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PIC10F200/202/204/206

6-Pin, 8-Bit Flash Microcontrollers

Devices Included In This Data Sheet:

- PIC10F200
- PIC10F202
- PIC10F204
- PIC10F206

High-Performance RISC CPU:

- Only 33 Single-Word Instructions to Learn
- All Single-Cycle Instructions except for Program Branches, which are Two-Cycle
- 12-Bit Wide Instructions
- 2-Level Deep Hardware Stack
- Direct, Indirect and Relative Addressing modes for Data and Instructions
- 8-Bit Wide Data Path
- Eight Special Function Hardware Registers
- Operating Speed:
  - 4 MHz internal clock
  - 1 µs instruction cycle

Special Microcontroller Features:

- 4 MHz Precision Internal Oscillator:
  - Factory calibrated to ±1%
- In-Circuit Serial Programming™ (ICSP™)
- In-Circuit Debugging (ICD) Support
- Power-on Reset (POR)
- Device Reset Timer (DRT)
- Watchdog Timer (WDT) with Dedicated On-Chip RC Oscillator for Reliable Operation
- Programmable Code Protection
- Multiplexed $\overline{\text{MCLR}}$ Input Pin
- Internal Weak Pull-ups on I/O Pins
- Power-Saving Sleep mode
- Wake-up from Sleep on Pin Change

Low-Power Features/CMOS Technology:

- Operating Current:
  - < 175 µA @ 2V, 4 MHz, typical
- Standby Current:
  - 100 nA @ 2V, typical
- Low-Power, High-Speed Flash Technology:
  - 100,000 Flash endurance
  - > 40 year retention
- Fully Static Design
- Wide Operating Voltage Range: 2.0V to 5.5V
- Wide Temperature Range:
  - Industrial: -40°C to +85°C
  - Extended: -40°C to +125°C

Peripheral Features (PIC10F200/202):

- Four I/O Pins:
  - Three I/O pins with individual direction control
  - One input-only pin
  - High current sink/source for direct LED drive
  - Wake-on-change
  - Weak pull-ups
- 8-Bit Real-Time Clock/Counter (TMR0) with 8-Bit Programmable Prescaler

Peripheral Features (PIC10F204/206):

- Four I/O Pins:
  - Three I/O pins with individual direction control
  - One input-only pin
  - High current sink/source for direct LED drive
  - Wake-on-change
  - Weak pull-ups
- 8-Bit Real-Time Clock/Counter (TMR0) with 8-Bit Programmable Prescaler
- One Comparator:
  - Internal absolute voltage reference
  - Both comparator inputs visible externally
  - Comparator output visible externally

TABLE 1: PIC10F20X MEMORY AND FEATURES

Device	Program Memory	Data Memory	I/O	Timers 8-bit	Comparator
	Flash (words)	SRAM (bytes)			
PIC10F200	256	16	4	1	0
PIC10F202	512	24	4	1	0
PIC10F204	256	16	4	1	1
PIC10F206	512	24	4	1	1

Pin Diagrams

FIGURE 1: 6-PIN SOT-23

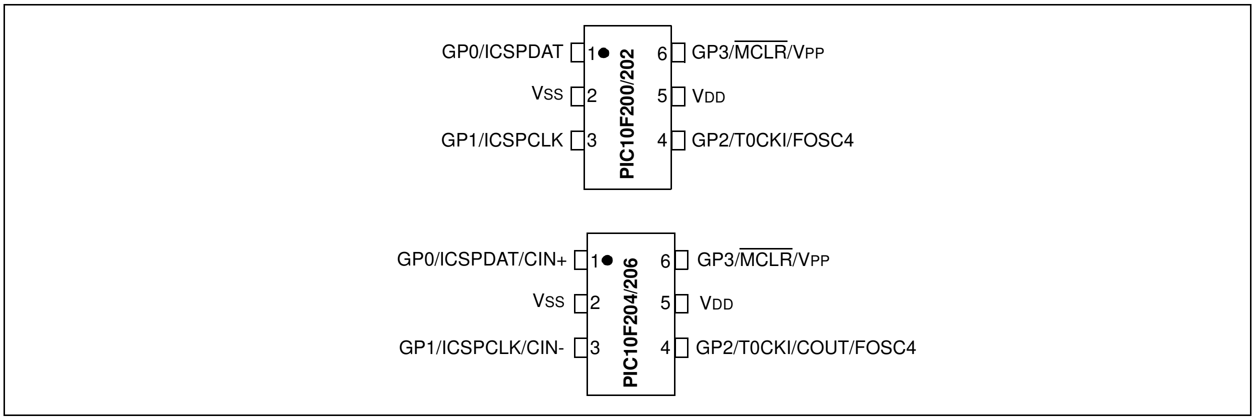


FIGURE 2: 8-PIN PDIP

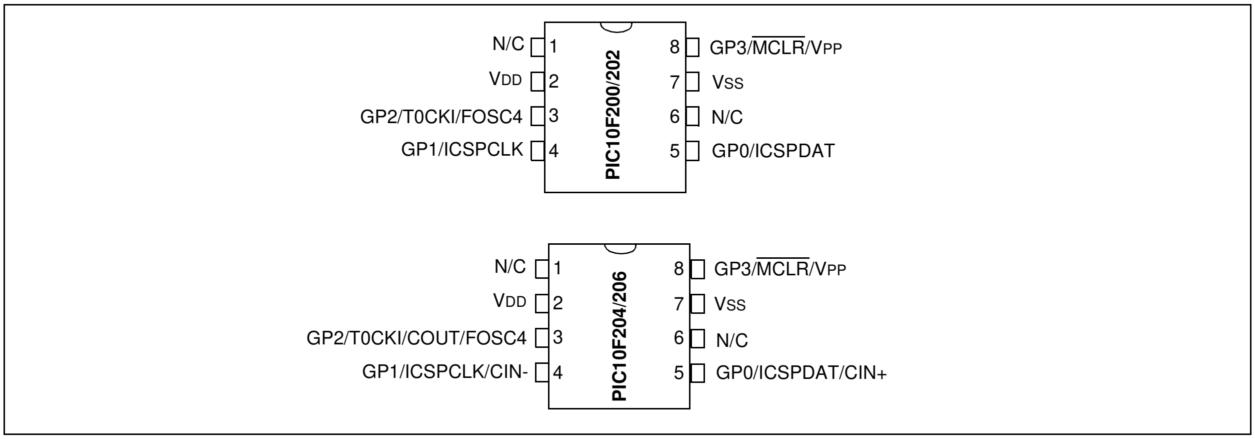
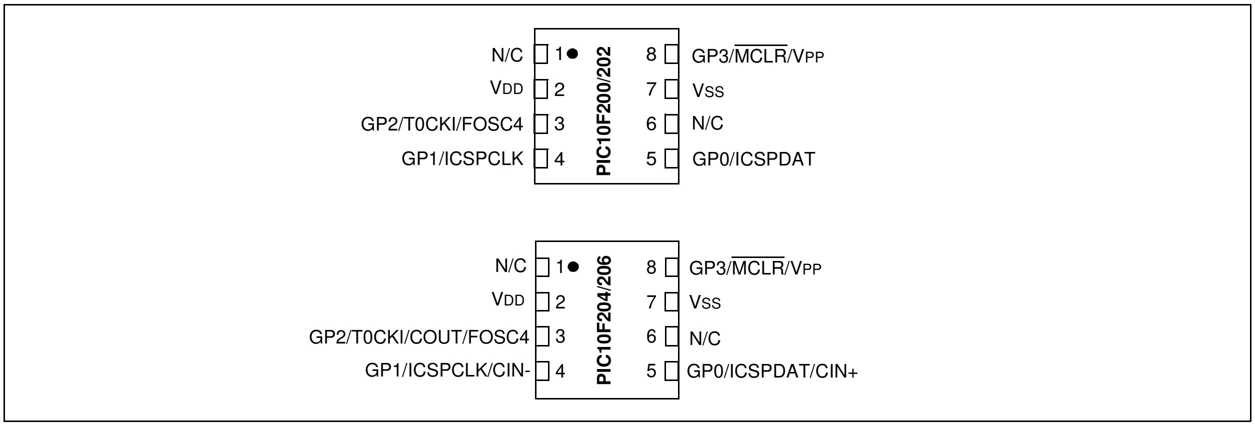


