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# Mobile SDRAM

## MT48LC8M32LF, MT48V8M32LF, MT48H8M32LF - 2 Meg x 32 x 4 banks

For the latest data sheet, refer to Micron's Web site: [www.micron.com/products/dram/mobile](http://www.micron.com/products/dram/mobile)

### Features

- Low voltage power supply
- Partial array self refresh power-saving mode
- Temperature Compensated Self Refresh (TCSR)
- Deep power-down mode
- Programmable output drive strength
- Fully synchronous; all signals registered on positive edge of system clock
- Internal pipelined operation; column address can be changed every clock cycle
- Internal banks for hiding row access/precharge
- Programmable burst lengths: 1, 2, 4, 8, or full page
- Auto precharge, includes concurrent auto precharge, and auto refresh modes
- Self-refresh mode; standard and low power
- 64ms, 4,096-cycle refresh
- LVTTTL-compatible inputs and outputs
- Commercial and industrial temperature ranges
- Supports CAS latency of 1, 2, 3

### Options

- V<sub>DD</sub>/V<sub>DDQ</sub>
  - 3.3V/3.3V LC
  - 2.5V/2.5V V
  - 1.8V/1.8V H
- Configurations
  - 8 Meg x 32 (2 Meg x 32 x 4 banks) 8M32
- Package/Ballout
  - 90-ball VFBGA (8mm x 13mm) (Standard) F5
  - 90-ball VFBGA (8mm x 13mm) (Lead-free) B5
- Timing (Cycle Time)
  - 7.5ns @ CL = 3 (133 MHz) -75
  - 7.5ns @ CL = 2 (104 MHz) -75
  - 8ns @ CL = 3 (125 MHz) -8
  - 8ns @ CL = 2 (104 MHz) -8
  - 10ns @ CL = 3 (100 MHz) -10
  - 10ns @ CL = 2 (83 MHz) -10
- Operating Temperature Range
  - Commercial (0° to +70°C) None
  - Industrial (-40°C to +85°C) IT

**Table 1: Addressing**

	8 Meg x 32
Configuration	2 Meg x 32 x 4 banks
Refresh Count	4K
Row Addressing	4K (A0–A11)
Bank Addressing	4 (BA0, BA1)
Column Addressing	512 (A0–A8)

**Table 2: Key Timing Parameters**

CL = CAS (READ) latency

Speed Grade	Clock Frequency	Access Time		Setup Time	Hold Time
		CL = 2	CL = 3		
-75	133 MHz	–	6ns	2.5ns	1ns
-8	125 MHz	–	7ns	2.5ns	1ns
-10	100 MHz	–	7ns	2.5ns	1ns
-75	133 MHz	7ns	–	2.5ns	1ns
-8	104 MHz	8ns	–	2.5ns	1ns
-10	83 MHz	8ns	-	2.5ns	1ns



## Table of Contents

Features	1
FBGA Part Number System	5
General Description	5
Ball Assignment	7
Ball Descriptions	8
Functional Description	9
Initialization	9
Register Definition	9
Mode Register	9
Burst Length	10
Burst Type	10
CAS Latency	12
Commands	17
Operation	21
Bank/Row Activation	21
READs	22
WRITEs	31
Truth Tables	41
Electrical Specifications	46
Absolute Maximum Ratings	46
Notes	54
Timing Diagrams	56
Package Dimensions	75

## List of Figures

Figure 1:	Functional Block Diagram 8 Meg x 32 SDRAM	6
Figure 2:	90-Ball VFBGA (Top View)	7
Figure 3:	Mode Register Definition	11
Figure 4:	CAS Latency	14
Figure 5:	Low Power Extended Mode Register Table	15
Figure 6:	Activating a Specific Row in a Specific Bank	21
Figure 7:	Example: Meeting tRCD (MIN) When $2 < tRCD (MIN)/tCK < 3$	22
Figure 8:	READ Command	23
Figure 9:	CAS Latency	24
Figure 10:	Consecutive READ Bursts	25
Figure 11:	Random READ Accesses	26
Figure 12:	READ to WRITE	28
Figure 13:	READ to WRITE with Extra Clock Cycle	28
Figure 14:	READ to PRECHARGE	29
Figure 15:	Terminating a READ Burst	30
Figure 16:	WRITE Command	31
Figure 17:	WRITE Burst	32
Figure 18:	WRITE to WRITE	32
Figure 19:	Random WRITE Cycles	33
Figure 20:	WRITE to READ	33
Figure 21:	WRITE to PRECHARGE	34
Figure 22:	Terminating a WRITE Burst	35
Figure 23:	PRECHARGE Command	36
Figure 24:	Power-Down	36
Figure 25:	Clock Suspend During WRITE Burst	38
Figure 26:	Clock Suspend During READ Burst	38
Figure 27:	READ With Auto Precharge Interrupted by a READ	39
Figure 28:	READ With Auto Precharge Interrupted by a WRITE	39
Figure 29:	WRITE With Auto Precharge Interrupted by a READ	40
Figure 30:	WRITE With Auto Precharge Interrupted by a WRITE	40
Figure 31:	Typical Self Refresh Current vs. Temperature – 3.3V Part	51
Figure 32:	Typical Self Refresh Current vs. Temperature – 2.5V Part	52
Figure 33:	Typical Self Refresh Current vs. Temperature – 1.8V Part	52
Figure 34:	Initialize and Load Mode Register <sup>1,2</sup>	56
Figure 35:	Power-Down Mode	57
Figure 36:	Clock Suspend Mode	58
Figure 37:	Auto Refresh Mode	59
Figure 38:	Self Refresh Mode	60
Figure 39:	READ – Without Auto Precharge <sup>1</sup>	61
Figure 40:	Read – With Auto Precharge <sup>1</sup>	62
Figure 41:	Single Read – Without Auto Precharge <sup>1</sup>	63
Figure 42:	Single Read – With Auto Precharge <sup>1</sup>	64
Figure 43:	Alternating Bank Read Accesses <sup>1</sup>	65
Figure 44:	Read – Full-page Burst <sup>1</sup>	66
Figure 45:	Read – DQM Operation <sup>1</sup>	67
Figure 46:	Write – Without Auto Precharge <sup>1</sup>	68
Figure 47:	Write – With Auto Precharge <sup>1</sup>	69
Figure 48:	Single Write – Without Auto Precharge <sup>1</sup>	70
Figure 49:	Single Write – With Auto Precharge <sup>1</sup>	71
Figure 50:	Alternating Bank Write Accesses <sup>1</sup>	72
Figure 51:	Write – Full-page Burst <sup>1</sup>	73
Figure 52:	Write – DQM Operation <sup>1</sup>	74
Figure 53:	90-Ball VFBGA (8mm x 13mm)	75

## List of Tables

Table 1:	Addressing . . . . .	1
Table 2:	Key Timing Parameters . . . . .	1
Table 3:	Cross Reference For VFBGA Device Marking . . . . .	5
Table 4:	Ball Descriptions . . . . .	8
Table 5:	Burst Definition Table . . . . .	12
Table 6:	CAS Latency . . . . .	14
Table 7:	Truth Table – Commands and DQM Operation . . . . .	17
Table 8:	Truth Table – CKE . . . . .	41
Table 9:	Truth Table – Current State Bank <i>n</i> , Command To Bank <i>n</i> . . . . .	42
Table 10:	Truth Table – Current State Bank <i>n</i> , Command To Bank <i>m</i> . . . . .	44
Table 11:	Absolute Maximum Ratings . . . . .	46
Table 12:	DC Electrical Characteristics and Operating Conditions (LC version) . . . . .	46
Table 13:	DC Electrical Characteristics and Operating Conditions (V version) . . . . .	47
Table 14:	DC Electrical Characteristics and Operating Conditions (H version) . . . . .	47
Table 15:	Electrical Characteristics and Recommended AC Operating Conditions . . . . .	48
Table 16:	AC Functional Characteristics . . . . .	49
Table 17:	IDD Specifications and Conditions (LC version) . . . . .	49
Table 18:	IDD Specifications and Conditions (V version) . . . . .	50
Table 19:	IDD Specifications and Conditions (H version) . . . . .	50
Table 20:	IDD7 – Self Refresh Current Options . . . . .	51
Table 21:	Capacitance . . . . .	53

