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## Contact us

Tel: +86-755-8981 8866 Fax: +86-755-8427 6832
Email \& Skype: info@chipsmall.com Web: www.chipsmall.com Address: A1208, Overseas Decoration Building, \#122 Zhenhua RD., Futian, Shenzhen, China

## FEATURES

- High-performance, $\mathbf{E}^{2}$ CMOS 3.3-V \& 5-V CPLD families
- Flexible architecture for rapid logic designs
- Excellent First-Time-Fit ${ }^{\mathrm{TM}}$ and refit feature
- SpeedLocking ${ }^{\text {TM }}$ performance for guaranteed fixed timing
- Central, input and output switch matrices for $100 \%$ routability and $100 \%$ pin-out retention
- High speed
- 5.0ns $\mathrm{t}_{\mathrm{PD}}$ Commercial and 7.5ns $\mathrm{t}_{\mathrm{PD}}$ Industrial
$-182 \mathrm{MHz} \mathrm{f} \mathrm{CNT}$
$\mathbf{3 2}$ to $\mathbf{5 1 2}$ macrocells; $\mathbf{3 2}$ to $\mathbf{7 6 8}$ registers
- 44 to 388 pins in PLCC, PQFP, TQFP, BGA, fpBGA and caBGA packages
- Flexible architecture for a wide range of design styles
- D/T registers and latches
- Synchronous or asynchronous mode
- Dedicated input registers
- Programmable polarity
- Reset/ preset swapping
- Advanced capabilities for easy system integration
- 3.3-V \& 5-V JEDEC-compliant operations
- JTAG (IEEE 1149.1) compliant for boundary scan testing
- 3.3-V \& 5-V JTAG in-system programming
- PCI compliant (-5/-55/-6/-65/-7/-10/-12 speed grades)
- Safe for mixed supply voltage system designs
- Programmable pull-up or Bus-Friendly ${ }^{\mathrm{TM}}$ inputs and I/Os
- Hot-socketing
- Programmable security bit
- Individual output slew rate control

Advanced E ${ }^{2}$ CMOS process provides high-performance, cost-effective solutions

- Lead-free package options

Table 1. ispMACH 4A Device Features

| 3.3 V Devices |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Feature | M4 A3-32 | M4 A3-64 | M4 A3-96 | M4 A3-128 | M4 A3-192 | M4 A3-256 | M4 A3-384 | M4 A3-512 |
| Macrocells | 32 | 64 | 96 | 128 | 192 | 256 | 384 | 512 |
| User I/O options | 32 | $32 / 64$ | 48 | 64 | 96 | 128/160/192 | $160 / 192$ | 160/192/256 |
| $\mathrm{t}_{\mathrm{PD}}$ ( ns ) | 5.0 | 5.5 | 5.5 | 5.5 | 6.0 | 5.5 | 6.5 | 7.5 |
| $\mathrm{f}_{\mathrm{CNT}}(\mathrm{MHz})$ | 182 | 167 | 167 | 167 | 160 | 167 | 154 | 125 |
| $\mathrm{t}_{\cos }(\mathrm{ns})$ | 4.0 | 4.0 | 4.0 | 4.0 | 4.5 | 4.0 | 4.5 | 5.5 |
| $\mathrm{t}_{\text {SS }}$ ( ns ) | 3.0 | 3.5 | 3.5 | 3.5 | 3.5 | 3.5 | 3.5 | 5.0 |
| Static Power (mA) | 20 | 25/52 | 40 | 55 | 85 | 110/150 | $149 / 155$ | 179 |
| JTAG Compliant | Yes | Yes | Yes | Yes | Yes | Yes | Yes | Yes |
| PCI Compliant | Yes | Yes | Yes | Yes | Yes | Yes | Yes | Yes |


| 5 V Devices |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Feature | M4 A5-32 | M4 A5-64 | M4 A5-96 | M4 A5-128 | M4 A5-192 | M4 A5-256 |
| Macrocells | 32 | 64 | 96 | 128 | 192 | 256 |
| User I/O options | 32 | 32 | 48 | 64 | 96 | 128 |
| $\mathrm{t}_{\mathrm{PD}}$ ( ns ) | 5.0 | 5.5 | 5.5 | 5.5 | 6.0 | 6.5 |
| $\mathrm{f}_{\mathrm{CNT}}(\mathrm{MHz})$ | 182 | 167 | 167 | 167 | 160 | 154 |
| $\mathrm{t}_{\mathrm{COS}}(\mathrm{ns})$ | 4.0 | 4.0 | 4.0 | 4.0 | 4.5 | 5.0 |
| $\mathrm{t}_{\text {SS }}$ ( ns ) | 3.0 | 3.5 | 3.5 | 3.5 | 3.5 | 3.5 |
| Static Power (mA) | 20 | 25 | 40 | 55 | 74 | 110 |
| JTAG Compliant | Yes | Yes | Yes | Yes | Yes | Yes |
| PCI Compliant | Yes | Yes | Yes | Yes | Yes | Yes |

## GENERAL DESCRIPTION

The ispMACH ${ }^{\mathrm{TM}} 4 \mathrm{~A}$ family from Lattice offers an exceptionally flexible architecture and delivers a superior Complex Programmable Logic Device (CPLD) solution of easy-to-use silicon products and software tools. The overall benefits for users are a guaranteed and predictable CPLD solution, faster time-to-market, greater flexibility and lower cost. The ispMACH 4A devices offer densities ranging from 32 to 512 macrocells with $100 \%$ utilization and $100 \%$ pin-out retention. The ispMACH 4A families offer 5-V (M4A5xxx ) and 3.3-V (M4A3-xxx) operation.
ispMACH 4A products are 5-V or 3.3-V in-system programmable through the JTAG (IEEE Std. 1149.1) interface. JTAG boundary scan testing also allows product testability on automated test equipment for device connectivity.