FIGURE 3: 8-PIN DFN



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1.0 GENERAL DESCRIPTION

The PIC10F200/202/204/206 devices from Microchip Technology are low-cost, high-performance, 8-bit, fully-static, Flash-based CMOS microcontrollers. They employ a RISC architecture with only 33 single-word/ single-cycle instructions. All instructions are single cycle (1 μs) except for program branches, which take two cycles. The PIC10F200/202/204/206 devices deliver performance in an order of magnitude higher than their competitors in the same price category. The 12-bit wide instructions are highly symmetrical, resulting in a typical 2:1 code compression over other 8-bit microcontrollers in its class. The easy-to-use and easy to remember instruction set reduces development time significantly.

The PIC10F200/202/204/206 products are equipped with special features that reduce system cost and power requirements. The Power-on Reset (POR) and Device Reset Timer (DRT) eliminate the need for external Reset circuitry. INTRC Internal Oscillator mode is provided, thereby preserving the limited number of I/O available. Power-Saving Sleep mode, Watchdog Timer and code protection features improve system cost, power and reliability.

The PIC10F200/202/204/206 devices are available in cost-effective Flash, which is suitable for production in any volume. The customer can take full advantage of Microchip's price leadership in Flash programmable microcontrollers, while benefiting from the Flash programmable flexibility.

The PIC10F200/202/204/206 products are supported by a full-featured macro assembler, a software simulator, an in-circuit debugger, a 'C' compiler, a low-cost development programmer and a full featured programmer. All the tools are supported on IBM® PC and compatible machines.

1.1 Applications

The PIC10F200/202/204/206 devices fit in applications ranging from personal care appliances and security systems to low-power remote transmitters/receivers. The Flash technology makes customizing application programs (transmitter codes, appliance settings, receiver frequencies, etc.) extremely fast and convenient. The small footprint packages, for through hole or surface mounting, make these microcontrollers well suited for applications with space limitations. Low cost, low power, high performance, ease-of-use and I/O flexibility make the PIC10F200/202/204/206 devices very versatile even in areas where no microcontroller use has been considered before (e.g., timer functions, logic and PLDs in larger systems and coprocessor applications).

TABLE 1-1: PIC10F200/202/204/206 DEVICES

		PIC10F200	PIC10F202	PIC10F204	PIC10F206
Clock	Maximum Frequency of Operation (MHz)	4	4	4	4
Memory	Flash Program Memory	256	512	256	512
	Data Memory (bytes)	16	24	16	24
Peripherals	Timer Module(s)	TMR0	TMR0	TMR0	TMR0
	Wake-up from Sleep on Pin Change	Yes	Yes	Yes	Yes
	Comparators	0	0	1	1
Features	I/O Pins	3	3	3	3
	Input-Only Pins	1	1	1	1
	Internal Pull-ups	Yes	Yes	Yes	Yes
	In-Circuit Serial Programming™	Yes	Yes	Yes	Yes
	Number of Instructions	33	33	33	33
	Packages	6-pin SOT-23 8-pin PDIP, DFN	6-pin SOT-23 8-pin PDIP, DFN	6-pin SOT-23 8-pin PDIP, DFN	6-pin SOT-23 8-pin PDIP, DFN

The PIC10F200/202/204/206 devices have Power-on Reset, selectable Watchdog Timer, selectable code-protect, high I/O current capability and precision internal oscillator.
The PIC10F200/202/204/206 devices use serial programming with data pin GP0 and clock pin GP1.

2.0 PIC10F200/202/204/206 DEVICE VARIETIES

A variety of packaging options are available. Depending on application and production requirements, the proper device option can be selected using the information in this section. When placing orders, please use the PIC10F200/202/204/206 Product Identification System at the back of this data sheet to specify the correct part number.

2.1 Quick Turn Programming (QTP) Devices

Microchip offers a QTP programming service for factory production orders. This service is made available for users who choose not to program medium-to-high quantity units and whose code patterns have stabilized. The devices are identical to the Flash devices but with all Flash locations and fuse options already programmed by the factory. Certain code and prototype verification procedures do apply before production shipments are available. Please contact your local Microchip Technology sales office for more details.

2.2 Serialized Quick Turn Programming^{SM} (SQTP^{SM}) Devices

Microchip offers a unique programming service, where a few user-defined locations in each device are programmed with different serial numbers. The serial numbers may be random, pseudo-random or sequential.

Serial programming allows each device to have a unique number, which can serve as an entry code, password or ID number.

3.0 ARCHITECTURAL OVERVIEW

The high performance of the PIC10F200/202/204/206 devices can be attributed to a number of architectural features commonly found in RISC microprocessors. To begin with, the PIC10F200/202/204/206 devices use a Harvard architecture in which program and data are accessed on separate buses. This improves bandwidth over traditional von Neumann architectures where program and data are fetched on the same bus. Separating program and data memory further allows instructions to be sized differently than the 8-bit wide data word. Instruction opcodes are 12 bits wide, making it possible to have all single-word instructions. A 12-bit wide program memory access bus fetches a 12-bit instruction in a single cycle. A two-stage pipeline overlaps fetch and execution of instructions. Consequently, all instructions (33) execute in a single cycle (1 μs @ 4 MHz) except for program branches.