**Table 3: Cross Reference For VFBGA Device Marking**

Part Number	VDD/VDDQ	Architecture	VFBGA	Production Marking
MT48LC8M32LFF5-75	3.3V/3.3V	8 Meg x 32	90-ball, 8 x 13mm	D9FMQ
MT48LC8M32LFF5-8	3.3V/3.3V	8 Meg x 32	90-ball, 8 x 13mm	D9CCH
MT48LC8M32LFF5-10	3.3V/3.3V	8 Meg x 32	90-ball, 8 x 13mm	D9CCK
MT48LC8M32LFB5-75	3.3V/3.3V	8 Meg x 32	90-ball, 8 x 13mm	D9FMX
MT48LC8M32LFB5-8	3.3V/3.3V	8 Meg x 32	90-ball, 8 x 13mm	D9CCW
MT48LC8M32LFB5-10	3.3V/3.3V	8 Meg x 32	90-ball, 8 x 13mm	D9CCZ
MT48V8M32LFF5-75	2.5V/2.5V	8 Meg x 32	90-ball, 8 x 13mm	D9FMS
MT48V8M32LFF5-8	2.5V/2.5V	8 Meg x 32	90-ball, 8 x 13mm	D9CCM
MT48V8M32LFF5-10	2.5V/2.5V	8 Meg x 32	90-ball, 8 x 13mm	D9CCP
MT48V8M32LFB5-75	2.5V/2.5V	8 Meg x 32	90-ball, 8 x 13mm	D9FMZ
MT48V8M32LFB5-8	2.5V/2.5V	8 Meg x 32	90-ball, 8 x 13mm	D9CDC
MT48V8M32LFB5-10	2.5V/2.5V	8 Meg x 32	90-ball, 8 x 13mm	D9CDF
MT48H8M32LFF5-75	1.8V/1.8V	8 Meg x 32	90-ball, 8 x 13mm	D9FMV
MT48H8M32LFF5-8	1.8V/1.8V	8 Meg x 32	90-ball, 8 x 13mm	D9CCR
MT48H8M32LFF5-10	1.8V/1.8V	8 Meg x 32	90-ball, 8 x 13mm	D9CCT
MT48H8M32LFB5-75	1.8V/1.8V	8 Meg x 32	90-ball, 8 x 13mm	D9FNB
MT48H8M32LFB5-8	1.8V/1.8V	8 Meg x 32	90-ball, 8 x 13mm	D9CDJ
MT48H8M32LFB5-10	1.8V/1.8V	8 Meg x 32	90-ball, 8 x 13mm	D9CDL

## FBGA Part Number System

Due to space limitations, FBGA-packaged components have an abbreviated part marking that is different from the part number. For a quick conversion of an FBGA code, see the FBGA part marking decoder on Micron's Web site, [www.micron.com/decoder](http://www.micron.com/decoder).

## General Description

The Micron<sup>®</sup> 256Mb SDRAM is a high-speed CMOS, dynamic random-access memory containing 268,435,456 bits. It is internally configured as a quad-bank DRAM with a synchronous interface (all signals are registered on the positive edge of the clock signal, CLK). Each of the 67,108,864-bit banks is organized as 4,096 rows by 512 columns by 32 bits.

Read and write accesses to the SDRAM are burst oriented; accesses start at a selected location and continue for a programmed number of locations in a programmed sequence. Accesses begin with the registration of an ACTIVE command, which is then followed by a READ or WRITE command. The address bits registered coincident with the ACTIVE command are used to select the bank and row to be accessed (BA0, BA1 select the bank; A0–A11 select the row). The address bits registered coincident with the READ or WRITE command are used to select the starting column location for the burst access.

The SDRAM provides for programmable read or write burst lengths of 1, 2, 4, or 8 locations, or the full page, with a burst terminate option. An auto precharge function may be enabled to provide a self-timed row precharge that is initiated at the end of the burst sequence.

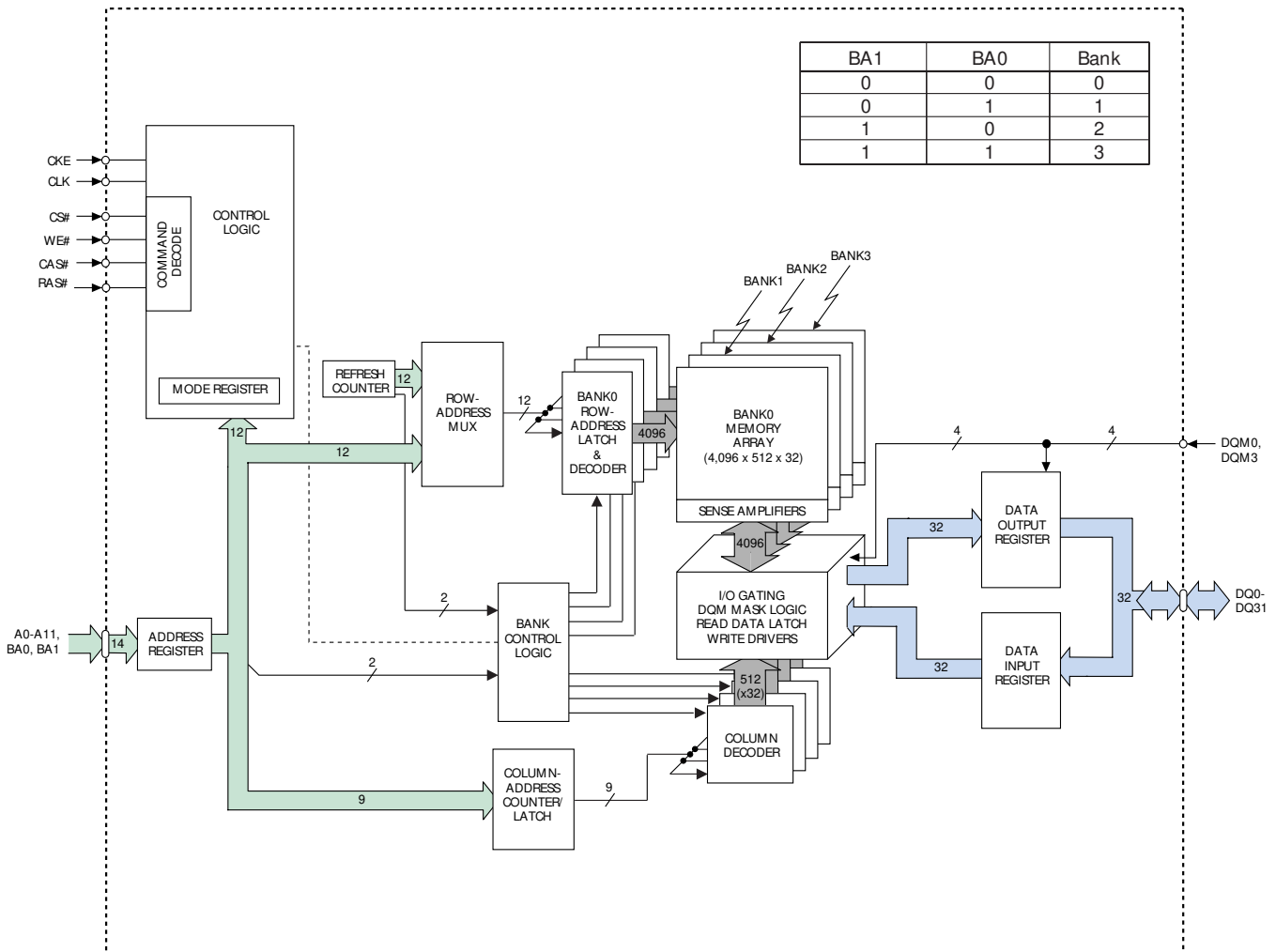
The 256Mb SDRAM uses an internal pipelined architecture to achieve high-speed operation. This architecture is compatible with the  $2n$  rule of prefetch architectures, but it also allows the column address to be changed on every clock cycle to achieve a high-

speed, fully random access. Precharging one bank while accessing one of the other three banks will hide the precharge cycles and provide seamless high-speed, random-access operation.

The 256Mb SDRAM is designed to operate in 3.3V, 2.5V, and 1.8V low-power memory systems. An auto refresh mode is provided, along with a power-saving, deep power-down mode. All inputs and outputs are LVTTTL-compatible.

