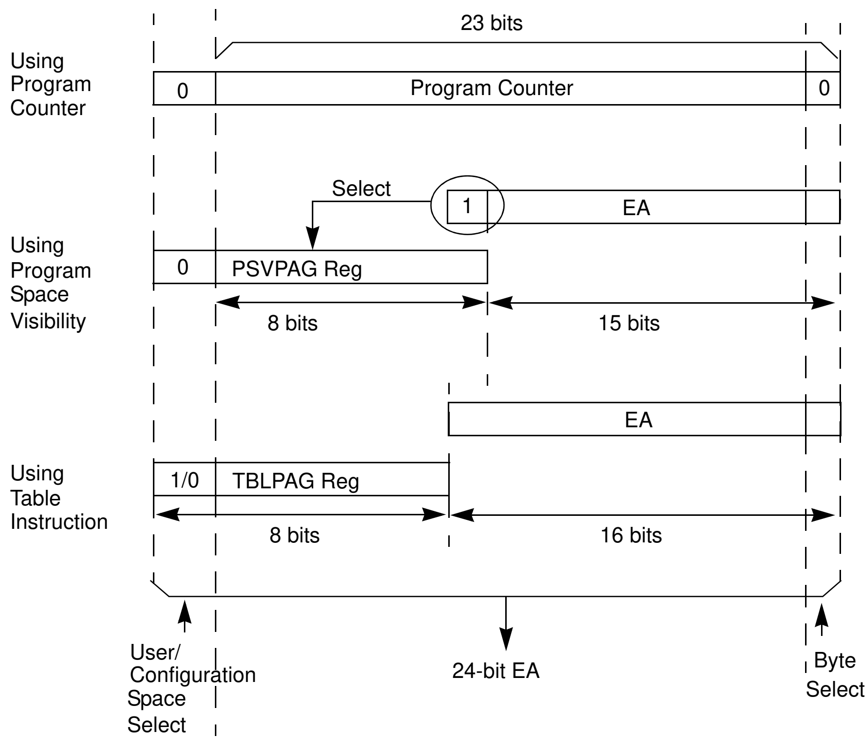


TABLE 3-1: PROGRAM SPACE ADDRESS CONSTRUCTION

Access Type	Access Space	Program Space Address				
		<23>	<22:16>	<15>	<14:1>	<0>
Instruction Access	User	0	PC<22:1>			0
`TBLRD/TBLWT`	User (TBLPAG<7> = 0)	TBLPAG<7:0>		Data EA <15:0>		
`TBLRD/TBLWT`	Configuration (TBLPAG<7> = 1)	TBLPAG<7:0>		Data EA <15:0>		
Program Space Visibility	User	0	PSVPAG<7:0>		Data EA <14:0>	

FIGURE 3-2: DATA ACCESS FROM PROGRAM SPACE ADDRESS GENERATION



Note: Program Space Visibility cannot be used to access bits <23:16> of a word in program memory.

3.1.1 DATA ACCESS FROM PROGRAM MEMORY USING TABLE INSTRUCTIONS

This architecture fetches 24-bit wide program memory. Consequently, instructions are always aligned. However, as the architecture is modified Harvard, data can also be present in program space.

There are two methods by which program space can be accessed: via special table instructions, or through the remapping of a 16K word program space page into the upper half of data space (see **Section 3.1.2 "Data Access from Program Memory Using Program Space Visibility"**). The `TBLRDL` and `TBLWTL` instructions offer a direct method of reading or writing the lsw of any address within program space, without going through data space. The `TBLRDH` and `TBLWTH` instructions are the only method whereby the upper 8 bits of a program space word can be accessed as data.

The PC is incremented by two for each successive 24-bit program word. This allows program memory addresses to directly map to data space addresses. Program memory can thus be regarded as two 16-bit word wide address spaces, residing side by side, each with the same address range. `TBLRDL` and `TBLWTL` access the space which contains the least significant data word, and `TBLRDH` and `TBLWTH` access the space which contains the Most Significant data Byte.

Figure 3-2 shows how the EA is created for table operations and data space accesses (PSV = `1`). Here, P<23:0> refers to a program space word, whereas D<15:0> refers to a data space word.