All ispMACH 4A family members deliver First-Time-Fit and easy system integration with pin-out retention after any design change and refit. For both $3.3-\mathrm{V}$ and $5-\mathrm{V}$ operation, ispMACH 4 A products can deliver guaranteed fixed timing as fast as $5.0 \mathrm{~ns}_{\mathrm{PD}}$ and $182 \mathrm{MHz} \mathrm{f}_{\mathrm{CNT}}$ through the SpeedLocking feature when using up to 20 product terms per output (Table 2).

Table 2. ispMACH 4A Speed Grades

| Device | Speed Grade |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | -5 | -55 | -6 | -65 | -7 | -10 | -12 | -14 |
| $\begin{aligned} & \text { M4 A3-32 } \\ & \text { M4 A5-32 } \end{aligned}$ | C |  |  |  | C, I | C, I | I |  |
| M4 A3-64 /32 <br> M4 A5-64 32 |  | C |  |  | C, I | C, I | I |  |
| M4A3-64/64 |  | C |  |  | C, I | C, I | I |  |
| $\begin{aligned} & \text { M4 A3-96 } \\ & \text { M4 A5-96 } \end{aligned}$ |  | C |  |  | C, I | C, I | I |  |
| M4A3-128 <br> M4 A5-128 |  | C |  |  | C, I | C, I | I |  |
| M4A3-192 <br> M4A5-192 |  |  | C |  | C, I | C, I | I |  |
| M4 A3-256/128 |  | C |  | C | C, I | C, I | I |  |
| M4 A5-256/128 |  |  |  | C | C | C, I | I |  |
| M4 A3-256/192 <br> M4 A3-256 /160 |  |  |  |  | C | C, I | I |  |
| M4 A3-384 |  |  |  | C |  | C, I | C, I | 1 |
| M4 A3-512 |  |  |  |  | C | C, I | C, I | I |

## Note:

1. $C=$ Commercial,$\quad I=$ Industrial

The ispMACH 4A family offers 20 density-I/O combinations in Thin Quad Flat Pack (TQFP), Plastic Quad Flat Pack (PQFP), Plastic Leaded Chip Carrier (PLCC), Ball Grid Array (BGA), fine-pitch BGA (fpBGA), and chip-array BGA (caBGA) packages ranging from 44 to 388 pins (Table 3). It also offers I/O safety features for mixed-voltage designs so that the $3.3-\mathrm{V}$ devices can accept $5-\mathrm{V}$ inputs, and $5-\mathrm{V}$ devices do not overdrive $3.3-\mathrm{V}$ inputs. Additional features include Bus-Friendly inputs and I/Os, a programmable power-down mode for extra power savings and individual output slew rate control for the highest speed transition or for the lowest noise transition.

Table 3. ispMACH 4A Package and I/O Options (Number of I/Os and dedicated inputs in Table)

| 3.3 V Devices |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Package | M4 A3-32 | M4 A3-64 | M4 A3-96 | M4 A3-128 | M4 A3-192 | M4 A3-256 | M4 A3-384 | M4 A3-512 |
| 44-pin PLCC | $32+2$ | $32+2$ |  |  |  |  |  |  |
| 44-pin TQFP | $32+2$ | $32+2$ |  |  |  |  |  |  |
| 48-pin TQFP | $32+2$ | $32+2$ |  |  |  |  |  |  |
| 100-pin TQFP |  | $64+6$ | $48+8$ | $64+6$ |  |  |  |  |
| 100-pin PQFP |  |  |  | $64+6$ |  |  |  |  |
| 100-ball caBGA |  |  |  | $64+6$ |  |  |  |  |
| 144-pin TQFP |  |  |  |  | $96+16$ |  |  |  |
| 144-ball fpBGA |  |  |  |  | $96+16$ |  |  |  |
| 208-pin PQFP |  |  |  |  |  | $128+14,160$ | 160 | 160 |
| 256-ball fpBGA |  |  |  |  |  | $128+14,192$ | 192 | 192 |
| 256-ball BGA |  |  |  |  |  | $128+14$ | 192 |  |
| 388-ball fpBGA |  |  |  |  |  |  |  | 256 |


| 5 V Devices |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Package | M4 A5-32 | M4 A5-64 | M4 A5-96 | M4 A5-128 | M4 A5-192 | M4 A5-256 |
| 44 -pin PLCC | $32+2$ | $32+2$ |  |  |  |  |
| 44 -pin TQFP | $32+2$ | $32+2$ |  |  |  |  |
| 48 -pin TQFP | $32+2$ | $32+2$ |  |  |  |  |
| 100 -pin TQFP |  |  | $48+8$ | $64+6$ |  |  |
| 100 -pin PQFP |  |  |  |  |  |  |
| 144 -pin TQFP |  |  |  | $96+16$ |  |  |
| 208 -pin PQFP |  |  |  |  |  |  |

## FUNCTIONAL DESCRIPTION

The fundamental architecture of ispMACH 4A devices (Figure 1) consists of multiple, optimized PAL ${ }^{\circledR}$ blocks interconnected by a central switch matrix. The central switch matrix allows communication between PAL blocks and routes inputs to the PAL blocks. Together, the PAL blocks and central switch matrix allow the logic designer to create large designs in a single device instead of having to use multiple devices.

The key to being able to make effective use of these devices lies in the interconnect schemes. In the ispMACH 4A architecture, the macrocells are flexibly coupled to the product terms through the logic allocator, and the I/O pins are flexibly coupled to the macrocells due to the output switch matrix. In addition, more input routing options are provided by the input switch matrix. These resources provide the flexibility needed to fit designs efficiently.


Figure 1. ispMACH 4A Block Diagram and PAL Block Structure

## Notes:

1. 16 for ispMACH $4 A$ devices with $1: 1$ macrocell-I/ O cell ratio (see next page).
2. Block clocks do not go to I/ O cells in M4A (3,5)-32/ 32 .
3. M4A (3,5)-192, M4A (3,5)-256, M4A 3-384, and M4A3-512 haze dediated clock pins uhich cannot be used as inputs and do not connect to the central switch matrix.