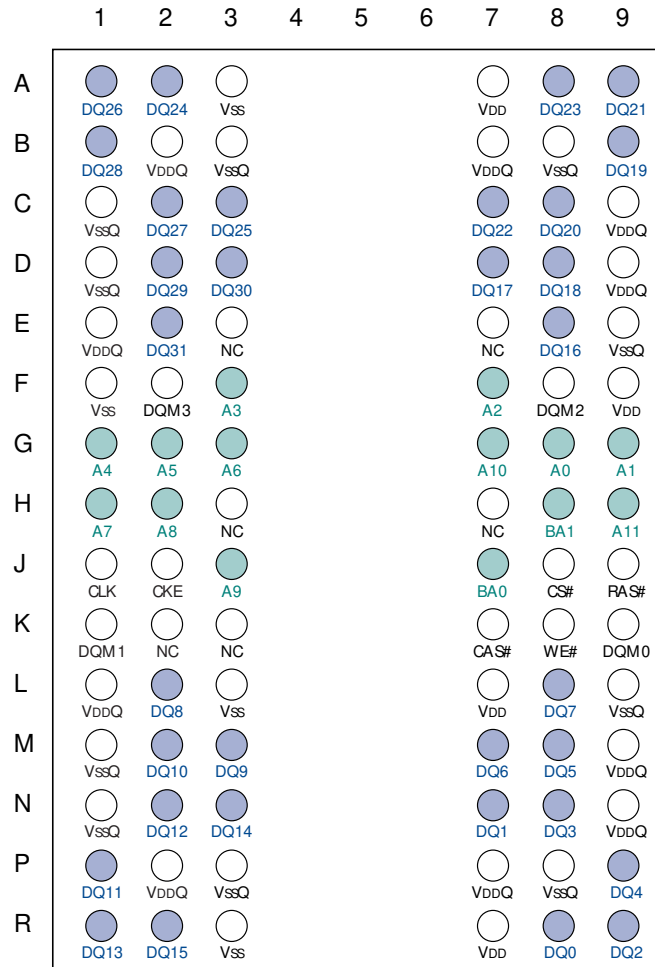
SDRAMs offer substantial advances in DRAM operating performance, including the ability to synchronously burst data at a high data rate with automatic column-address generation, the ability to interleave between internal banks in order to hide precharge time, and the capability to randomly change column addresses on each clock cycle during a burst access.

**Figure 1: Functional Block Diagram 8 Meg x 32 SDRAM**



## Ball Assignment

Figure 2: 90-Ball VFBGA (Top View)





**Ball Descriptions**
**Table 4: Ball Descriptions**

90-Ball VFBGA	Symbol	Type	Description
J1	CLK	Input	Clock: CLK is driven by the system clock. All SDRAM input signals are sampled on the positive edge of CLK. CLK also increments the internal burst counter and controls the output registers.
J2	CKE	Input	Clock Enable: CKE activates (HIGH) and deactivates (LOW) the CLK signal. Deactivating the clock provides PRECHARGE POWER-DOWN and SELF REFRESH operation (all banks idle), ACTIVE POWER-DOWN (row active in any bank), DEEP POWER DOWN (all banks idle), or CLOCK SUSPEND operation (burst/access in progress). CKE is synchronous except after the device enters power-down and self refresh modes, where CKE becomes asynchronous until after exiting the same mode. The input buffers, including CLK, are disabled during power-down and self refresh modes, providing low standby power.
J8	CS#	Input	Chip Select: CS# enables (registered LOW) and disables (registered HIGH) the command decoder. All commands are masked when CS# is registered HIGH. CS# provides for external bank selection on systems with multiple banks. CS# is considered part of the command code.
J9, K7, K8	RAS#, CAS#, WE#	Input	Command Inputs: RAS#, CAS#, and WE# (along with CS#) define the command being entered.
K9, K1, F8, F2	DQM0-3	Input	Input/Output Mask: DQM is sampled HIGH and is an input mask signal for write accesses and an output enable signal for read accesses. Input data is masked during a WRITE cycle. The output buffers are placed in a High-Z state (two-clock latency) when during a READ cycle. DQM0 corresponds to DQ0-DQ7, DQM1 corresponds to DQ8-DQ15, DQM2 corresponds to DQ16-DQ23, and DQM3 corresponds to DQ24-DQ31. DQM0-3 are considered same state when referenced as DQM.
J7, H8	BA0, BA1	Input	Bank Address Input(s): BA0 and BA1 define to which bank the ACTIVE, READ, WRITE, or PRECHARGE command is being applied. These balls also provide the op-code during a LOAD MODE REGISTER command
G8, G9, F7, F3, G1, G2, G3, H1, H2, J3, G7, H9	A0-A11	Input	Address Inputs: A0-A11 are sampled during the ACTIVE command (row-address A0-A11) and READ/WRITE command (column-address A0-A8; with A10 defining auto precharge) to select one location out of the memory array in the respective bank. A10 is sampled during a PRECHARGE command to determine if all banks are to be precharged (A10 HIGH) or bank selected by BA0, BA1. The address inputs also provide the op-code during a LOAD MODE REGISTER command.
F8, N7, R9, N8, P9, M8, M7, L8, L2, M3, M2, P1, N2, R1, N3, F2, E8, D7, D8, B9, C8, A9, C7, A8, A2, C3, A1, C2, B1, D2, D3, E2	DQ0-DQ31	I/O	Data Input/Output: Data bus.
E3, E7, H3, H7, K2, K3	NC	-	Internally Not Connected: These could be left unconnected, but it is recommended they be connected to Vss. H3 is a no connect for this part, but may be used as A12 in future designs.
B2, B7, C9, D9, E1, L1, M9, N9, P2, P7	VDDQ	Supply	DQ Power: Provide isolated power to DQs for improved noise immunity.
B3, B8, C1, D1, E9, L9, M1, N1, P3, P8	VSSQ	Supply	DQ Ground: Provide isolated ground to DQs for improved noise immunity.
A7, F9, L7, R7	VDD	Supply	Core Power Supply.
A3, F1, L3, R3	VSS	Supply	Ground.

## Functional Description

In general, the 256Mb SDRAMs (2 Meg x 32 x 4 banks) are quad-bank DRAMs that operate at 3.3V, 2.5V, and 1.8V and include a synchronous interface (all signals are registered on the positive edge of the clock signal, CLK). Each of the 67,108,864-bit banks is organized as 4,096 rows by 512 columns by 32 bits.

Read and write accesses to the SDRAM are burst oriented; accesses start at a selected location and continue for a programmed number of locations in a programmed sequence. Accesses begin with the registration of an ACTIVE command, which is then followed by a READ or WRITE command. The address bits registered coincident with the ACTIVE command are used to select the bank and row to be accessed (BA0 and BA1 select the bank, A0–A11 select the row). The address bits (A0–A8) registered coincident with the READ or WRITE command are used to select the starting column location for the burst access.

Prior to normal operation, the SDRAM must be initialized. The following sections provide detailed information covering device initialization, register definition, command descriptions, and device operation.

## Initialization

SDRAMs must be powered up and initialized in a predefined manner. Operational procedures other than those specified may result in undefined operation. Once the power is applied to VDD and VDDQ (simultaneously) and the clock is stable (stable clock is defined as a signal cycling within timing constraints specified for the clock ball), the SDRAM requires a 100µs delay prior to issuing any command other than a COMMAND INHIBIT or NOP. Starting at some point during this 100µs period and continuing at least through the end of this period, COMMAND INHIBIT or NOP commands should be applied.

Once the 100µs delay has been satisfied with at least one COMMAND INHIBIT or NOP command having been applied, a PRECHARGE command should be applied. All banks must then be precharged, thereby placing the device in the all banks idle state.

Once in the idle state, two AUTO REFRESH cycles must be performed. After the AUTO REFRESH cycles are complete, the SDRAM is ready for mode register programming. Because the mode register will power up in an unknown state, it should be loaded prior to applying any operational command.

## Register Definition

### Mode Register

In order to achieve low power consumption, there are two mode registers in the component: mode register and extended mode register. Extended mode register is illustrated in Figure 5. The mode register is used to define the specific mode of operation of the SDRAM. This definition includes the selection of a burst length, a burst type, a CAS latency, an operating mode and a write burst mode, as shown in Figure 3. The mode register is programmed via the LOAD MODE REGISTER command and will retain the stored information until it is programmed again or the device loses power.