The table below lists program memory (Flash) and data memory (RAM) for the PIC10F200/202/204/206 devices.

TABLE 3-1: PIC10F2XX MEMORY

Device	Memory	
	Program	Data
PIC10F200	256 x 12	16 x 8
PIC10F202	512 x 12	24 x 8
PIC10F204	256 x 12	16 x 8
PIC10F206	512 x 12	24 x 8

The PIC10F200/202/204/206 devices can directly or indirectly address its register files and data memory. All Special Function Registers (SFR), including the PC, are mapped in the data memory. The PIC10F200/202/204/206 devices have a highly orthogonal (symmetrical) instruction set that makes it possible to carry out any operation, on any register, using any addressing mode. This symmetrical nature and lack of “special optimal situations” make programming with the PIC10F200/202/204/206 devices simple, yet efficient. In addition, the learning curve is reduced significantly.

The PIC10F200/202/204/206 devices contain an 8-bit ALU and working register. The ALU is a general purpose arithmetic unit. It performs arithmetic and Boolean functions between data in the working register and any register file.

The ALU is 8 bits wide and capable of addition, subtraction, shift and logical operations. Unless otherwise mentioned, arithmetic operations are two’s complement in nature. In two-operand instructions, one operand is typically the W (working) register. The other operand is either a file register or an immediate constant. In single operand instructions, the operand is either the W register or a file register.

The W register is an 8-bit working register used for ALU operations. It is not an addressable register.

Depending on the instruction executed, the ALU may affect the values of the Carry (C), Digit Carry (DC) and Zero (Z) bits in the STATUS register. The C and DC bits operate as a $\overline{\text{borrow}}$ and $\overline{\text{digit borrow}}$ out bit, respectively, in subtraction. See the `SUBWF` and `ADDWF` instructions for examples.

A simplified block diagram is shown in Figure 3-1 and Figure 3-2, with the corresponding device pins described in Table 3-2.