A set of Table Instructions are provided to move byte or word-sized data to and from program space.

1. **`TBLRDL:`** Table Read Low
   Word: Read the least significant word of the program address;
   P<15:0> maps to D<15:0>.
   Byte: Read one of the LSBs of the program address;
   P<7:0> maps to the destination byte when byte select = `0`;
   P<15:8> maps to the destination byte when byte select = `1`.
2. **`TBLWTL:`** Table Write Low (refer to **Section 6.0 "Flash Program Memory"** for details on Flash Programming).
3. **`TBLRDH:`** Table Read High
   Word: Read the most significant word of the program address;
   P<23:16> maps to D<7:0>; D<15:8> always be = `0`.
   Byte: Read one of the MSBs of the program address;
   P<23:16> maps to the destination byte when byte select = `0`;
   The destination byte will always be = `0` when byte select = `1`.
4. **`TBLWTH:`** Table Write High (refer to **Section 6.0 "Flash Program Memory"** for details on Flash Programming).

FIGURE 3-3: PROGRAM DATA TABLE ACCESS (LEAST SIGNIFICANT WORD)

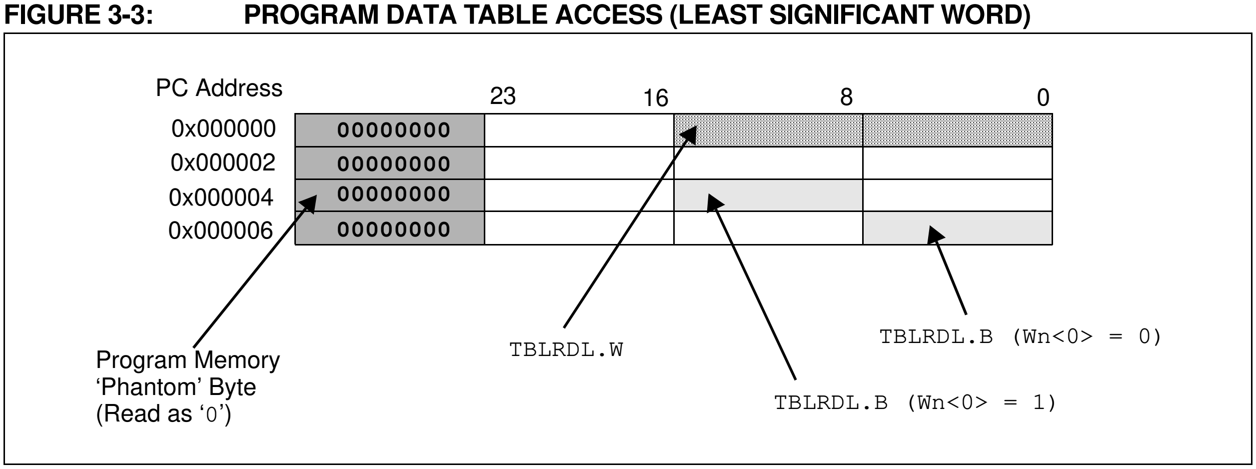
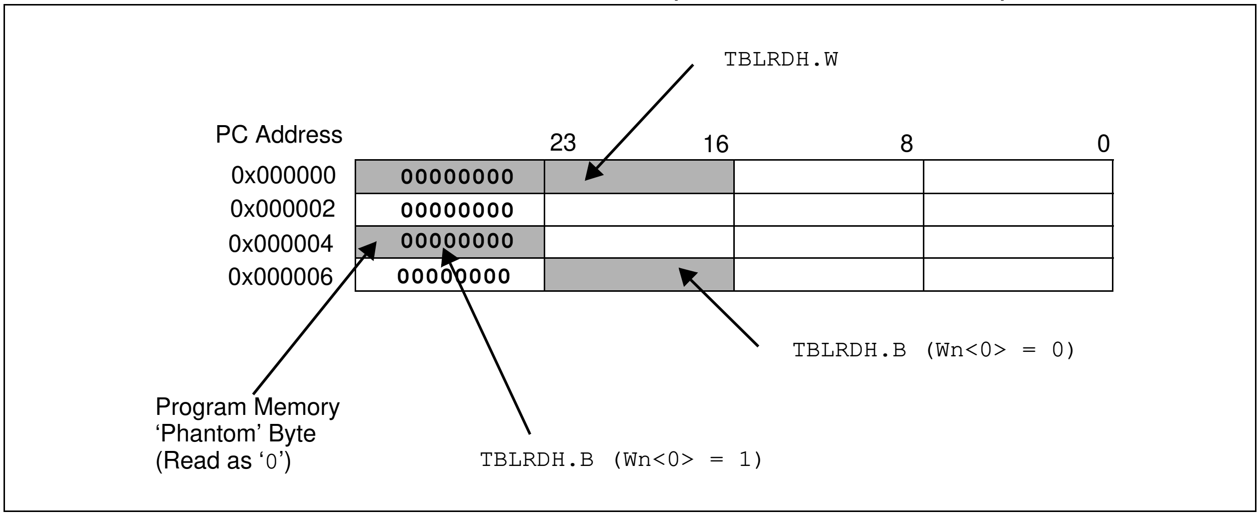