Mode register bits M0–M2 specify the burst length, M3 specifies the type of burst (sequential or interleaved), M4–M6 specify the CAS latency, M7 and M8 specify the operating mode, M9 specifies the write burst mode, and M10 and M11 should be set to zero. M12 and M13 should be set to zero to prevent the extended mode register from being programmed.

The mode register must be loaded when all banks are idle, and the controller must wait the specified time before initiating the subsequent operation. Violating either of these requirements will result in unspecified operation.

## Burst Length

Read and write accesses to the SDRAM are burst oriented, with the burst length being programmable, as shown in Figure 3. The burst length determines the maximum number of column locations that can be accessed for a given READ or WRITE command. Burst lengths of 1, 2, 4, or 8 locations are available for both the sequential and the interleaved burst types, and a full-page burst is available for the sequential type. The full-page burst is used in conjunction with the BURST TERMINATE command to generate arbitrary burst lengths.

Reserved states should not be used, as unknown operation or incompatibility with future versions may result.

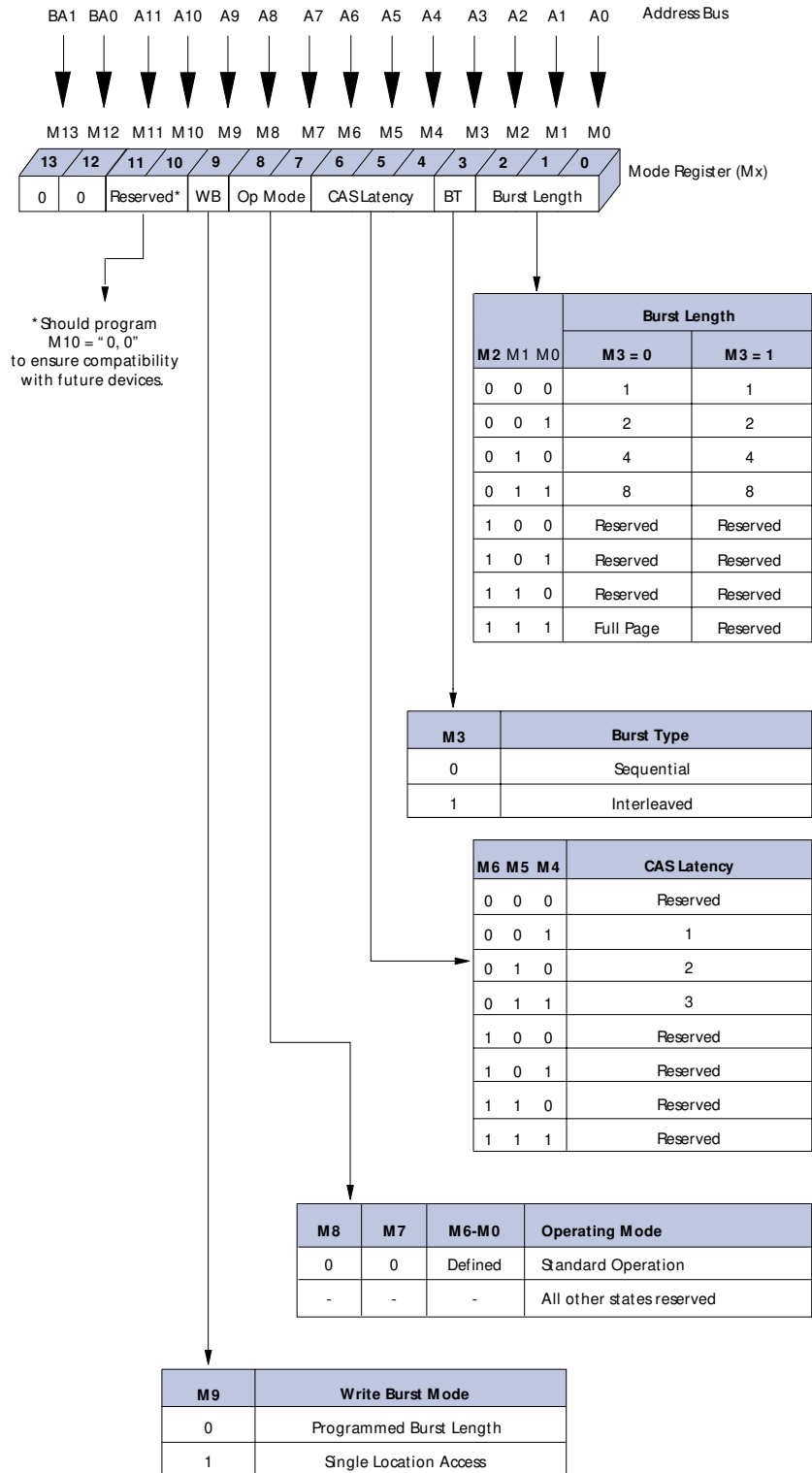
When a READ or WRITE command is issued, a block of columns equal to the burst length is effectively selected. All accesses for that burst take place within this block, meaning that the burst will wrap within the block if a boundary is reached. The block is uniquely selected by A1–A8 when BL = 2, A2–A8 when BL = 4, and A3–A8 when BL = 8. The remaining (least significant) address bit(s) is (are) used to select the starting location within the block. Full-page bursts wrap within the page if the boundary is reached.

## Burst Type

Accesses within a given burst may be programmed to be either sequential or interleaved; this is referred to as the burst type and is selected via bit M3.

The ordering of accesses within a burst is determined by the burst length, the burst type, and the starting column address, as shown in Table 4.

Figure 3: Mode Register Definition



**Table 5: Burst Definition Table**

Burst Length	Starting Column Address			Order of Accesses Within a Burst	
				Type = Sequential	Type = Interleaved
2			A0		
			0	0-1	0-1
			1	1-0	1-0
4		A1	A0		
		0	0	0-1-2-3	0-1-2-3
		0	1	1-2-3-0	1-0-3-2
		1	0	2-3-0-1	2-3-0-1
		1	1	3-0-1-2	3-2-1-0
8	A2	A1	A0		
	0	0	0	0-1-2-3-4-5-6-7	0-1-2-3-4-5-6-7
	0	0	1	1-2-3-4-5-6-7-0	1-0-3-2-5-4-7-6
	0	1	0	2-3-4-5-6-7-0-1	2-3-0-1-6-7-4-5
	0	1	1	3-4-5-6-7-0-1-2	3-2-1-0-7-6-5-4
	1	0	0	4-5-6-7-0-1-2-3	4-5-6-7-0-1-2-3
	1	0	1	5-6-7-0-1-2-3-4	5-4-7-6-1-0-3-2
	1	1	0	6-7-0-1-2-3-4-5	6-7-4-5-2-3-0-1
1	1	1	7-0-1-2-3-4-5-6	7-6-5-4-3-2-1-0	
Full Page (y)	n = A0-A11/9/8 (location 0-y)			Cn, Cn + 1, Cn + 2 Cn + 3, Cn + 4... ...Cn - 1, Cn...	Not Supported

- Notes:
1. For full-page accesses:  $y = 512$ .
  2. For BL = 2, A1–A8 select the block-of-two burst; A0 selects the starting column within the block.
  3. For BL = 4, A2–A8 select the block-of-four burst; A0-A1 select the starting column within the block.
  4. For BL = 8, A3–A8 select the block-of-eight burst; A0-A2 select the starting column within the block.
  5. For a full-page burst, the full row is selected and A0–A8 select the starting column.
  6. Whenever a boundary of the block is reached within a given sequence above, the following access wraps within the block.
  7. For BL = 1, A0–A8 select the unique column to be accessed, and mode register bit M3 is ignored.

## CAS Latency

The CAS latency is the delay, in clock cycles, between the registration of a READ command and the availability of the first piece of output data. The latency can be set to one, two, or three clocks.