FIGURE 3-1: PIC10F200/202 BLOCK DIAGRAM

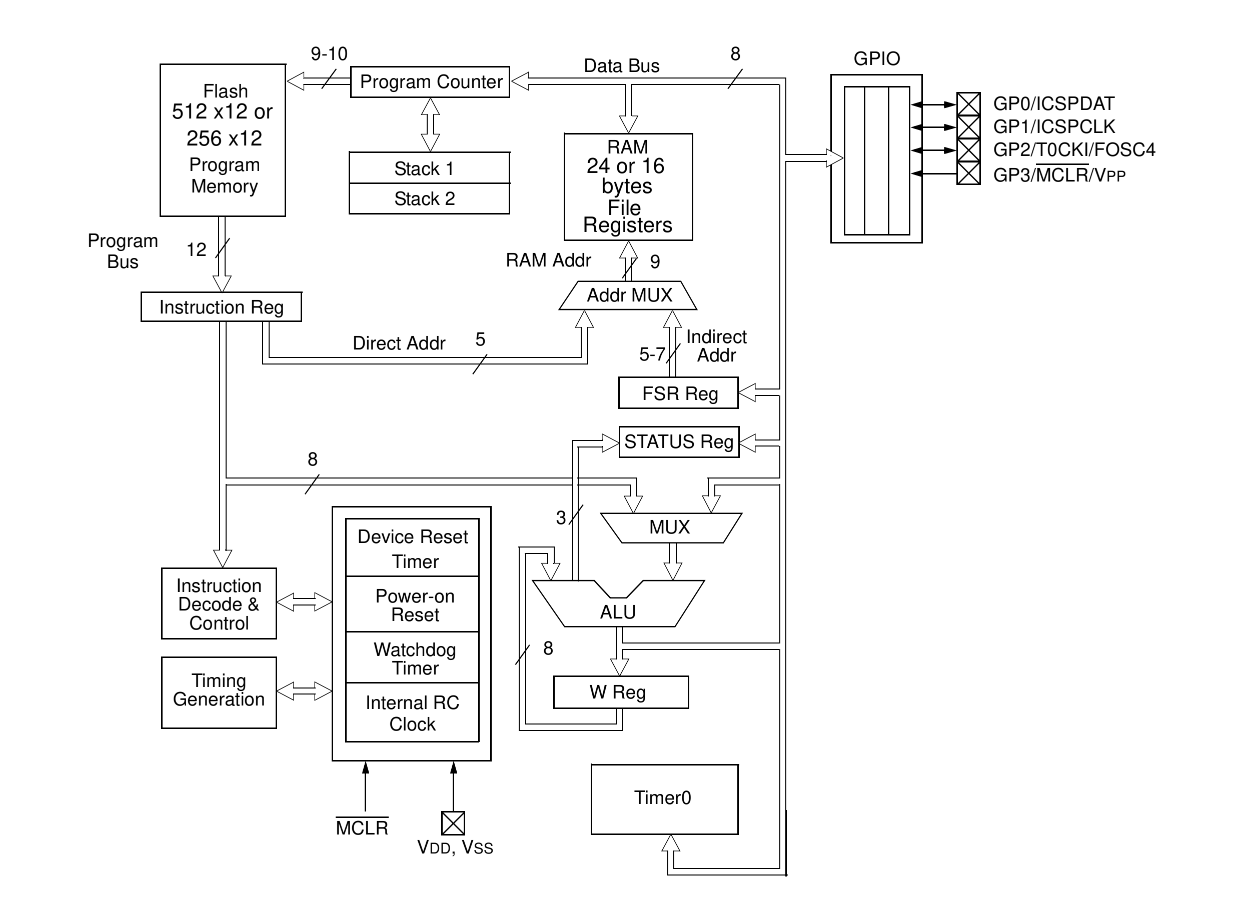


FIGURE 3-2: PIC10F204/206 BLOCK DIAGRAM

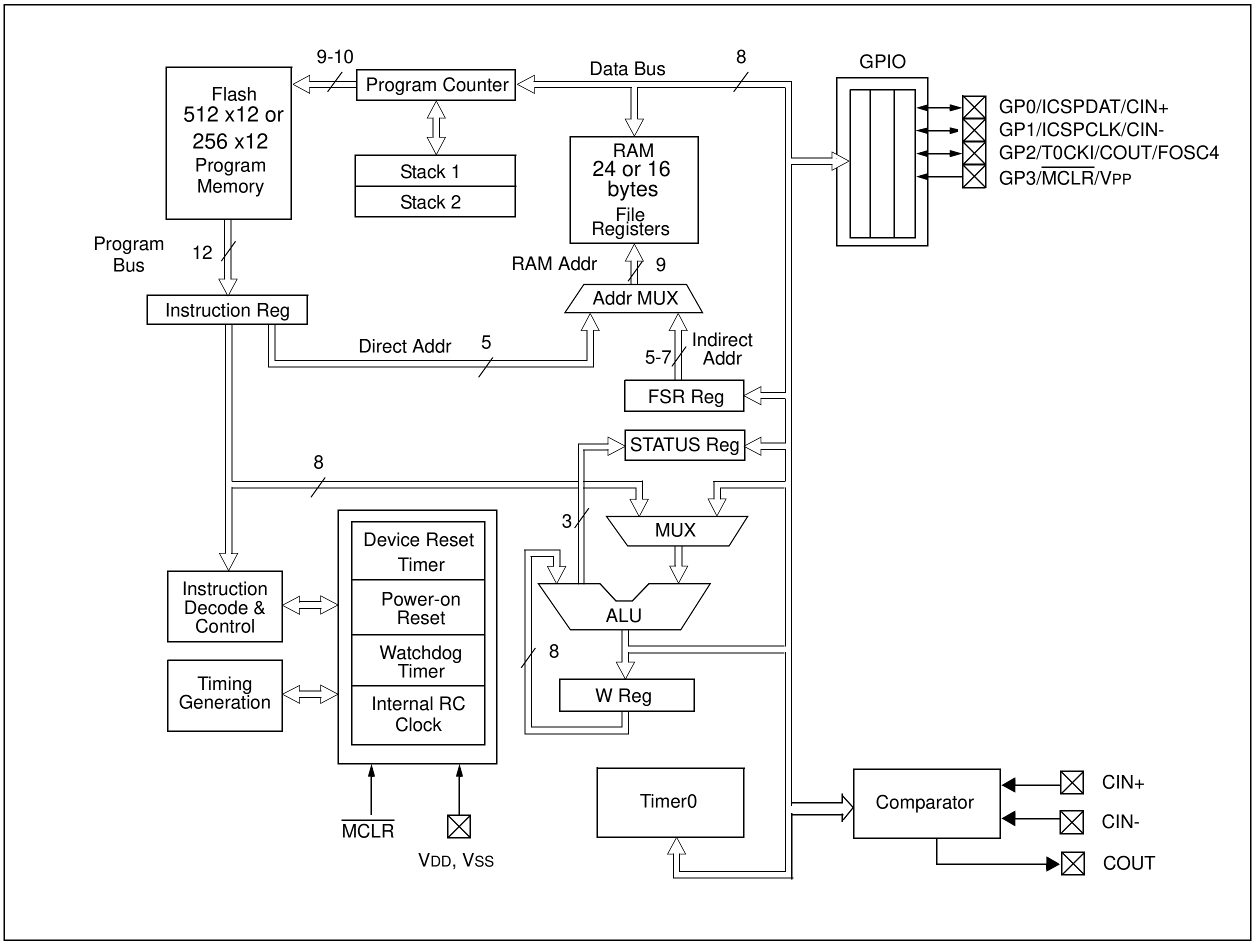


TABLE 3-2: PIC10F200/202/204/206 PINOUT DESCRIPTION

Name	Function	Input Type	Output Type	Description
GP0/ICSPDAT/CIN+	GP0	TTL	CMOS	Bidirectional I/O pin. Can be software programmed for internal weak pull-up and wake-up from Sleep on pin change.
	ICSPDAT	ST	CMOS	In-Circuit Serial Programming™ data pin.
	CIN+	AN	—	Comparator input (PIC10F204/206 only).
GP1/ICSPCLK/CIN-	GP1	TTL	CMOS	Bidirectional I/O pin. Can be software programmed for internal weak pull-up and wake-up from Sleep on pin change.
	ICSPCLK	ST	CMOS	In-Circuit Serial Programming clock pin.
	CIN-	AN	—	Comparator input (PIC10F204/206 only).
GP2/T0CKI/COUT/ FOSC4	GP2	TTL	CMOS	Bidirectional I/O pin.
	T0CKI	ST	—	Clock input to TMR0.
	COUT	—	CMOS	Comparator output (PIC10F204/206 only).
	FOSC4	—	CMOS	Oscillator/4 output.
GP3/$\overline{\text{MCLR}}$/VPP	GP3	TTL	—	Input pin. Can be software programmed for internal weak pull-up and wake-up from Sleep on pin change.
	$\overline{\text{MCLR}}$	ST	—	Master Clear (Reset). When configured as $\overline{\text{MCLR}}$, this pin is an active-low Reset to the device. Voltage on GP3/$\overline{\text{MCLR}}$/VPP must not exceed VDD during normal device operation or the device will enter Programming mode. Weak pull-up always on if configured as $\overline{\text{MCLR}}$.
	VPP	HV	—	Programming voltage input.
VDD	VDD	P	—	Positive supply for logic and I/O pins.
VSS	VSS	P	—	Ground reference for logic and I/O pins.

Legend: I = Input, O = Output, I/O = Input/Output, P = Power, — = Not used, TTL = TTL input, ST = Schmitt Trigger input, AN = Analog input

3.1 Clocking Scheme/Instruction Cycle

The clock is internally divided by four to generate four non-overlapping quadrature clocks, namely Q1, Q2, Q3 and Q4. Internally, the PC is incremented every Q1 and the instruction is fetched from program memory and latched into the instruction register in Q4. It is decoded and executed during the following Q1 through Q4. The clocks and instruction execution flow is shown in Figure 3-3 and Example 3-1.

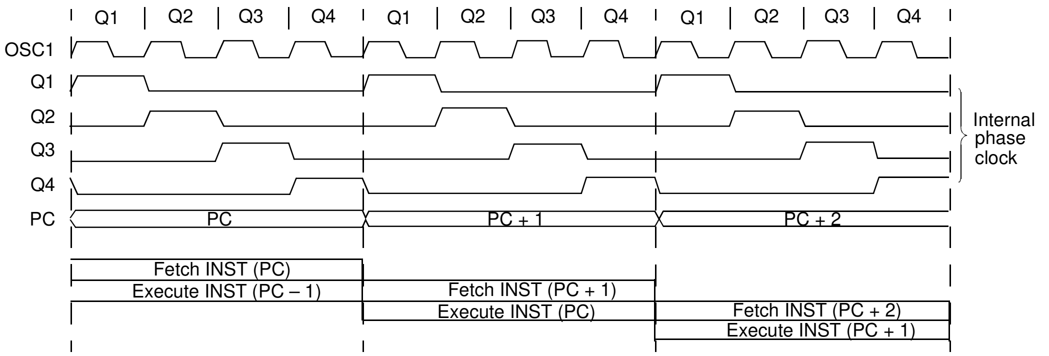
3.2 Instruction Flow/Pipelining

An instruction cycle consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle, while decode and execute take another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the PC to change (e.g., `GOTO`), then two cycles are required to complete the instruction (Example 3-1).