FIGURE 3-4: PROGRAM DATA TABLE ACCESS (MOST SIGNIFICANT BYTE)



3.1.2 DATA ACCESS FROM PROGRAM MEMORY USING PROGRAM SPACE VISIBILITY

The upper 32 Kbytes of data space may optionally be mapped into any 16K word program space page. This provides transparent access of stored constant data from X data space, without the need to use special instructions (i.e., `TBLRDL/H`, `TBLWTL/H` instructions).

Program space access through the data space occurs if the MSb of the data space EA is set and program space visibility is enabled, by setting the PSV bit in the Core Control register (CORCON). The functions of CORCON are discussed in **Section 2.4 "DSP Engine"**.

Data accesses to this area add an additional cycle to the instruction being executed, since two program memory fetches are required.

Note that the upper half of addressable data space is always part of the X data space. Therefore, when a DSP operation uses program space mapping to access this memory region, Y data space should typically contain state (variable) data for DSP operations, whereas X data space should typically contain coefficient (constant) data.

Although each data space address, 0x8000 and higher, maps directly into a corresponding program memory address (see Figure 3-5), only the lower 16-bits of the 24-bit program word are used to contain the data. The upper 8 bits should be programmed to force an illegal instruction to maintain machine robustness. Refer to the "*16-bit MCU and DSC Programmer's Reference Manual*" (DS70157) for details on instruction encoding.

Note that by incrementing the PC by 2 for each program memory word, the Least Significant 15 bits of data space addresses directly map to the Least Significant 15 bits in the corresponding program space addresses. The remaining bits are provided by the Program Space Visibility Page register, PSVPAG<7:0>, as shown in Figure 3-5.

Note:	PSV access is temporarily disabled during table reads/writes.