If a READ command is registered at clock edge  $n$ , and the latency is  $m$  clocks, the data will be available by clock edge  $n + m$ . The DQs will start driving as a result of the clock edge one cycle earlier ( $n + m - 1$ ), and provided that the relevant access times are met, the data will be valid by clock edge  $n + m$ . For example, assuming that the clock cycle time is such that all relevant access times are met, if a read command is registered at T0

and the latency is programmed to two clocks, the DQs will start driving after T1 and the data will be valid by T2, as shown in Figure 4. Table 5 indicates the operating frequencies at which each CAS latency setting can be used.

Reserved states should not be used as unknown operation or incompatibility with future versions may result.

### **Operating Mode**

The normal operating mode is selected by setting M7 and M8 to zero; the other combinations of values for M7 and M8 are reserved for future use and/or test modes. The programmed burst length applies to both read and write bursts.

Test modes and reserved states should not be used because unknown operation or incompatibility with future versions may result.

### **Write Burst Mode**

When M9 = 0, the burst length programmed via M0-M2 applies to both READ and WRITE bursts; when M9 = 1, the programmed burst length applies to READ bursts, but write accesses are single-location (nonburst) accesses.

### **Low-Power Extended Mode Register Definition**

The low-power extended mode register controls the functions beyond those controlled by the mode register. These additional functions are special features of the mobile device. They include temperature compensated self refresh (TCSR) control, partial array self refresh (PASR), and output drive strength. Not programming the extended mode register upon initialization will result in default settings for the low-power features. The extended mode will default with the temperature sensor enabled, full drive strength, and full array refresh.

The low-power extended mode register is programmed via the MODE REGISTER SET command (BA1 = 1, BA0 = 0) and retains the stored information until it is programmed again or the device loses power.

Figure 4: CAS Latency

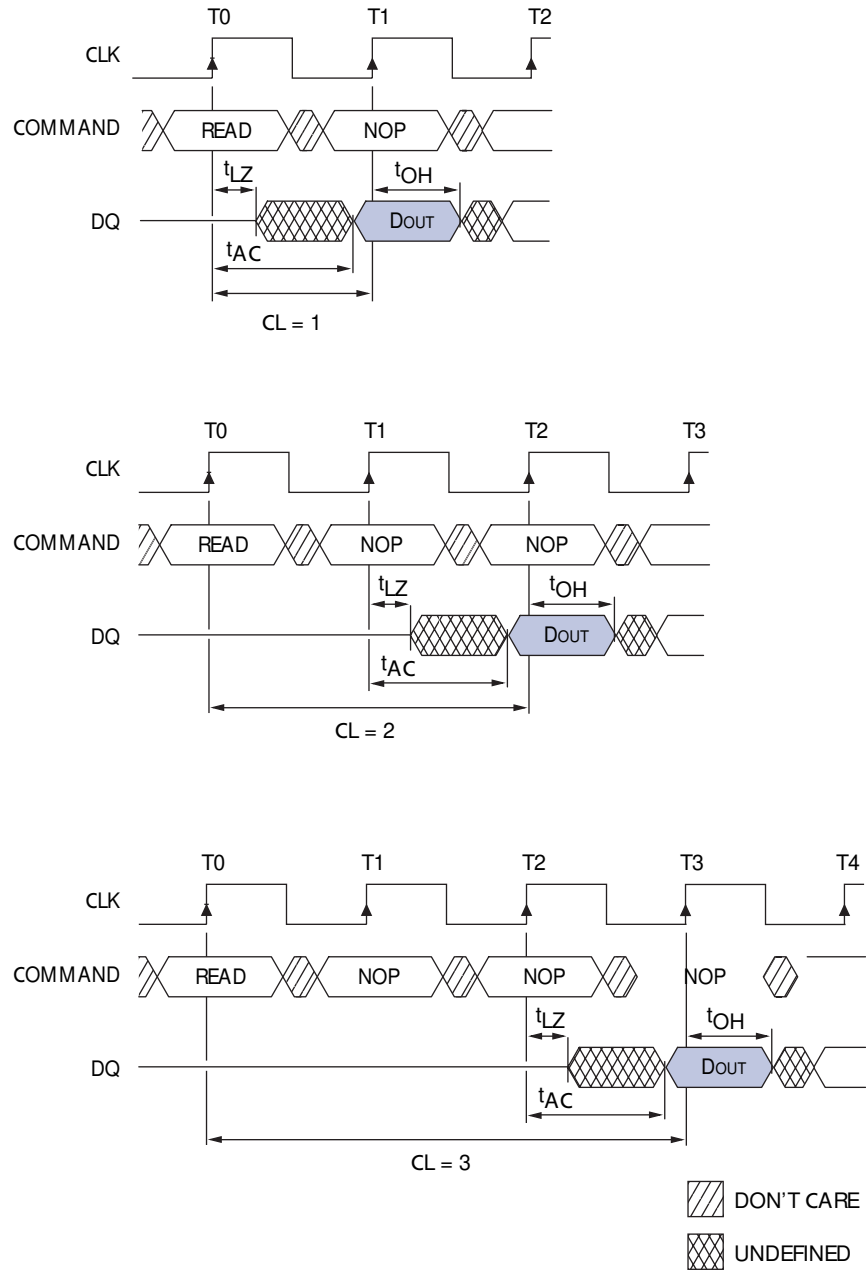
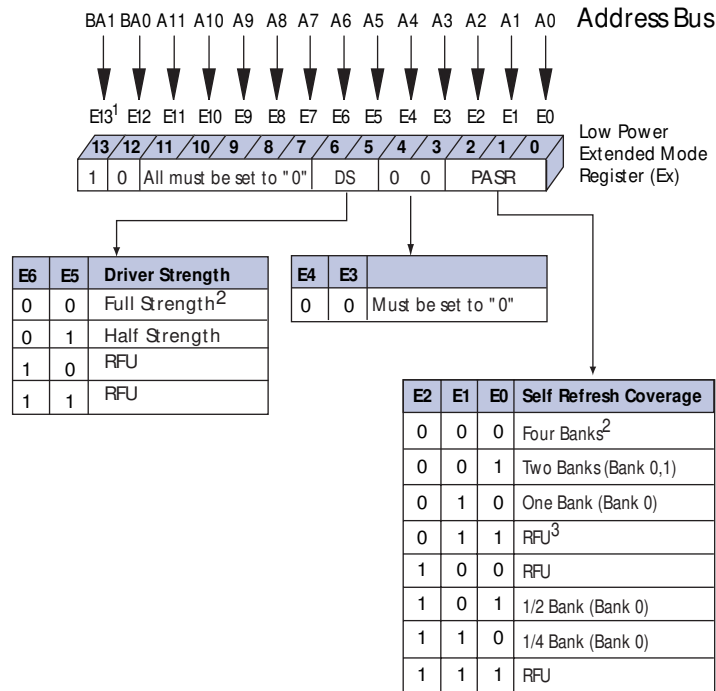


Table 6: CAS Latency

Speed	Allowable Operating Frequency (MHz)		
	CL = 1	CL = 2	CL = 3
-75	–	≤104	≤133
- 8	≤50	≤104	≤125
-10	≤50	≤83.3	≤100

Figure 5: Low Power Extended Mode Register Table



- Notes:
1. E13 and E12 (BA1 and BA0) must be "1, 0" to select the extended mode register (vs. the base mode register).
  2. Default EMR values are full array for PASR, full drive strength.
  3. RFU: reserved for future use.
  4. E4 and E3 are "Don't Care."

The low-power extended mode register must be programmed with E7 through E11 set to "0". It must be loaded when all banks are idle and no bursts are in progress, and the controller must wait the specified time before initiating any subsequent operation. Violating either of these requirements results in unspecified operation. Once the values are entered, the extended mode register settings will be retained even after exiting deep power-down mode.

### Temperature Compensated Self Refresh

Temperature compensated self refresh (TCSR) allows the controller to program the refresh interval during self refresh mode, according to the case temperature of the Mobile device. This allows great power savings during self refresh during most operating temperature ranges. Only during extreme temperatures would the controller have to select the maximum TCSR level. This would guarantee data during self refresh.

Every cell in the DRAM requires refreshing due to the capacitor losing its charge over time. The refresh rate is dependent on temperature. At higher temperatures a capacitor loses charge quicker than at lower temperatures, requiring the cells to be refreshed more often. Historically, during self refresh, the refresh rate has been set to accommodate the worst case, or highest temperature range expected.