Table 4. Architectural Summary of ispMACH 4A devices

|  | ispMACH 4 A Devices |  |
| :---: | :---: | :---: |
|  | $\begin{gathered} \text { M4 A3-64 } 32, \text { M4 A5-64 } 32 \\ \text { M4 A3-96 } / 48, \text { M4 A5-96 /48 } \\ \text { M4A3-128 /64, M4 A5-128 /64 } \\ \text { M4A3-192 /96, M4 A5-192/96 } \\ \text { M4A3-256 } 128, \text { M4 A5-256/128 } \\ \text { M4A3-384 } \\ \text { M4A3-512 } \end{gathered}$ | $\begin{gathered} \text { M4 A3-32 /32 } \\ \text { M4 A5-32 /32 } \\ \text { M4 A3-64 /64 } \\ \text { M4 A3-256 /160 } \\ \text { M4 A3-256/192 } \end{gathered}$ |
| Macrocell-I/O Cell Ratio | 2:1 | 1:1 |
| Input Switch Matrix | Yes | Yes ${ }^{1}$ |
| Input Registers | Yes | No |
| Central Switch Matrix | Yes | Yes |
| Output Switch Matrix | Yes | Yes |

The Macrocell-I/O cell ratio is defined as the number of macrocells versus the number of I/O cells internally in a PAL block (Table 4).
The central switch matrix takes all dedicated inputs and signals from the input switch matrices and routes them as needed to the PAL blocks. Feedback signals that return to the same PAL block still must go through the central switch matrix. This mechanism ensures that PAL blocks in ispMACH 4A devices communicate with each other with consistent, predictable delays.
The central switch matrix makes a ispMACH 4A device more advanced than simply several PAL devices on a single chip. It allows the designer to think of the device not as a collection of blocks, but as a single programmable device; the software partitions the design into PAL blocks through the central switch matrix so that the designer does not have to be concerned with the internal architecture of the device.

## Each PAL block consists of:

- Product-term array
- Logic allocator
- Macrocells
- Output switch matrix
- I/O cells
- Input switch matrix
- Clock generator


## Notes:

1. M4A 3-64/ 64 internal switch matrix functionality embedded in central switch matrix.

## Product-Term Array

The product-term array consists of a number of product terms that form the basis of the logic being implemented. The inputs to the AND gates come from the central switch matrix (Table 5), and are provided in both true and complement forms for efficient logic implementation.

Table 5. PAL Block Inputs

| Device | Number of Inputs to PAL Block |
| :---: | :---: |
| M4 A3-32 32 and M4 A5-32 32 | 33 |
| M4 A3-64 32 and M4 A5-64 32 | 33 |
| M4 A3-64/64 | 33 |
| M4 A3-96/48 and M4 A5-96/48 | 33 |
| M4 A3-128/64 and M4 A5-128/64 | 33 |
| M4 A3-192/96 and M4 A5-192/96 | 34 |
| M4 A3-256/128 and M4 A5-256/128 | 34 |
| M4 A3-256/160 and M4 A3-256/192 | 36 |
| M4 A3-384 | 36 |
| M4 A3-512 | 36 |

## Logic Allocator

Within the logic allocator, product terms are allocated to macrocells in "product term clusters." The availability and distribution of product term clusters are automatically considered by the software as it fits functions within a PAL block. The size of a product term cluster has been optimized to provide high utilization of product terms, making complex functions using many product terms possible. Yet when few product terms are used, there will be a minimal number of unused-or wasted-product terms left over. The product term clusters available to each macrocell within a PAL block are shown in Tables 6 and 7.

Each product term cluster is associated with a macrocell. The size of a cluster depends on the configuration of the associated macrocell. When the macrocell is used in synchronous mode
(Figure 2a), the basic cluster has 4 product terms. When the associated macrocell is used in asynchronous mode (Figure 2b), the cluster has 2 product terms. Note that if the product term cluster is routed to a different macrocell, the allocator configuration is not determined by the mode of the macrocell actually being driven. The configuration is always set by the mode of the macrocell that the cluster will drive if not routed away, regardless of the actual routing.

In addition, there is an extra product term that can either join the basic cluster to give an extended cluster, or drive the second input of an exclusive-OR gate in the signal path. If included with the basic cluster, this provides for up to 20 product terms on a synchronous function that uses four extended 5-product-term clusters. A similar asynchronous function can have up to 18 product terms.
When the extra product term is used to extend the cluster, the value of the second XOR input can be programmed as a 0 or a 1, giving polarity control. The possible configurations of the logic allocator are shown in Figures 3 and 4.

Table 6. Logic Allocator for All ispMACH 4A Devices (except M4A(3,5)-32/32)

| Output Macrocell | Available Clusters | Output Macrocell | Available Clusters |
| :---: | :---: | :---: | :---: |
| $\mathrm{M}_{0}$ | $\mathrm{C}_{0}, \mathrm{C}_{1}, \mathrm{C}_{2}$ | $\mathrm{M}_{8}$ | $\mathrm{C}_{7}, \mathrm{C}_{8}, \mathrm{C}_{9}, \mathrm{C}_{10}$ |
| $\mathrm{M}_{1}$ | $\mathrm{C}_{0}, \mathrm{C}_{1}, \mathrm{C}_{2}, \mathrm{C}_{3}$ | $\mathrm{M}_{9}$ | $\mathrm{C}_{8}, \mathrm{C}_{9}, \mathrm{C}_{10}, \mathrm{C}_{11}$ |
| $\mathrm{M}_{2}$ | $\mathrm{C}_{1}, \mathrm{C}_{2}, \mathrm{C}_{3}, \mathrm{C}_{4}$ | $\mathrm{M}_{10}$ | $\mathrm{C}_{9}, \mathrm{C}_{10}, \mathrm{C}_{11}, \mathrm{C}_{12}$ |
| $\mathrm{M}_{3}$ | $\mathrm{C}_{2}, \mathrm{C}_{3}, \mathrm{C}_{4}, \mathrm{C}_{5}$ | $\mathrm{M}_{11}$ | $\mathrm{C}_{10}, \mathrm{C}_{11}, \mathrm{C}_{12}, \mathrm{C}_{13}$ |
| $\mathrm{M}_{4}$ | $\mathrm{C}_{3}, \mathrm{C}_{4}, \mathrm{C}_{5}, \mathrm{C}_{6}$ | $\mathrm{M}_{12}$ | $\mathrm{C}_{11}, \mathrm{C}_{12}, \mathrm{C}_{13}, \mathrm{C}_{14}$ |
| $\mathrm{M}_{5}$ | $\mathrm{C}_{4}, \mathrm{C}_{5}, \mathrm{C}_{6}, \mathrm{C}_{7}$ | $\mathrm{M}_{13}$ | $\mathrm{C}_{12}, \mathrm{C}_{13}, \mathrm{C}_{14}, \mathrm{C}_{15}$ |
| $\mathrm{M}_{6}$ | $\mathrm{C}_{5}, \mathrm{C}_{6}, \mathrm{C}_{7}, \mathrm{C}_{8}$ | $\mathrm{M}_{14}$ | $\mathrm{C}_{13}, \mathrm{C}_{14}, \mathrm{C}_{15}$ |
| $\mathrm{M}_{7}$ | $\mathrm{C}_{6}, \mathrm{C}_{7}, \mathrm{C}_{8}, \mathrm{C}_{9}$ | $\mathrm{M}_{15}$ | $C_{14}, C_{15}$ |