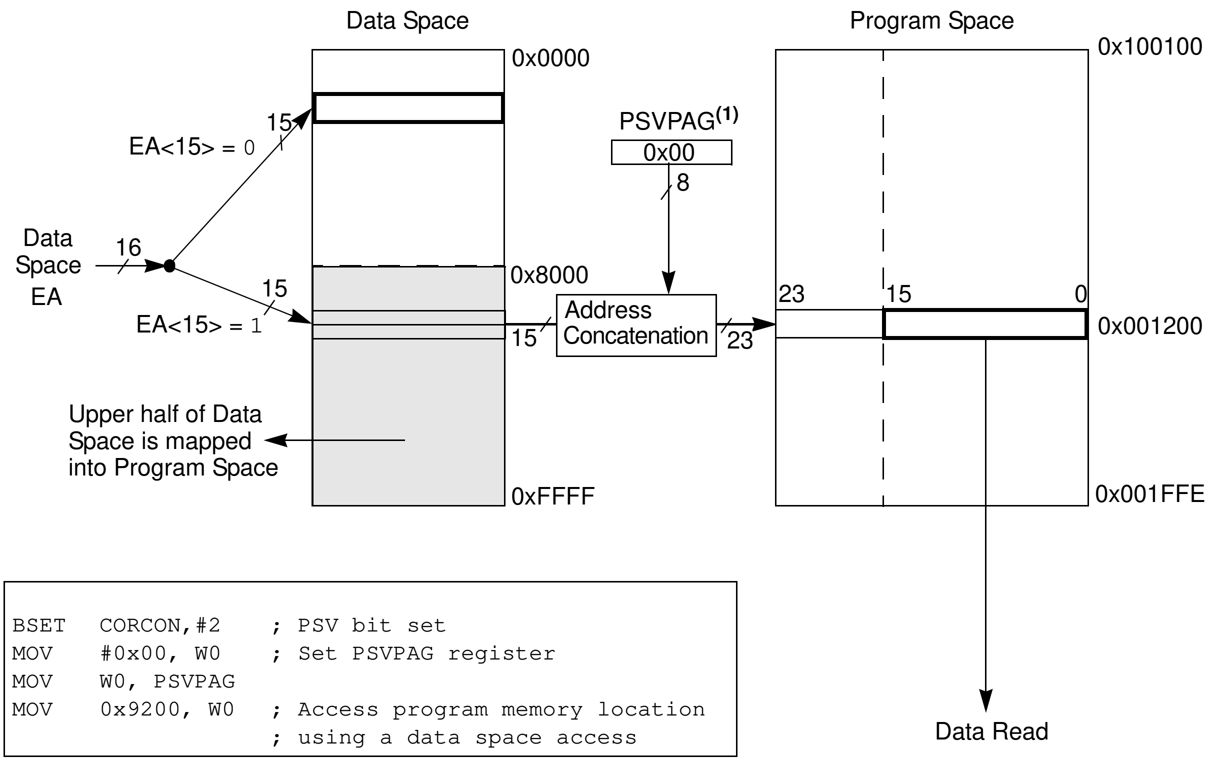
For instructions that use PSV which are executed outside a `REPEAT` loop:

- The following instructions will require one instruction cycle in addition to the specified execution time:
  - `MAC` class of instructions with data operand prefetch
  - `MOV` instructions
  - `MOV.D` instructions
- All other instructions will require two instruction cycles in addition to the specified execution time of the instruction.

For instructions that use PSV which are executed inside a `REPEAT` loop:

- The following instances will require two instruction cycles in addition to the specified execution time of the instruction:
  - Execution in the first iteration
  - Execution in the last iteration
  - Execution prior to exiting the loop due to an interrupt
  - Execution upon re-entering the loop after an interrupt is serviced
- Any other iteration of the `REPEAT` loop will allow the instruction, accessing data using PSV, to execute in a single cycle.

FIGURE 3-5: DATA SPACE WINDOW INTO PROGRAM SPACE OPERATION




```
BSET   CORCON,#2     ; PSV bit set
MOV    #0x00, W0     ; Set PSVPAG register
MOV    W0, PSVPAG
MOV    0x9200, W0    ; Access program memory location
                     ; using a data space access
```

Note: PSVPAG is an 8-bit register, containing bits <22:15> of the program space address (i.e., it defines the page in program space to which the upper half of data space is being mapped).

3.2 Data Address Space

The core has two data spaces. The data spaces can be considered either separate (for some DSP instructions), or as one unified linear address range (for MCU instructions). The data spaces are accessed using two Address Generation Units (AGUs) and separate data paths.