Thus, during ambient temperatures, the power consumed during refresh was unnecessarily high because the refresh rate was set to accommodate the higher temperatures. This SDRAM has an on-chip temperature sensor that automatically adjusts refresh rate according to die temperature. The default setting for the TCSR is with the temperature sensor enabled.

### **Partial Array Self Refresh**

For further power savings during self refresh, the partial array self refresh (PASR) feature allows the controller to select the amount of memory that will be refreshed during self refresh. The refresh options are all banks (banks 0, 1, 2, and 3), two banks (banks 0 and 1), and one bank (bank 0). Also included in the refresh options are the half-bank and quarter-bank partial array self refresh (bank 0). WRITE and READ commands occur to any bank selected during standard operation, but only the selected banks in PASR will be refreshed during self refresh. It is important to note that data in banks 2 and 3 will be lost when the two bank option is used. Data will be lost in banks 1, 2, and 3 when the one bank option is used.

### **Driver Strength**

Bits E5 and E6 of the extended mode register can be used to select the driver strength of the DQ outputs. This value should be set according to the application's requirements. Full drive strength was carried over from standard SDRAM and is suitable to drive higher load systems.

## Commands

Table 7 provides a quick reference of available commands. This is followed by a written description of each command. Three additional Truth Tables appear following "Operation" on page 21; these tables provide current state/next state information.

**Table 7: Truth Table – Commands and DQM Operation**

CKE is HIGH for all commands shown except SELF REFRESH and DEEP POWER DOWN

Name (Function)	CS#	RAS#	CAS#	WE#	DQM	ADDR	DQs	Notes
COMMAND INHIBIT (NOP)	H	X	X	X	X	X	X	
NO OPERATION (NOP)	L	H	H	H	X	X	X	
ACTIVE (Select bank and activate row)	L	L	H	H	X	Bank/Row	X	1
READ (Select bank and column, and start READ burst)	L	H	L	H	L/H <sup>8</sup>	Bank/Col	X	2
WRITE (Select bank and column, and start WRITE burst)	L	H	L	L	L/H <sup>8</sup>	Bank/Col	Valid	2
BURST TERMINATE or DEEP POWER DOWN (Enter deep power-down mode)	L	H	H	L	X	X	X	3, 4
PRECHARGE (Deactivate row in bank or banks)	L	L	H	L	X	Code	X	5
AUTO REFRESH or SELF REFRESH (Enter self refresh mode)	L	L	L	H	X	X	X	6, 7
LOAD MODE REGISTER	L	L	L	L	X	Op-Code	X	8
Write Enable/Output Enable	X	X	X	X	L	X	Active	9
Write Inhibit/Output High-Z	X	X	X	X	H	X	High-Z	9

- Notes:
1. A0–A11 provide row address, and BA0, BA1 determine which bank is made active.
  2. A0–A8 provide column address; A10 HIGH enables the auto precharge feature (non persistent), while A10 LOW disables the auto precharge feature; BA0, BA1 determine which bank is being read from or written to.
  3. This command is BURST TERMINATE when CKE is HIGH and DEEP POWER DOWN when CKE is LOW.
  4. The purpose of the BURST TERMINATE command is to stop a data burst, thus the command could coincide with data on the bus. However the DQs column reads a don't care state to illustrate that the BURST TERMINATE command can occur when there is no data present.
  5. A10 LOW: BA0, BA1 determine the bank being precharged. A10 HIGH: All banks precharged and BA0, BA1 are "Don't Care."
  6. This command is AUTO REFRESH if CKE is HIGH, SELF REFRESH if CKE is LOW.
  7. Internal refresh counter controls row addressing; all inputs and I/Os are "Don't Care" except for CKE.
  8. A0–A11 define the op-code written to the mode and extended mode register.
  9. Activates or deactivates the DQs during WRITES (zero-clock delay) and READS (two-clock delay). DQM0 controls DQ0–DQ7; DQM1 controls DQ8–DQ15; DQM2 controls DQ16–DQ23; and DQM3 control DQ24–DQ31.

### COMMAND INHIBIT

The COMMAND INHIBIT function prevents new commands from being executed by the SDRAM, regardless of whether the CLK signal is enabled. The SDRAM is effectively deselected and the DQ balls tri-state. Operations already in progress are not affected.

### NO OPERATION (NOP)

The NO OPERATION (NOP) command is used to perform a NOP to an SDRAM which is selected (CS# is LOW). This prevents unwanted commands from being registered during idle or wait states. Operations already in progress are not affected.

## LOAD MODE REGISTER

The mode register is loaded via inputs A0, BA0, and BA1. (See "Mode Register" on page 9.) The LOAD MODE REGISTER and LOAD EXTENDED MODE REGISTER commands can only be issued when all banks are idle, and a subsequent executable command cannot be issued until  $t^{\text{MRD}}$  is met.

The values of the load mode register and extended mode register will be retained even when exiting deep power-down mode.

## ACTIVE

The ACTIVE command is used to open (or activate) a row in a particular bank for a subsequent access. The value on the BA0, BA1 inputs selects the bank, and the address provided on inputs A0–A11 selects the row. This row remains active (or open) for accesses until a PRECHARGE command is issued to that bank. A PRECHARGE command must be issued before opening a different row in the same bank.

## READ

The READ command is used to initiate a burst read access to an active row. The value on the BA0, BA1 inputs selects the bank, and the address provided on inputs A0–A8 selects the starting column location. The value on input A10 determines whether or not auto precharge is used. If auto precharge is selected, the row being accessed will be precharged at the end of the read burst; if auto precharge is not selected, the row will remain open for subsequent accesses. Read data appears on the DQs subject to the logic level on the DQM inputs two clocks earlier. If a given DQM signal was registered HIGH, the corresponding DQs will be High-Z two clocks later; if the DQM signal was registered LOW, the DQs will provide valid data.

## WRITE

The WRITE command is used to initiate a burst write access to an active row. The value on the BA0, BA1 inputs selects the bank, and the address provided on inputs A0–A8 selects the starting column location. The value on input A10 determines whether or not auto precharge is used. If auto precharge is selected, the row being accessed will be precharged at the end of the write burst; if auto precharge is not selected, the row will remain open for subsequent accesses. Input data appearing on the DQs is written to the memory array subject to the DQM input logic level appearing coincident with the data. If a given DQM signal is registered LOW, the corresponding data will be written to memory; if the DQM signal is registered HIGH, the corresponding data inputs will be ignored, and a write will not be executed to that byte/column location.

## PRECHARGE

The PRECHARGE command is used to deactivate the open row in a particular bank or the open row in all banks. The bank(s) will be available for a subsequent row access a specified time ( $t^{\text{RP}}$ ) after the precharge command is issued. Input A10 determines whether one or all banks are to be precharged, and in the case where only one bank is to be precharged, inputs BA0, BA1 select the bank. Otherwise BA0, BA1 are treated as "Don't Care." Once a bank has been precharged, it is in the idle state and must be activated prior to any READ or WRITE commands being issued to that bank.

## Auto Precharge

Auto precharge is a feature which performs the same individual-bank precharge function described above, without requiring an explicit command. This is accomplished by using A10 to enable auto precharge in conjunction with a specific READ or WRITE command. A precharge of the bank/row that is addressed with the READ or WRITE com-

mand is automatically performed upon completion of the READ or WRITE burst, except in the full-page burst mode, where auto precharge does not apply. Auto precharge is non persistent in that it is either enabled or disabled for each individual READ or WRITE command.

Auto precharge ensures that the precharge is initiated at the earliest valid stage within a burst. The user must not issue another command to the same bank until the precharge time ( $t_{RP}$ ) is completed. This is determined as if an explicit PRECHARGE command was issued at the earliest possible time, as described for each burst type in "Operation" on page 21.

## **BURST TERMINATE**

The BURST TERMINATE command is used to truncate either fixed-length or full-page bursts. The most recently registered READ or WRITE command prior to the BURST TERMINATE command will be truncated, as shown in "Operation" on page 21.

## **AUTO REFRESH**

AUTO REFRESH is used during normal operation of the SDRAM and is analogous to CAS#-BEFORE-RAS# (CBR) refresh in conventional DRAMs. This command is non persistent, so it must be issued each time a refresh is required. All active banks must be PRECHARGED prior to issuing an AUTO REFRESH command. The AUTO REFRESH command should not be issued until the minimum  $t_{RP}$  has been met after the PRECHARGE command, as shown in "Operation" on page 21.