Table 7. Logic Allocator for M4A(3,5)-32/32

| Output Macrocell | Available Clusters | Output Macrocell | Available Clusters |
| :---: | :---: | :---: | :---: |
| $\mathrm{M}_{0}$ | $\mathrm{C}_{0}, \mathrm{C}_{1}, \mathrm{C}_{2}$ | $\mathrm{M}_{8}$ | $\mathrm{C}_{8}, \mathrm{C}_{9}, \mathrm{C}_{10}$ |
| $\mathrm{M}_{1}$ | $\mathrm{C}_{0}, \mathrm{C}_{1}, \mathrm{C}_{2}, \mathrm{C}_{3}$ | $\mathrm{M}_{9}$ | $\mathrm{C}_{8}, \mathrm{C}_{9}, \mathrm{C}_{10}, \mathrm{C}_{11}$ |
| $\mathrm{M}_{2}$ | $\mathrm{C}_{1}, \mathrm{C}_{2}, \mathrm{C}_{3}, \mathrm{C}_{4}$ | $\mathrm{M}_{10}$ | $\mathrm{C}_{9}, \mathrm{C}_{10}, \mathrm{C}_{11}, \mathrm{C}_{12}$ |
| $\mathrm{M}_{3}$ | $\mathrm{C}_{2}, \mathrm{C}_{3}, \mathrm{C}_{4}, \mathrm{C}_{5}$ | $\mathrm{M}_{11}$ | $\mathrm{C}_{10}, \mathrm{C}_{11}, \mathrm{C}_{12}, \mathrm{C}_{13}$ |
| $\mathrm{M}_{4}$ | $\mathrm{C}_{3}, \mathrm{C}_{4}, \mathrm{C}_{5}, \mathrm{C}_{6}$ | $\mathrm{M}_{12}$ | $\mathrm{C}_{11}, \mathrm{C}_{12}, \mathrm{C}_{13}, \mathrm{C}_{14}$ |
| $\mathrm{M}_{5}$ | $\mathrm{C}_{4}, \mathrm{C}_{5}, \mathrm{C}_{6}, \mathrm{C}_{7}$ | $\mathrm{M}_{13}$ | $\mathrm{C}_{12}, \mathrm{C}_{13}, \mathrm{C}_{14}, \mathrm{C}_{15}$ |
| $\mathrm{M}_{6}$ | $\mathrm{C}_{5}, \mathrm{C}_{6}, \mathrm{C}_{7}$ | $\mathrm{M}_{14}$ | $\mathrm{C}_{13}, \mathrm{C}_{14}, C_{15}$ |
| $\mathrm{M}_{7}$ | $\mathrm{C}_{6}, \mathrm{C}_{7}$ | $\mathrm{M}_{15}$ | $\mathrm{C}_{14}, C_{15}$ |



17466G-005
a. Synchronous Mode


Figure 2. Logic Allocator: Configuration of Cluster " $n$ " Set by Mode of Macrocell " $n$ "

a. Basic cluster with XOR

b. Extended cluster, active high

c. Extended cluster, active low

d. Basic cluster routed away; single-product-term, active high

e. Extended cluster routed away

Figure 3. Logic Allocator Configurations: Synchronous Mode

a. Basic cluster with XOR

b. Extended cluster, active high

c. Extended cluster, active low

d. Basic cluster routed away; single-product-term, active high

e. Extended cluster routed away

Figure 4. Logic Allocator Configurations: Asynchronous Mode

Note that the configuration of the logic allocator has absolutely no impact on the speed of the signal. All configurations have the same delay. This means that designers do not have to decide between optimizing resources or speed; both can be optimized.
If not used in the cluster, the extra product term can act in conjunction with the basic cluster to provide XOR logic for such functions as data comparison, or it can work with the D-,T-type flip-flop to provide for J-K, and S-R register operation. In addition, if the basic cluster is routed to another macrocell, the extra product term is still available for logic. In this case, the first XOR input will be a logic 0 . This circuit has the flexibility to route product terms elsewhere without giving up the use of the macrocell.
Product term clusters do not "wrap" around a PAL block. This means that the macrocells at the ends of the block have fewer product terms available.

## Macrocell

The macrocell consists of a storage element, routing resources, a clock multiplexer, and initialization control. The macrocell has two fundamental modes: synchronous and asynchronous (Figure 5). The mode chosen only affects clocking and initialization in the macrocell.


Figure 5. Macrocell

In either mode, a combinatorial path can be used. For combinatorial logic, the synchronous mode will generally be used, since it provides more product terms in the allocator.