A fetch cycle begins with the PC incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the Instruction Register (IR) in cycle Q1. This instruction is then decoded and executed during the Q2, Q3 and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

FIGURE 3-3: CLOCK/INSTRUCTION CYCLE



EXAMPLE 3-1: INSTRUCTION PIPELINE FLOW

1. `MOVLW 03H`	Fetch 1	Execute 1				
2. `MOVWF GPIO`		Fetch 2	Execute 2			
3. `CALL SUB_1`			Fetch 3	Execute 3		
4. `BSF GPIO, BIT1`				Fetch 4	Flush	
					Fetch SUB_1	Execute SUB_1

All instructions are single cycle, except for any program branches. These take two cycles, since the fetch instruction is “flushed” from the pipeline, while the new instruction is being fetched and then executed.

4.0 MEMORY ORGANIZATION

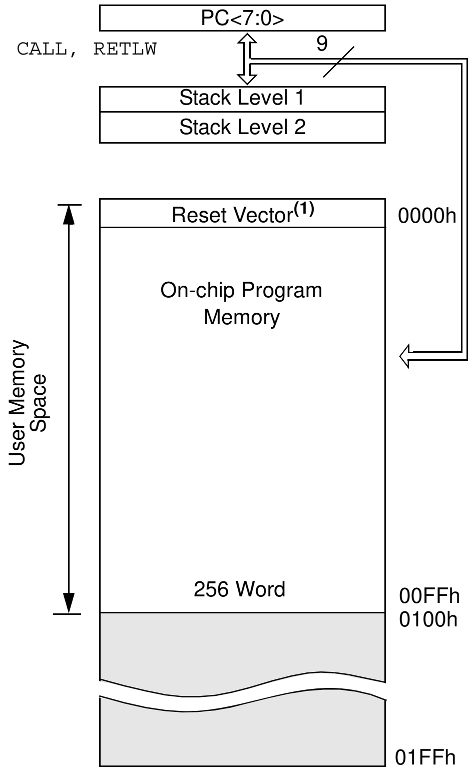
The PIC10F200/202/204/206 memories are organized into program memory and data memory. Data memory banks are accessed using the File Select Register (FSR).

4.1 Program Memory Organization for the PIC10F200/204

The PIC10F200/204 devices have a 9-bit Program Counter (PC) capable of addressing a 512 x 12 program memory space.

Only the first 256 x 12 (0000h-00FFh) for the PIC10F200/204 are physically implemented (see Figure 4-1). Accessing a location above these boundaries will cause a wraparound within the first 256 x 12 space (PIC10F200/204). The effective Reset vector is at 0000h (see Figure 4-1). Location 00FFh (PIC10F200/204) contains the internal clock oscillator calibration value. This value should never be overwritten.

FIGURE 4-1: PROGRAM MEMORY MAP AND STACK FOR THE PIC10F200/204



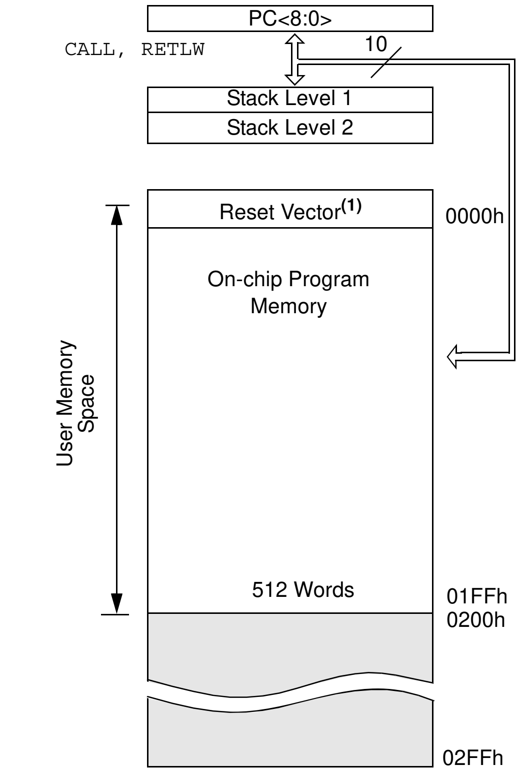
Note 1: Address 0000h becomes the effective Reset vector. Location 00FFh contains the `MOVLW XX` internal oscillator calibration value.

4.2 Program Memory Organization for the PIC10F202/206

The PIC10F202/206 devices have a 10-bit Program Counter (PC) capable of addressing a 1024 x 12 program memory space.

Only the first 512 x 12 (0000h-01FFh) for the PIC10F202/206 are physically implemented (see Figure 4-2). Accessing a location above these boundaries will cause a wraparound within the first 512 x 12 space (PIC10F202/206). The effective Reset vector is at 0000h (see Figure 4-2). Location 01FFh (PIC10F202/206) contains the internal clock oscillator calibration value. This value should never be overwritten.

FIGURE 4-2: PROGRAM MEMORY MAP AND STACK FOR THE PIC10F202/206



Note 1: Address 0000h becomes the effective Reset vector. Location 01FFh contains the `MOVLW XX` internal oscillator calibration value.

4.3 Data Memory Organization

Data memory is composed of registers or bytes of RAM. Therefore, data memory for a device is specified by its register file. The register file is divided into two functional groups: Special Function Registers (SFR) and General Purpose Registers (GPR).

The Special Function Registers include the TMR0 register, the Program Counter (PCL), the STATUS register, the I/O register (GPIO) and the File Select Register (FSR). In addition, Special Function Registers are used to control the I/O port configuration and prescaler options.

The General Purpose registers are used for data and control information under command of the instructions.

For the PIC10F200/204, the register file is composed of seven Special Function registers and 16 General Purpose registers (see Figure 4-3 and Figure 4-4).

For the PIC10F202/206, the register file is composed of eight Special Function registers and 24 General Purpose registers (see Figure 4-4).

4.3.1 GENERAL PURPOSE REGISTER FILE

The General Purpose Register file is accessed, either directly or indirectly, through the File Select Register (FSR). See **Section 4.9 “Indirect Data Addressing: INDF and FSR Registers”**.

FIGURE 4-3: PIC10F200/204 REGISTER FILE MAP

File Address	
00h	INDF[1]
01h	TMR0
02h	PCL
03h	STATUS
04h	FSR
05h	OSCCAL
06h	GPIO
07h	CMCON0[2]
08h 0Fh	Unimplemented[3]
10h 1Fh	General Purpose Registers

Note 1: Not a physical register. See **Section 4.9 "Indirect Data Addressing: INDF and FSR Registers"**.