The addressing is generated by the internal refresh controller. This makes the address bits "Don't Care" during an AUTO REFRESH command. The 256Mb SDRAM requires 4,096 AUTO REFRESH cycles every 64ms ( $t_{REF}$ ). Providing a distributed AUTO REFRESH command every 15.625 $\mu$ s will meet the refresh requirement and ensure that each row is refreshed. Alternatively, 4,096 AUTO REFRESH commands can be issued in a burst at the minimum cycle rate ( $t_{RFC}$ ), once every 64ms.

## **SELF REFRESH**

The SELF REFRESH command can be used to retain data in the SDRAM, even if the rest of the system is powered down. When in the self refresh mode, the SDRAM retains data without external clocking. The SELF REFRESH command is initiated like an AUTO REFRESH command, except CKE is disabled (LOW). Once the SELF REFRESH command is registered, all the inputs to the SDRAM become "Don't Care" with the exception of CKE, which must remain LOW.

Once self refresh mode is engaged, the SDRAM provides its own internal clocking, causing it to perform its own auto refresh cycles. The SDRAM must remain in self refresh mode for a minimum period equal to  $t_{RAS}$  and may remain in self refresh mode for an indefinite period beyond that.

The procedure for exiting self refresh requires a sequence of commands. First, CLK must be stable (stable clock is defined as a signal cycling within timing constraints specified for the clock ball) prior to CKE going back HIGH. Once CKE is HIGH, the SDRAM must have NOP commands issued (a minimum of two clocks) for  $t_{XSR}$  because time is required for the completion of any internal refresh in progress.

Upon exiting the self refresh mode, AUTO REFRESH commands must be issued every 15.625 $\mu$ s or less as both SELF REFRESH and AUTO REFRESH utilize the row refresh counter.

**DEEP POWER-DOWN**

DEEP POWER-DOWN is an operating mode to achieve maximum power reduction by eliminating the power of the whole memory array of the devices. Array data will not be retained once the device enters deep power-down mode. The settings in the mode and extended mode register will be retained during deep power-down.

This mode is entered by having all banks idle then CS# and WE# held LOW with RAS# and CAS# held HIGH at the rising edge of the clock, while CKE is LOW. This mode is exited by asserting CKE HIGH.

## Operation

### Bank/Row Activation

Before any READ or WRITE commands can be issued to a bank within the SDRAM, a row in that bank must be “opened.” This is accomplished via the ACTIVE command, which selects both the bank and the row to be activated (see Figure 5 on page 15).

After opening a row (issuing an ACTIVE command), a READ or WRITE command may be issued to that row, subject to the  $t_{RCD}$  specification.  $t_{RCD}(\text{MIN})$  should be divided by the clock period and rounded up to the next whole number to determine the earliest clock edge after the ACTIVE command on which a READ or WRITE command can be entered. For example, a  $t_{RCD}$  specification of 20ns with a 125 MHz clock (8ns period) results in 2.5 clocks, rounded to 3. This is reflected in Figure 6 on page 21, which covers any case where  $2 < t_{RCD}(\text{MIN})/t_{CK} \leq 3$ . (The same procedure is used to convert other specification limits from time units to clock cycles.)

A subsequent ACTIVE command to a different row in the same bank can only be issued after the previous active row has been “closed” (precharged). The minimum time interval between successive ACTIVE commands to the same bank is defined by  $t_{RC}$ .

A subsequent ACTIVE command to another bank can be issued while the first bank is being accessed, which results in a reduction of total row-access overhead. The minimum time interval between successive ACTIVE commands to different banks is defined by  $t_{RRD}$ .

**Figure 6: Activating a Specific Row in a Specific Bank**

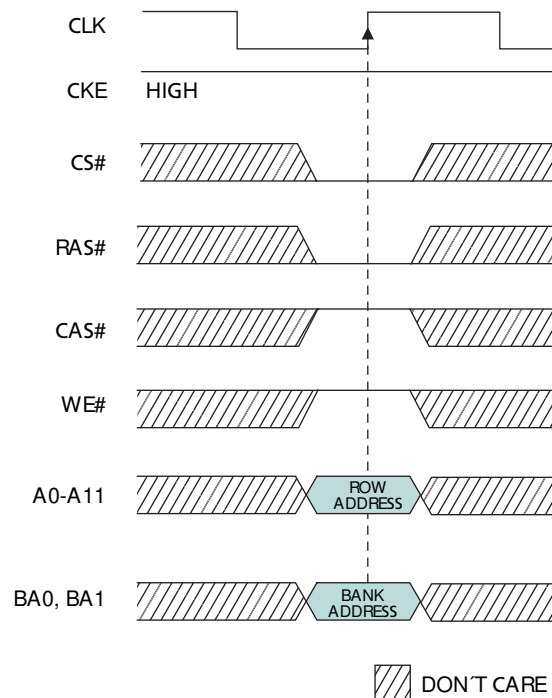
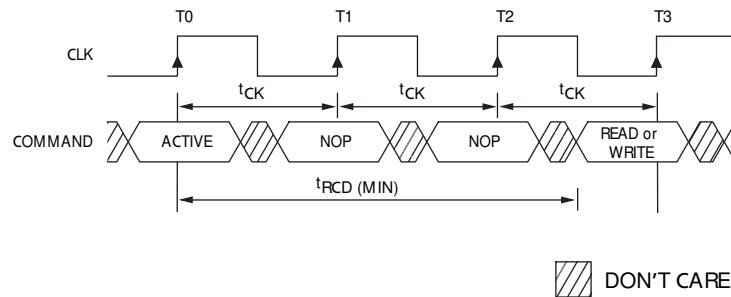


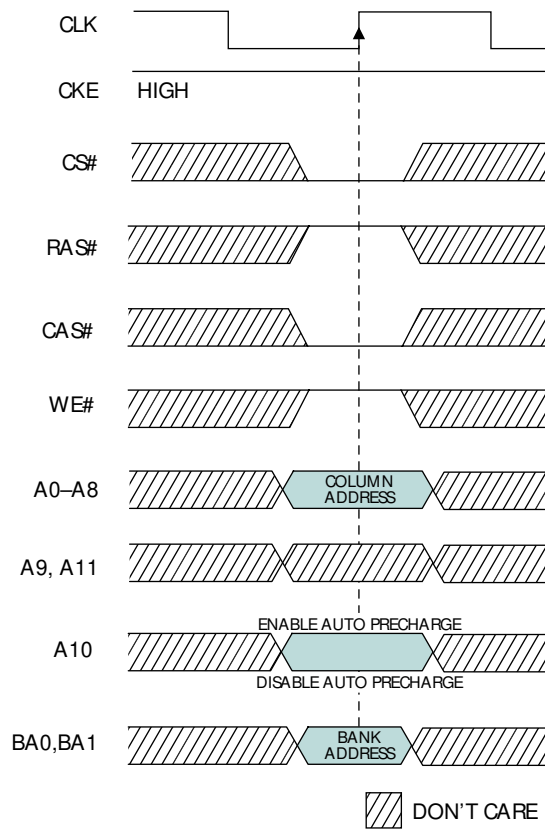
Figure 7: Example: Meeting  $t_{RCD} (MIN)$  When  $2 < t_{RCD} (MIN)/t_{CK} \leq 3$



## READs

READ bursts are initiated with a READ command, as shown in Figure 8. The starting column and bank addresses are provided with the READ command, and auto precharge is either enabled or disabled for that burst access. If auto precharge is enabled, the row being accessed is precharged at the completion of the burst. For the generic READ commands used in the following illustrations, auto precharge is disabled. During READ bursts, the valid data-out element from the starting column address will be available following the CAS latency after the READ command. Each subsequent data-out element will be valid by the next positive clock edge. Figure 9 shows general timing for each possible CAS latency setting.