The flip-flop can be configured as a D-type or T-type latch. J-K or S-R registers can be synthesized. The primary flip-flop configurations are shown in Figure 6, although others are possible. Flip-flop functionality is defined in Table 8. Note that a J-K latch is inadvisable as it will cause oscillation if both J and K inputs are HIGH.

a. D-type with XOR

c. Latch with XOR

e. T-type with programmable T polarity

b. D-type with programmable D polarity

d. Latch with programmable $D$ polarity

f. Combinatorial with XOR

g. Combinatorial with programmable polarity

Figure 6. Primary Macrocell Configurations

Table 8. Register/Latch Operation

| Configuration | Input(s) | CLK/LE ${ }^{1}$ | Q+ |
| :---: | :---: | :---: | :---: |
| D-type Register | D $=\mathrm{X}$ | $0,1, \downarrow(\uparrow)$ | Q |
|  | $\mathrm{D}=0$ | $\uparrow(\downarrow)$ | 0 |
|  | $\mathrm{D}=1$ | $\uparrow(\downarrow)$ | 1 |
| T-type Register | $\mathrm{T}=\mathrm{X}$ | $0,1, \downarrow(\uparrow)$ | Q |
|  | $\mathrm{T}=0$ | $\uparrow(\downarrow)$ | Q |
|  | $\mathrm{T}=1$ | $\uparrow(\downarrow)$ | $\overline{\text { Q }}$ |
| D-type Latch | $\mathrm{D}=\mathrm{X}$ | 1 (0) | Q |
|  | $\mathrm{D}=0$ | 0 (1) | 0 |
|  | $\mathrm{D}=1$ | 0 (1) | 1 |

## Note:

1. Polarity of CLK/LE an be programmed

Although the macrocell shows only one input to the register, the XOR gate in the logic allocator allows the D-, T-type register to emulate J-K, and S-R behavior. In this case, the available product terms are divided between J and K (or S and R). When configured as J-K, S-R, or T-type, the extra product term must be used on the XOR gate input for flip-flop emulation. In any register type, the polarity of the inputs can be programmed.
The clock input to the flip-flop can select any of the four PAL block clocks in synchronous mode, with the additional choice of either polarity of an individual product term clock in the asynchronous mode.

The initialization circuit depends on the mode. In synchronous mode (Figure 7), asynchronous reset and preset are provided, each driven by a product term common to the entire PAL block.


Figure 7. Synchronous Mode Initialization Configurations

A reset/preset swapping feature in each macrocell allows for reset and preset to be exchanged, providing flexibility. In asynchronous mode (Figure 8), a single individual product term is provided for initialization. It can be selected to control reset or preset.

a. Reset

b. Preset

17466G-014
17466G-015
Figure 8. Asynchronous Mode Initialization Configurations
Note that the reset/preset swapping selection feature effects power-up reset as well. The initialization functionality of the flip-flops is illustrated in Table 9. The macrocell sends its data to the output switch matrix and the input switch matrix. The output switch matrix can route this data to an output if so desired. The input switch matrix can send the signal back to the central switch matrix as feedback.

Table 9. Asynchronous Reset/Preset Operation

| AR | AP | CLK/LE $^{\mathbf{1}}$ | Q+ |
| :---: | :---: | :---: | :---: |
| 0 | 0 | X | See Table 8 |
| 0 | 1 | x | 1 |
| 1 | 0 | x | 0 |
| 1 | 1 | X | 0 |

## Note:

1. Transparent latch is unaffected by $A R, A P$

## Output Switch Matrix

The output switch matrix allows macrocells to be connected to any of several I/O cells within a PAL block. This provides high flexibility in determining pinout and allows design changes to occur without effecting pinout.

In ispMACH 4A devices with 2:1 Macrocell-I/O cell ratio, each PAL block has twice as many macrocells as I/O cells. The ispMACH 4A output switch matrix allows for half of the macrocells to drive I/O cells within a PAL block, in combinations according to Figure 9. Each I /O cell can choose from eight macrocells; each macrocell has a choice of four I/O cells. The ispMACH 4A devices with 1:1 Macrocell-I/O cell ratio allow each macrocell to drive one of eight I/O cells (Figure 9).


Figure 9. ispMACH 4A Output Switch Matrix

Table 10. Output Switch Matrix Combinations for ispMACH 4A Devices with 2:1 Macrocell-I/O Cell Ratio

| Macrocell | Routable to I/O Cells |
| :---: | :--- |
| M0, M1 | $\mathrm{I} / \mathrm{O} 0, \mathrm{I} / \mathrm{O} 5, \mathrm{I} / \mathrm{O} 6, \mathrm{I} / \mathrm{O} 7$ |
| M2, M3 | $\mathrm{I} / \mathrm{O} 0, \mathrm{I} / \mathrm{O} 1, \mathrm{I} / \mathrm{O}, \mathrm{I} / \mathrm{O} 7$ |
| M4, M5 | $\mathrm{I} / \mathrm{O} 0, \mathrm{I} / \mathrm{O} 1, \mathrm{I} / \mathrm{O} 2, \mathrm{I} / \mathrm{O} 7$ |
| M6, M7 | $\mathrm{I} / \mathrm{O} 0, \mathrm{I} / \mathrm{O} 1, \mathrm{I} / \mathrm{O} 2, \mathrm{I} / \mathrm{O} 3$ |
| M8, M9 | $\mathrm{I} / \mathrm{O} 1, \mathrm{I} / \mathrm{O} 2, \mathrm{I} / \mathrm{O} 3, \mathrm{I} / \mathrm{O} 4$ |
| M10, M11 | $\mathrm{I} / \mathrm{O} 2, \mathrm{I} / \mathrm{O} 3, \mathrm{I} / \mathrm{O} 4, \mathrm{I} / \mathrm{O} 5$ |

Table 10. Output Switch Matrix Combinations for ispMACH 4A Devices with 2:1 Macrocell-I/O Cell Ratio

| Macrocell | Routable to I/O Cells |
| :---: | :---: |
| M12, M13 | I/O3, I/O4, I/O5, I/O6 |
| M14, M15 | I/O4, I/O5, I/O6, I/O7 |
|  |  |
| I/O Cell | Available Macrocells |
| I/O0 | M0, M1, M2, M3, M4, M5, M6, M7 |
| I/O1 | M2, M3, M4, M5, M6, M7, M8, M9 |
| I/O2 | M4, M5, M6, M7, M8, M9, M10, M11 |
| I/O3 | M6, M7, M8, M9, M10, M11, M12, M13 |
| I/O4 | M8, M9, M10, M11, M12, M13, M14, M15 |
| I/O5 | M0, M1, M10, M11, M12, M13, M14, M15 |
| I/O6 | M0, M1, M2, M3, M12, M13, M14, M15 |
| I/O7 | M0, M1, M2, M3, M4, M5, M14, M15 |