2: PIC10F204 only. Unimplemented on the PIC10F200 and reads as 00h.

3: Unimplemented, read as 00h.

FIGURE 4-4: PIC10F202/206 REGISTER FILE MAP

File Address	
00h	INDF[1]
01h	TMR0
02h	PCL
03h	STATUS
04h	FSR
05h	OSCCAL
06h	GPIO
07h	CMCON0[2]
08h 1Fh	General Purpose Registers

Note 1: Not a physical register. See **Section 4.9 "Indirect Data Addressing: INDF and FSR Registers"**.

2: PIC10F206 only. Unimplemented on the PIC10F202 and reads as 00h.

4.3.2 SPECIAL FUNCTION REGISTERS

The Special Function Registers (SFRs) are registers used by the CPU and peripheral functions to control the operation of the device (Table 4-1).

The Special Function Registers can be classified into two sets. The Special Function Registers associated with the “core” functions are described in this section. Those related to the operation of the peripheral features are described in the section for each peripheral feature.

TABLE 4-1: SPECIAL FUNCTION REGISTER (SFR) SUMMARY (PIC10F200/202/204/206)

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-On Reset[(2)]	Register on Page
00h	INDF	Uses Contents of FSR to Address Data Memory (not a physical register)								`xxxx xxxx`	19
01h	TMR0	8-bit Real-Time Clock/Counter								`xxxx xxxx`	23, 27
02h[(1)]	PCL	Low-order 8 bits of PC								`1111 1111`	18
03h	STATUS	GPWUF	CWUF[(5)]	—	$\overline{\text{TO}}$	$\overline{\text{PD}}$	Z	DC	C	`00-1 1xxx`[(3)]	15
04h	FSR	Indirect Data Memory Address Pointer								`111x xxxx`	19
05h	OSCCAL	CAL6	CAL5	CAL4	CAL3	CAL2	CAL1	CAL0	FOSC4	`1111 1110`	17
06h	GPIO	—	—	—	—	GP3	GP2	GP1	GP0	`---- xxxx`	20
07h[(4)]	CMCON0	CMPOUT	$\overline{\text{COUTEN}}$	POL	$\overline{\text{CMPT0CS}}$	CMPON	CNREF	CPREF	$\overline{\text{CWU}}$	`1111 1111`	28
N/A	TRISGPIO	—	—	—	—	I/O Control Register				`---- 1111`	31
N/A	OPTION	$\overline{\text{GPWU}}$	$\overline{\text{GPPU}}$	T0CS	T0SE	PSA	PS2	PS1	PS0	`1111 1111`	16

Legend: – = unimplemented, read as ‘0’, x = unknown, u = unchanged, q = value depends on condition.

Note 1: The upper byte of the Program Counter is not directly accessible. See **Section 4.7 “Program Counter”** for an explanation of how to access these bits.

2: Other (non Power-up) Resets include external Reset through $\overline{\text{MCLR}}$, Watchdog Timer and wake-up on pin change Reset.

3: See Table 9-1 for other Reset specific values.

4: PIC10F204/206 only.

5: PIC10F204/206 only. On all other devices, this bit is reserved and should not be used.

4.4 STATUS Register

This register contains the arithmetic status of the ALU, the Reset status and the page preselect bit.

The STATUS register can be the destination for any instruction, as with any other register. If the STATUS register is the destination for an instruction that affects the Z, DC or C bits, then the write to these three bits is disabled. These bits are set or cleared according to the device logic. Furthermore, the $\overline{\text{TO}}$ and $\overline{\text{PD}}$ bits are not writable. Therefore, the result of an instruction with the STATUS register as destination may be different than intended.

For example, `CLRF STATUS`, will clear the upper three bits and set the Z bit. This leaves the STATUS register as `000u u1uu` (where `u` = unchanged).

Therefore, it is recommended that only `BCF`, `BSF` and `MOVWF` instructions be used to alter the STATUS register. These instructions do not affect the Z, DC or C bits from the STATUS register. For other instructions which do affect Status bits, see **Section 10.0 "Instruction Set Summary"**.

REGISTER 4-1: STATUS REGISTER

R/W-0	R/W-0	U-1	R-1	R-1	R/W-x	R/W-x	R/W-x
GPWUF	CWUF[1]	—	$\overline{\text{TO}}$	$\overline{\text{PD}}$	Z	DC	C
bit 7							bit 0

Legend:

R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 7 **GPWUF:** GPIO Reset bit
1 = Reset due to wake-up from Sleep on pin change
0 = After power-up or other Reset

bit 6 **CWUF:** Comparator Wake-up on Change Flag bit[1]
1 = Reset due to wake-up from Sleep on comparator change
0 = After power-up or other Reset conditions.

bit 5 **Reserved:** Do not use. Use of this bit may affect upward compatibility with future products.

bit 4 **$\overline{\text{TO}}$:** Time-out bit
1 = After power-up, `CLRWDT` instruction or `SLEEP` instruction
0 = A WDT time-out occurred

bit 3 **$\overline{\text{PD}}$:** Power-down bit
1 = After power-up or by the `CLRWDT` instruction
0 = By execution of the `SLEEP` instruction

bit 2 **Z:** Zero bit
1 = The result of an arithmetic or logic operation is zero
0 = The result of an arithmetic or logic operation is not zero

bit 1 **DC:** Digit Carry/$\overline{\text{Borrow}}$ bit (for `ADDWF` and `SUBWF` instructions)
`ADDWF`:
1 = A carry from the 4th low-order bit of the result occurred
0 = A carry from the 4th low-order bit of the result did not occur
`SUBWF`:
1 = A borrow from the 4th low-order bit of the result did not occur
0 = A borrow from the 4th low-order bit of the result occurred

bit 0 **C:** Carry/$\overline{\text{Borrow}}$ bit (for `ADDWF`, `SUBWF` and `RRF`, `RLF` instructions)

`ADDWF`:	`SUBWF`:	`RRF` or `RLF`:
1 = A carry occurred	1 = A borrow did not occur	Load bit with LSb or MSb, respectively
0 = A carry did not occur	0 = A borrow occurred	

Note 1: This bit is used on the PIC10F204/206. For code compatibility do not use this bit on the PIC10F200/202.

4.5 OPTION Register

The OPTION register is a 8-bit wide, write-only register, which contains various control bits to configure the Timer0/WDT prescaler and Timer0.

By executing the `OPTION` instruction, the contents of the W register will be transferred to the OPTION register. A Reset sets the OPTION<7:0> bits.

Note:	If TRIS bit is set to ‘`0`’, the wake-up on change and pull-up functions are disabled for that pin (i.e., note that TRIS overrides Option control of $\overline{\text{GPPU}}$ and $\overline{\text{GPWU}}$).