Figure 8: READ Command

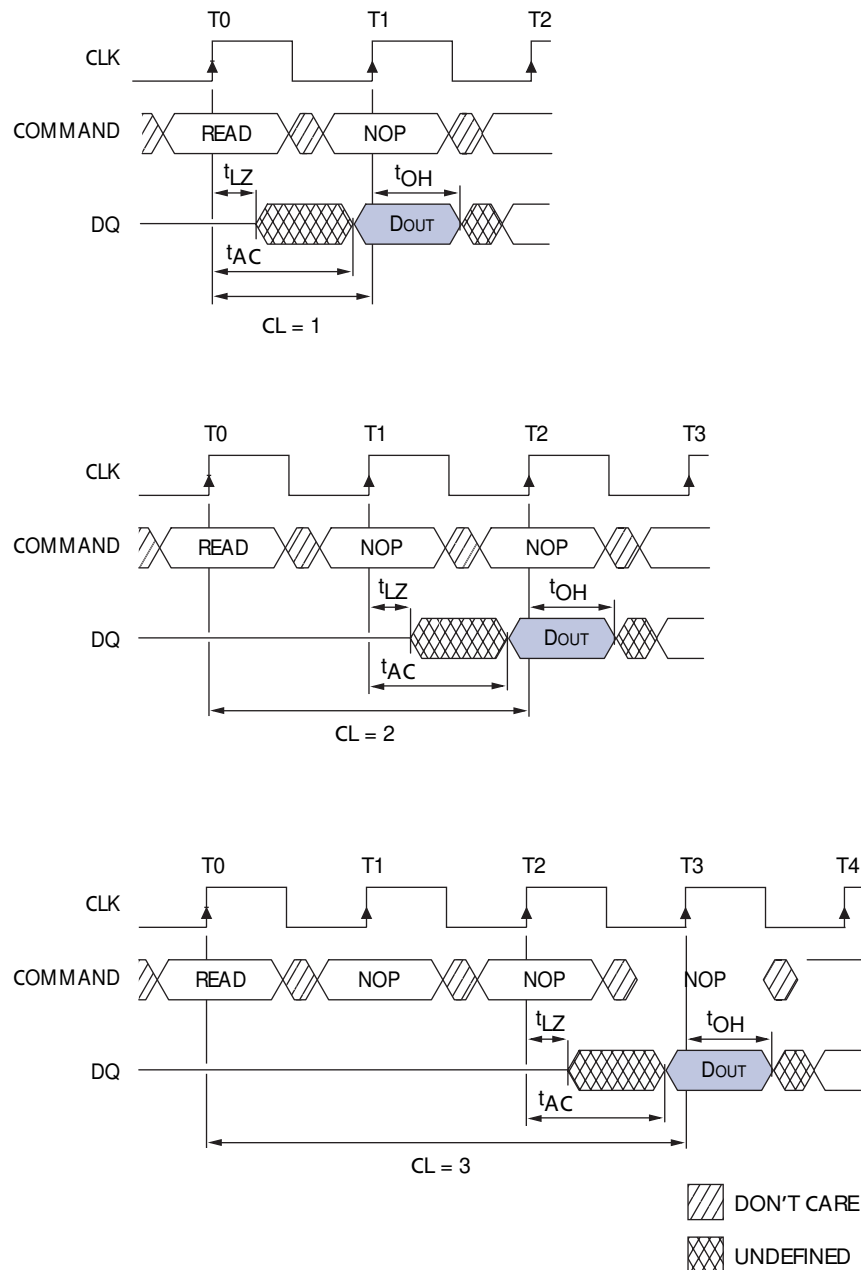


Upon completion of a burst, assuming no other commands have been initiated, the DQs will go High-Z. A full-page burst will continue until terminated. (At the end of the page, it will wrap to column 0 and continue.)

Data from any READ burst may be truncated with a subsequent READ command, and data from a fixed-length READ burst may be immediately followed by data from a READ command. In either case, a continuous flow of data can be maintained. The first data element from the new burst follows either the last element of a completed burst or the last desired data element of a longer burst that is being truncated. The new READ command should be issued  $x$  cycles before the clock edge at which the last desired data element is valid, where  $x$  equals the CAS latency minus one.

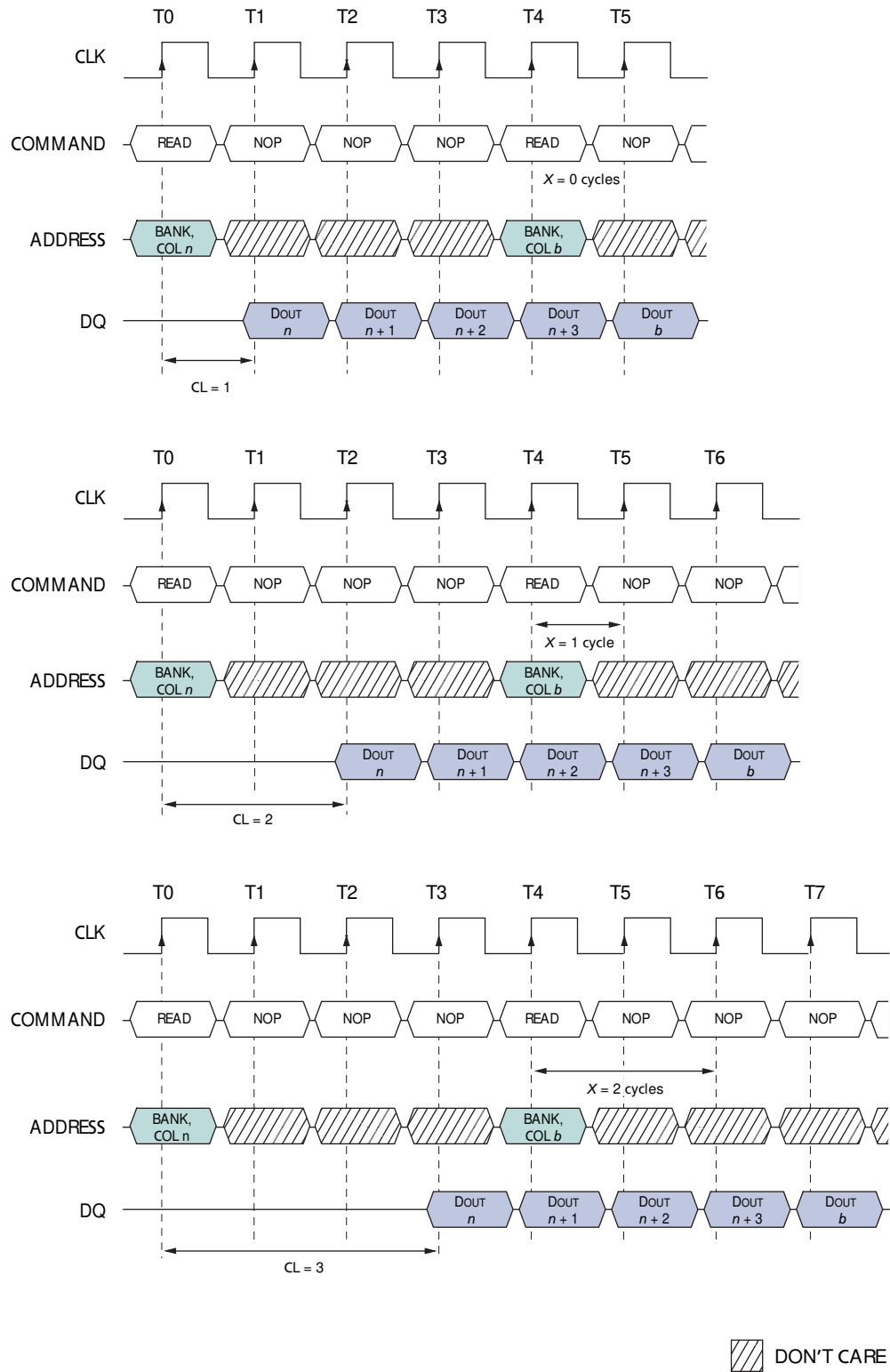


Figure 9: CAS Latency



This is shown in Figure 10 for CAS latencies of one, two, and three; data element  $n + 3$  is either the last of a burst of four or the last desired of a longer burst. The 256Mb SDRAM uses a pipelined architecture and therefore does not require the  $2n$  rule associated with a prefetch architecture. A READ command can be initiated on any clock cycle following a previous READ command. Full-speed random read accesses can be performed to the same bank, as shown in Figure 10 on page 25, or each subsequent READ may be performed to a different bank.

Figure 10: Consecutive READ Bursts



Note: Each READ command may be to either bank. DQM is LOW.