Table 11. Output Switch Matrix Combinations for M4A3-256/160 and M4A3-256/192

| $\frac{\text { Macrocell }}{\text { M0 }}$ | Routable to I/O Cells |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M1 | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M2 | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M3 | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M4 | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M5 | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M6 | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M7 | I/O0 | I/O1 | I/O2 | I/O3 | I/O4 | I/O5 | I/O6 | I/O7 |
| M8 | I/O8 | I/O9 | I/O10 | I/O11 | I/O12 | I/O13 | I/O14 | I/O15 |
| M9 | I/08 | I/O9 | I/O10 | I/O11 | I/O12 | I/O13 | I/O14 | I/O15 |
| M10 | I/O8 | I/O9 | I/O10 | I/O11 | I/O12 | I/013 | I/O14 | I/O15 |
| M11 | I/08 | I/O9 | I/O10 | I/O11 | I/O12 | I/013 | I/O14 | I/O15 |
| M12 | I/O8 | I/O9 | I/O10 | I/O11 | I/O12 | I/013 | I/O14 | I/O15 |
| M13 | I/O8 | I/O9 | I/O10 | I/O11 | I/O12 | I/O13 | I/O14 | I/O15 |
| M14 | I/O8 | I/O9 | I/O10 | I/O11 | I/O12 | I/O13 | I/O14 | I/O15 |
| M15 | I/O8 | I/O9 | I/O10 | I/O11 | I/O12 | I/O13 | I/O14 | I/O15 |


| I/O Cell | Available Macrocells |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| I/O0 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |
| I/O1 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |
| I/O2 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |
| I/O3 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |
| I/O4 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |
| I/O5 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |
| I/O6 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |
| I/O7 | M0 | M1 | M2 | M3 | M4 | M5 | M6 | M7 |

Table 11. Output Switch Matrix Combinations for M4A3-256/160 and M4A3-256/192

| Macrocell | Routable to I/O Cells |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| I/O8 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |
| I/O9 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |
| I/O10 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |
| I/O11 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |
| I/O12 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |
| I/O13 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |
| I/O14 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |
| I/O15 | M8 | M9 | M10 | M11 | M12 | M13 | M14 | M15 |

Table 12. Output Switch Matrix Combinations for M4A(3,5)-32/32

| Macrocell | Routable to I/O Cells |
| :---: | :---: |
| M0, M1, M2, M3, M4, M5, M6, M7 | I/O0, I/O1, I/O2, I/O3, I/O4, I/O5, I/O6, I/O7 |
| M8, M9, M10, M11, M12, M13, M14, M15 | $\mathrm{I} / \mathrm{O} 8, \mathrm{I} / \mathrm{O} 9, \mathrm{I} / \mathrm{O} 10, \mathrm{I} / \mathrm{O} 11, \mathrm{I} / \mathrm{O} 12, \mathrm{I} / \mathrm{O} 13, \mathrm{I} / \mathrm{O} 14, \mathrm{I} / \mathrm{O} 15$ |


| I/O Cell | Available Macrocells |
| :---: | :---: |
| $\mathrm{I} / \mathrm{O} 0, \mathrm{I} / \mathrm{O} 1, \mathrm{I} / \mathrm{O} 2, \mathrm{I} / \mathrm{O} 3, \mathrm{I} / \mathrm{O} 4, \mathrm{I} / \mathrm{O} 5, \mathrm{I} / \mathrm{O} 6, \mathrm{I} / \mathrm{O} 7$ | $\mathrm{M} 0, \mathrm{M} 1, \mathrm{M} 2, \mathrm{M} 3, \mathrm{M} 4, \mathrm{M} 5, \mathrm{M} 6, \mathrm{M} 7$ |
| $\mathrm{I} / \mathrm{O} 8, \mathrm{I} / \mathrm{O} 9, \mathrm{I} / \mathrm{O} 10, \mathrm{I} / \mathrm{O} 11, \mathrm{I} / \mathrm{O} 12, \mathrm{I} / \mathrm{O} 13, \mathrm{I} / \mathrm{O} 14, \mathrm{I} / \mathrm{O} 15$ | $\mathrm{M} 8, \mathrm{M} 9, \mathrm{M} 10, \mathrm{M} 11, \mathrm{M} 12, \mathrm{M} 13, \mathrm{M} 14, \mathrm{M} 15$ |

Table 13. Output Switch Matrix Combinations for M4A3-64/64

| Macrocell | Routable to I/O Cells |
| :---: | :---: |
| MO, M1 | I/O0, I/O1, I/O10, I/O11, I/O12, I/O13, I/O14, I/O15 |
| M2, M3 | I/O0, I/O1, I/O2, I/O3, I/O12, I/O13, I/O14, I/O15 |
| M4, M5 | I/O0, I/O1, I/O2, I/O3, I/O4, I/O5, I/O14, I/O15 |
| M6, M7 | I/O0, I/O1, I/O2, I/O3, I/O4, I/O5, I/O6, I/O7 |
| M8, M9 | I/O2, I/O3, I/O4, I/O5, I/O6, I/O7, I/O8, I/O9 |
| M10, M11 | I/O4, I/O5, I/O6, I/O7, I/O8, I/O9, I/O10, I/O11 |
| M12, M13 | I/O6, I/O7, I/O8, I/O9, I/O10, I/O11, I/O12, I/O13 |
| M14, M15 | I/O8, I/O9, I/O10, I/O11, I/O12, I/O13, I/O14, I/O15 |
| I/O Cell | Available Macrocells |
| I/O0, I/O1 | M0, M1, M2, M3, M4, M5, M6, M7 |
| I/O2, I/O3 | M2, M3, M4, M5, M6, M7, M8, M9 |
| I/O4, I/O5 | M4, M5, M6, M7, M8, M9, M10, M11 |
| I/O6, I/O7 | M6, M7, M8, M9, M10, M11, M12, M13 |
| I/O8, I/O9 | M8, M9, M10, M11, M12, M13, M14, M15 |
| I/O10, I/O11 | M0, M1, M10, M11, M12, M13, M14, M15 |
| I/O12, I/O13 | M0, M1, M2, M3, M12, M13, M14, M15 |
| I/O14, I/O15 | M0, M1, M2, M3, M4, M5, M14, M15 |

## I/O Cell

The I/O cell (Figures 10 and 11) simply consists of a programmable output enable, a feedback path, and flip-flop (except ispMACH 4A devices with 1:1 macrocell-I/O cell ratio). An individual output enable product term is provided for each I/O cell. The feedback signal drives the input switch matrix.