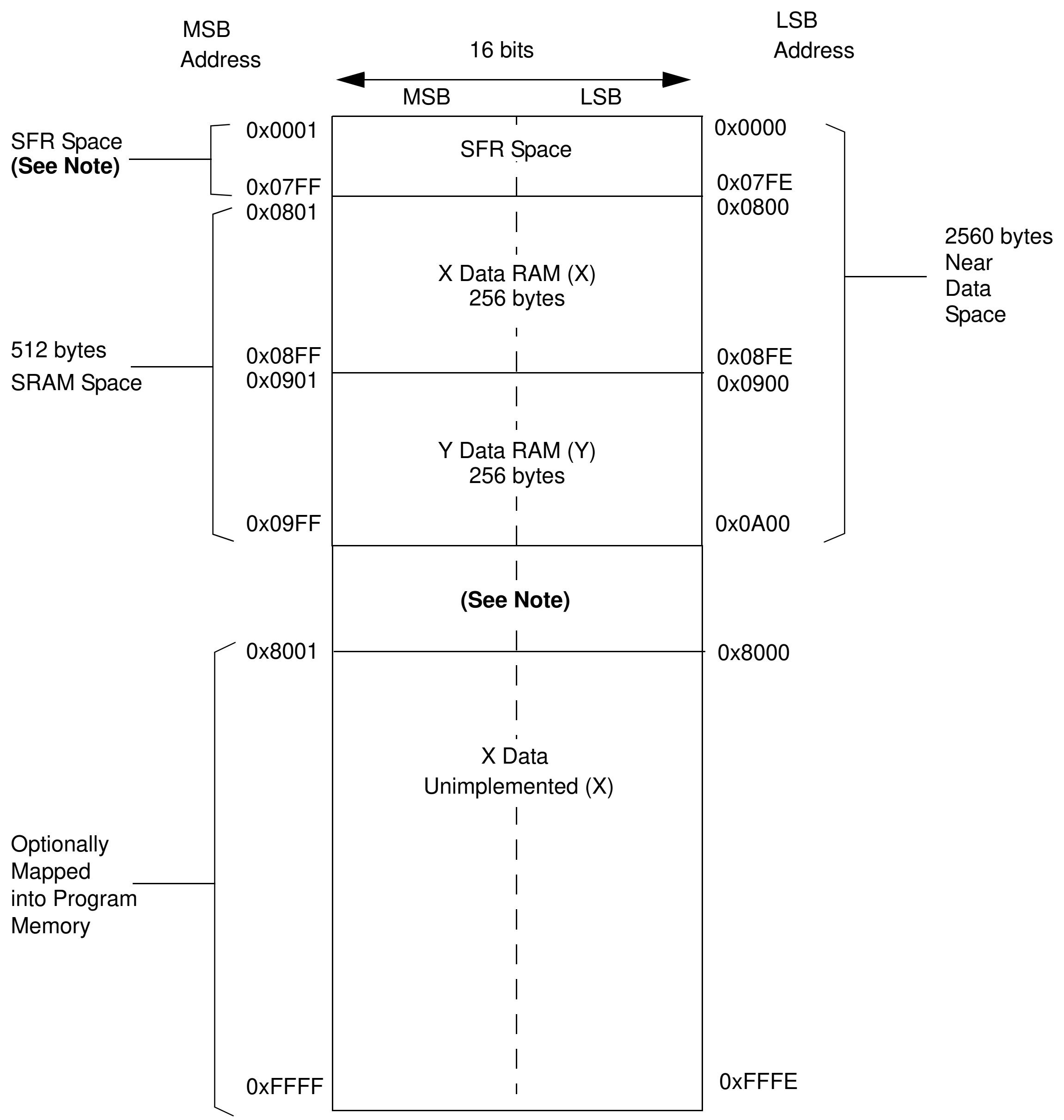
3.2.1 DATA SPACE MEMORY MAP

The data space memory is split into two blocks, X and Y data space. A key element of this architecture is that Y space is a subset of X space, and is fully contained within X space. In order to provide an apparent linear addressing space, X and Y spaces have contiguous addresses.

When executing any instruction other than one of the `MAC` class of instructions, the X block consists of the 256 byte data address space (including all Y addresses). When executing one of the `MAC` class of instructions, the X block consists of the 256 bytes data address space excluding the Y address block (for data reads only). In other words, all other instructions regard the entire data memory as one composite address space. The `MAC` class instructions extract the Y address space from data space and address it using EAs sourced from W10 and W11. The remaining X data space is addressed using W8 and W9. Both address spaces are concurrently accessed only with the `MAC` class instructions.

A data space memory map is shown in Figure 3-6.

FIGURE 3-6: DATA SPACE MEMORY MAP



Note: Unimplemented SFR or SRAM locations read as ‘0’.

FIGURE 3-7: DATA SPACE FOR MCU AND DSP (MAC CLASS) INSTRUCTIONS