Note:	If the T0CS bit is set to ‘`1`’, it will override the TRIS function on the T0CKI pin.

REGISTER 4-2: OPTION REGISTER

W-1	W-1	W-1	W-1	W-1	W-1	W-1	W-1
$\overline{\text{GPWU}}$	$\overline{\text{GPPU}}$	T0CS	T0SE	PSA	PS2	PS1	PS0
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
-n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

bit 7 **$\overline{\text{GPWU}}$:** Enable Wake-up on Pin Change bit (GP0, GP1, GP3)
`1` = Disabled
`0` = Enabled

bit 6 **$\overline{\text{GPPU}}$:** Enable Weak Pull-ups bit (GP0, GP1, GP3)
`1` = Disabled
`0` = Enabled

bit 5 **T0CS:** Timer0 Clock Source Select bit
`1` = Transition on T0CKI pin (overrides TRIS on the T0CKI pin)
`0` = Transition on internal instruction cycle clock, FOSC/4

bit 4 **T0SE:** Timer0 Source Edge Select bit
`1` = Increment on high-to-low transition on the T0CKI pin
`0` = Increment on low-to-high transition on the T0CKI pin

bit 3 **PSA:** Prescaler Assignment bit
`1` = Prescaler assigned to the WDT
`0` = Prescaler assigned to Timer0

bit 2-0 **PS<2:0>:** Prescaler Rate Select bits

Bit Value	Timer0 Rate	WDT Rate
`000`	1 : 2	1 : 1
`001`	1 : 4	1 : 2
`010`	1 : 8	1 : 4
`011`	1 : 16	1 : 8
`100`	1 : 32	1 : 16
`101`	1 : 64	1 : 32
`110`	1 : 128	1 : 64
`111`	1 : 256	1 : 128

.

4.6 OSCCAL Register

The Oscillator Calibration (OSCCAL) register is used to calibrate the internal precision 4 MHz oscillator. It contains seven bits for calibration.

> **Note:** Erasing the device will also erase the pre-programmed internal calibration value for the internal oscillator. The calibration value must be read prior to erasing the part so it can be reprogrammed correctly later.

After you move in the calibration constant, do not change the value. See **Section 9.2.2 "Internal 4 MHz Oscillator"**.

REGISTER 4-3: OSCCAL REGISTER

R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-0
CAL6	CAL5	CAL4	CAL3	CAL2	CAL1	CAL0	FOSC4
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 7-1 **CAL<6:0>:** Oscillator Calibration bits

```
0111111 = Maximum frequency
•
•
•
0000001
0000000 = Center frequency
1111111
•
•
•
1000000 = Minimum frequency
```

bit 0 **FOSC4:** INTOSC/4 Output Enable bit[1]

1 = INTOSC/4 output onto GP2
0 = GP2/T0CKI/COUT applied to GP2

Note 1: Overrides GP2/T0CKI/COUT control registers when enabled.

4.7 Program Counter

As a program instruction is executed, the Program Counter (PC) will contain the address of the next program instruction to be executed. The PC value is increased by one every instruction cycle, unless an instruction changes the PC.

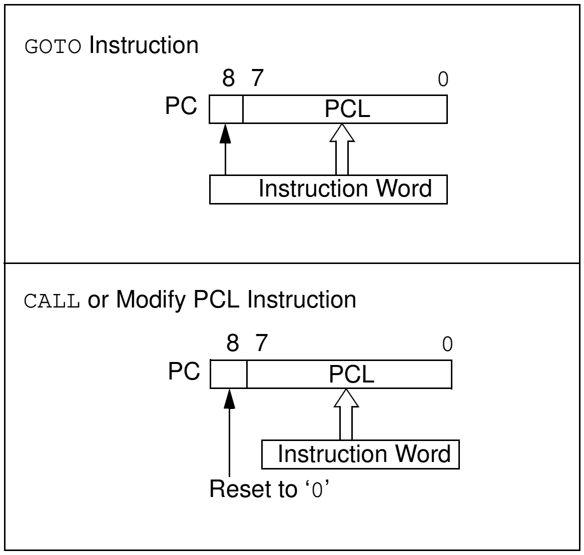
For a GOTO instruction, bits 8-0 of the PC are provided by the GOTO instruction word. The Program Counter Low (PCL) is mapped to PC<7:0>.

For a CALL instruction, or any instruction where the PCL is the destination, bits 7:0 of the PC again are provided by the instruction word. However, PC<8> does not come from the instruction word, but is always cleared (Figure 4-5).

Instructions where the PCL is the destination, or modify PCL instructions, include MOVWF PC, ADDWF PC and BSF PC,5.

Note: Because PC<8> is cleared in the CALL instruction or any modify PCL instruction, all subroutine calls or computed jumps are limited to the first 256 locations of any program memory page (512 words long).

FIGURE 4-5: LOADING OF PC BRANCH INSTRUCTIONS



4.7.1 EFFECTS OF RESET

The PC is set upon a Reset, which means that the PC addresses the last location in program memory (i.e., the oscillator calibration instruction). After executing MOVLW XX, the PC will roll over to location 0000h and begin executing user code.

4.8 Stack

The PIC10F200/204 devices have a 2-deep, 8-bit wide hardware PUSH/POP stack.

The PIC10F202/206 devices have a 2-deep, 9-bit wide hardware PUSH/POP stack.

A CALL instruction will PUSH the current value of Stack 1 into Stack 2 and then PUSH the current PC value, incremented by one, into Stack Level 1. If more than two sequential CALLs are executed, only the most recent two return addresses are stored.

A RETLW instruction will POP the contents of Stack Level 1 into the PC and then copy Stack Level 2 contents into level 1. If more than two sequential RETLWs are executed, the stack will be filled with the address previously stored in Stack Level 2.

Note 1: The W register will be loaded with the literal value specified in the instruction. This is particularly useful for the implementation of the data look-up tables within the program memory.

2: There are no Status bits to indicate stack overflows or stack underflow conditions.

3: There are no instruction mnemonics called PUSH or POP. These are actions that occur from the execution of the CALL and RETLW instructions.

4.9 Indirect Data Addressing: INDF and FSR Registers

The INDF register is not a physical register. Addressing INDF actually addresses the register whose address is contained in the FSR register (FSR is a *pointer*). This is indirect addressing.

4.10 Indirect Addressing

- Register file 09 contains the value 10h
- Register file 0A contains the value 0Ah
- Load the value 09 into the FSR register
- A read of the INDF register will return the value of 10h
- Increment the value of the FSR register by one (FSR = 0A)
- A read of the INDR register now will return the value of 0Ah.

Reading INDF itself indirectly (FSR = 0) will produce 00h. Writing to the INDF register indirectly results in a no operation (although Status bits may be affected).

A simple program to clear RAM locations 10h-1Fh using indirect addressing is shown in Example 4-1.