17466G-017
17466G-018
Figure 10. I/O Cell for ispMACH 4A Devices with 2:1 Macrocell-I/O Cell Ratio

Figure 11. I/O Cell for ispMACH 4A Devices with 1:1 Macrocell-I/O Cell Ratio

The I/O cell (Figure 10) contains a flip-flop, which provides the capability for storing the input in a D-type register or latch. The clock can be any of the PAL block clocks. Both the direct and registered versions of the input are sent to the input switch matrix. This allows for such functions as "time-domain-multiplexed" data comparison, where the first data value is stored, and then the second data value is put on the I/O pin and compared with the previous stored value.

Note that the flip-flop used in the ispMACH 4A I /O cell is independent of the flip-flops in the macrocells. It powers up to a logic low.

## Zero-Hold-Time Input Register

The ispMACH 4A devices have a zero-hold-time (ZHT) fuse which controls the time delay associated with loading data into all I/O cell registers and latches. When programmed, the ZHT fuse increases the data path setup delays to input storage elements, matching equivalent delays in the clock path. When the fuse is erased, the setup time to the input storage element is minimized. This feature facilitates doing worst-case designs for which data is loaded from sources which have low (or zero) minimum output propagation delays from clock edges.

Input Switch Matrix
The input switch matrix (Figures 12 and 13) optimizes routing of inputs to the central switch matrix. Without the input switch matrix, each input and feedback signal has only one way to enter the central switch matrix. The input switch matrix provides additional ways for these signals to enter the central switch matrix.


## PAL Block Clock Generation

Each ispMACH 4A device has four clock pins that can also be used as inputs. These pins drive a clock generator in each PAL block (Figure 14). The clock generator provides four clock signals that can be used anywhere in the PAL block. These four PAL block clock signals can consist of a large number of combinations of the true and complement edges of the global clock signals. Table 14 lists the possible combinations.


Figure 14. PAL Block Clock Generator ${ }^{1}$

1. M4A(3,5)-32/ 32 and M4A(3,5)-64/ 32 have only two dock pins, GCL K0 and GCLK1. GCL K2 is tied to GCL K0, and GCL K3 is tied to GCL K1.

Table 14. PAL Block Clock Combinations ${ }^{1}$

| Block CLK0 | Block CLK1 | Block CLK2 | Block CLK3 |
| :---: | :---: | :---: | :---: |
| GCLK0 | GCLK1 | X | x |
| GCLK1 | GCLK1 | X | X |
| GCLK0 | GCLK0 | x | x |
| GCLK1 | GCLK0 | X | X |
| X | X | GCLK2 (GCLK0) | GCLK3 (GCLK1) |
| x | x | GCLK3 ( GCLK1) | GCLK3 (GCLK1) |
| x | x | GCLK2 (GCLK0) | GCLK2 (GCLK0) |
| X | x | $\overline{\text { GCLK3 }}$ ( $\overline{\text { GCLK1 }}$ ) | $\overline{\text { GCLK2 }}$ ( $\overline{\text { GCLK0 }}$ ) |

## Note:

1. Values in parentheses are for the $M 4 A(3,5)-32 / 32$ and $M 4 A(3,5)-64 / 32$.

This feature provides high flexibility for partitioning state machines and dual-phase clocks. It also allows latches to be driven with either polarity of latch enable, and in a master-slave configuration.

## ispMACH 4A TIMING MODEL

The primary focus of the ispMACH 4A timing model is to accurately represent the timing in a ispMACH 4A device, and at the same time, be easy to understand. This model accurately describes all combinatorial and registered paths through the device, making a distinction between internal feedback and external feedback. A signal uses internal feedback when it is fed back into the switch matrix or block without having to go through the output buffer. The input register specifications are also reported as internal feedback. When a signal is fed back into the switch matrix after having gone through the output buffer, it is using external feedback.

The parameter, $\mathrm{t}_{\mathrm{BUF}}$, is defined as the time it takes to go from feedback through the output buffer to the I/O pad. If a signal goes to the internal feedback rather than to the I/O pad, the parameter designator is followed by an " i ". By adding $\mathrm{t}_{\mathrm{BUF}}$ to this internal parameter, the external parameter is derived. For example, $\mathrm{t}_{\mathrm{PD}}=\mathrm{t}_{\mathrm{PDi}}+\mathrm{t}_{\mathrm{BUF}}$. A diagram representing the modularized ispMACH 4A timing model is shown in Figure 15. Refer to the application note entitled MA CH 4 Timing and High Speed Design for a more detailed discussion about the timing parameters.


17466G-025
Figure 15. ispMACH 4A Timing Model

## SPEEDLOCKING FOR GUARANTEED FIXED TIMING

The ispMACH 4A architecture allows allocation of up to 20 product terms to an individual macrocell with the assistance of an XOR gate without incurring additional timing delays.

The design of the switch matrix and PAL blocks guarantee a fixed pin-to-pin delay that is independent of the logic required by the design. Other competitive CPLDs incur serious timing delays as product terms expand beyond their typical 4 or 5 product term limits. Speed and SpeedLocking combine to give designs easy access to the performance required in today's designs.

## IEEE 1149.1-COMPLIANT BOUNDARY SCAN TESTABILITY

All ispMACH 4A devices have boundary scan cells and are compliant to the IEEE 1149.1 standard. This allows functional testing of the circuit board on which the device is mounted through a serial scan path that can access all critical logic nodes. Internal registers are linked internally, allowing test data to be shifted in and loaded directly onto test nodes, or test node data to be captured and shifted out for verification. In addition, these devices can be linked into a board-level serial scan path for more complete board-level testing.

## IEEE 1149.1-COMPLIANT IN-SYSTEM PROGRAMMING

Programming devices in-system provides a number of significant benefits including: rapid prototyping, lower inventory levels, higher quality, and the ability to make in-field modifications. All ispMACH 4A devices provide In-System Programming (ISP) capability through their Boundary ScanTest Access Ports. This capability has been implemented in a manner that ensures that the port remains compliant to the IE EE 1149.1 standard. By using IEEE 1149.1 as the communication interface through which ISP is achieved, customers get the benefit of a standard, well-defined interface.
ispMACH 4A devices can be programmed across the commercial temperature and voltage range. The PCbased ispVM ${ }^{\mathrm{TM}}$ software facilitates in-system programming of ispMACH 4A devices. ispVM takes the JEDEC file output produced by the design implementation software, along with information about the JTAG chain, and creates a set of vectors that are used to drive the JTAG chain. ispVM software can use these vectors to drive a JTAG chain via the parallel port of a PC. Alternatively, ispVM software can output files in formats understood by common automated test equipment. This equpment can then be used to program ispMACH 4A devices during the testing of a circuit board.

## PCI COMPLIANT

ispMACH 4A devices in the $-5 /-55 /-6 /-65 /-7 /-10 /-12$ speed grades are compliant with the PCI Local Bus Specification version 2.1, published by the PCI Special Interest Group (SIG). The 5-V devices are fully PCIcompliant. The 3.3-V devices are mostly compliant but do not meet the PCI condition to clamp the inputs as they rise above $\mathrm{V}_{\mathrm{CC}}$ because of their $5-\mathrm{V}$ input tolerant feature.

## SAFE FOR MIXED SUPPLY VOLTAGE SYSTEM DESIGNS

Both the $3.3-\mathrm{V}$ and $5-\mathrm{V} \mathrm{V}_{\mathrm{CC}}$ ispMACH 4 A devices are safe for mixed supply voltage system designs. The $5-\mathrm{V}$ devices will not overdrive $3.3-\mathrm{V}$ devices above the output voltage of 3.3 V , while they accept inputs from other $3.3-\mathrm{V}$ devices. The $3.3-\mathrm{V}$ device will accept inputs up to 5.5 V . Both the $5-\mathrm{V}$ and $3.3-\mathrm{V}$ versions have the same high-speed performance and provide easy-to-use mixed-voltage design capability.

## PULL UP OR BUS-FRIENDLY INPUTS AND I/Os

All ispMACH 4A devices have inputs and I/Os which feature the Bus-Friendly circuitry incorporating two inverters in series which loop back to the input. This double inversion weakly holds the input at its last driven logic state. While it is good design practice to tie unused pins to a known state, the Bus-Friendly input structure pulls pins away from the input threshold voltage where noise can cause high-frequency switching. At power-up, the Bus-Friendly latches are reset to a logic level " 1. " For the circuit diagram, please refer to the document entitled MA CH E ndurance Characteristics on the Lattice Data Book CD-ROM or Lattice web site.

All ispMACH 4A devices have a programmable bit that configures all inputs and I/Os with either pull-up or Bus-Friendly characteristics. If the device is configured in pull-up mode, all inputs and I/O pins are
weakly pulled up. For the circuit diagram, please refer to the document entitled MACH E ndurance Characteristics on the Lattice Data Book CD-ROM or Lattice web site.

## POWER MANAGEMENT

Each individual PAL block in ispMACH 4A devices features a programmable low-power mode, which results in power savings of up to $50 \%$. The signal speed paths in the low-power PAL block will be slower than those in the non-low-power PAL block. This feature allows speed critical paths to run at maximum frequency while the rest of the signal paths operate in the low-power mode.

## PROGRAMMABLE SLEW RATE

Each ispMACH 4A device I/O has an individually programmable output slew rate control bit. Each output can be individually configured for the higher speed transition ( $3 \mathrm{~V} / \mathrm{ns}$ ) or for the lower noise transition (1 $\mathrm{V} / \mathrm{ns}$ ). For high-speed designs with long, unterminated traces, the slow-slew rate will introduce fewer reflections, less noise, and keep ground bounce to a minimum. For designs with short traces or well terminated lines, the fast slew rate can be used to achieve the highest speed. The slew rate is adjusted independent of power.

## POWER-UP RESET/SET

All flip-flops power up to a known state for predictable system initialization. If a macrocell is configured to SET on a signal from the control generator, then that macrocell will be SET during device power-up. If a macrocell is configured to RESET on a signal from the control generator or is not configured for set/reset, then that macrocell will RESET on power-up. To guarantee initialization values, the $\mathrm{V}_{\mathrm{CC}}$ rise must be monotonic, and the clock must be inactive until the reset delay time has elapsed.

## SECURITY BIT

A programmable security bit is provided on the ispMACH 4A devices as a deterrent to unauthorized copying of the array configuration patterns. Once programmed, this bit defeats readback of the programmed pattern by a device programmer, securing proprietary designs from competitors. Programming and verification are also defeated by the security bit. The bit can only be reset by erasing the entire device.

## HOT SOCKETING

ispMACH 4A devices are well-suited for those applications that require hot socketing capability. Hot socketing a device requires that the device, when powered down, can tolerate active signals on the I/Os and inputs without being damaged. Additionally, it requires that the effects of the powered-down MACH devices be minimal on active signals.


Figure 16. PAL Block for ispMACH 4A with 2:1 Macrocell - I/O Cell Ratio


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Figure 17. PAL Block for ispMACH 4A Devices with 1:1 Macrocell-I/O Cell Ratio (except M4A (3,5)-32/32)


Figure 18. PAL Block for M4A (3,5)-32/32
17466H-042