EXAMPLE 4-1: HOW TO CLEAR RAM USING INDIRECT ADDRESSING

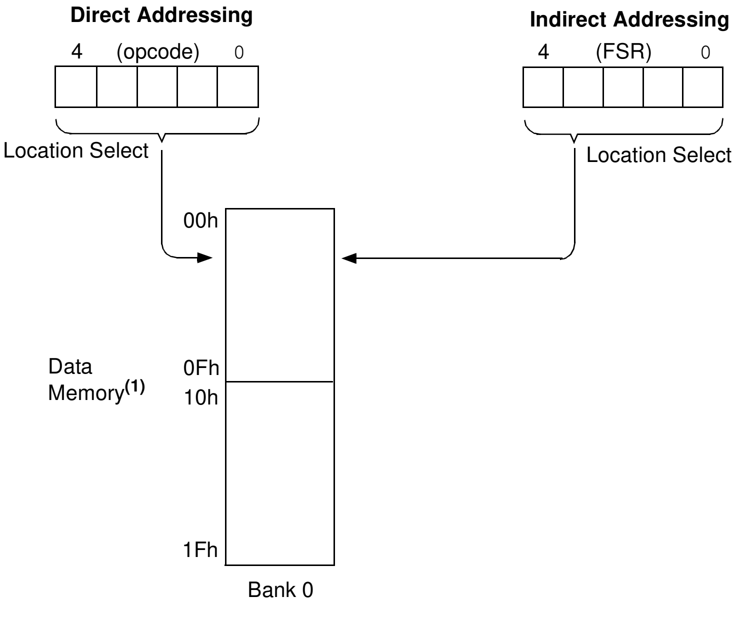
```
        MOVLW   0x10    ;initialize pointer
        MOVWF   FSR     ;to RAM
NEXT    CLRF    INDF    ;clear INDF
                        ;register
        INCF    FSR,F   ;inc pointer
        BTFSC   FSR,4   ;all done?
        GOTO    NEXT    ;NO, clear next
CONTINUE
                :       ;YES, continue
                :
```

The FSR is a 5-bit wide register. It is used in conjunction with the INDF register to indirectly address the data memory area.

The FSR<4:0> bits are used to select data memory addresses 00h to 1Fh.

> **Note:** PIC10F200/202/204/206 – Do not use banking. FSR <7:5> are unimplemented and read as '1's.

FIGURE 4-6: DIRECT/INDIRECT ADDRESSING (PIC10F200/202/204/206)



Note 1: For register map detail, see **Section 4.3 "Data Memory Organization"**.

5.0 I/O PORT

As with any other register, the I/O register(s) can be written and read under program control. However, read instructions (e.g., `MOVF GPIO, W`) always read the I/O pins independent of the pin's Input/Output modes. On Reset, all I/O ports are defined as input (inputs are at high-impedance) since the I/O control registers are all set.

5.1 GPIO

GPIO is an 8-bit I/O register. Only the low-order 4 bits are used (GP<3:0>). Bits 7 through 4 are unimplemented and read as '0's. Please note that GP3 is an input-only pin. Pins GP0, GP1 and GP3 can be configured with weak pull-ups and also for wake-up on change. The wake-up on change and weak pull-up functions are not pin selectable. If GP3/MCLR is configured as MCLR, weak pull-up is always on and wake-up on change for this pin is not enabled.

5.2 TRIS Registers

The Output Driver Control register is loaded with the contents of the W register by executing the `TRIS f` instruction. A '1' from a TRIS register bit puts the corresponding output driver in a High-Impedance mode. A '0' puts the contents of the output data latch on the selected pins, enabling the output buffer. The exceptions are GP3, which is input-only and the GP2/T0CKI/COUT/FOSC4 pin, which may be controlled by various registers. See Table 5-1.

> **Note:** A read of the ports reads the pins, not the output data latches. That is, if an output driver on a pin is enabled and driven high, but the external system is holding it low, a read of the port will indicate that the pin is low.

The TRIS registers are "write-only" and are set (output drivers disabled) upon Reset.

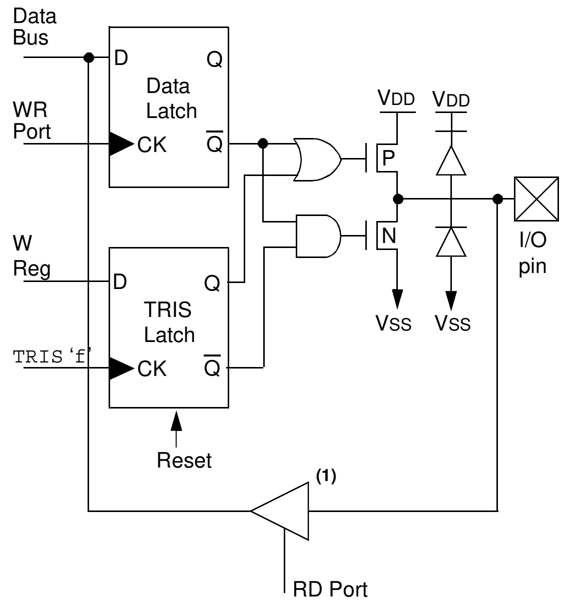
TABLE 5-1: ORDER OF PRECEDENCE FOR PIN FUNCTIONS

Priority	GP0	GP1	GP2	GP3
1	CIN+	CIN-	FOSC4	I/MCLR
2	TRIS GPIO	TRIS GPIO	COUT	—
3	—	—	T0CKI	—
4	—	—	TRIS GPIO	—

5.3 I/O Interfacing

The equivalent circuit for an I/O port pin is shown in Figure 5-1. All port pins, except GP3 which is input-only, may be used for both input and output operations. For input operations, these ports are non-latching. Any input must be present until read by an input instruction (e.g., `MOVF GPIO, W`). The outputs are latched and remain unchanged until the output latch is rewritten. To use a port pin as output, the corresponding direction control bit in TRIS must be cleared (= 0). For use as an input, the corresponding TRIS bit must be set. Any I/O pin (except GP3) can be programmed individually as input or output.

FIGURE 5-1: PIC10F200/202/204/206 EQUIVALENT CIRCUIT FOR A SINGLE I/O PIN



Note 1: See Table 3-2 for buffer type.

TABLE 5-2: SUMMARY OF PORT REGISTERS

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-On Reset	Value on All Other Resets
N/A	TRISGPIO	—	—	—	—	I/O Control Register				`---- 1111`	`---- 1111`
N/A	OPTION	$\overline{\text{GPWU}}$	$\overline{\text{GPPU}}$	T0CS	T0SE	PSA	PS2	PS1	PS0	`1111 1111`	`1111 1111`
03h	STATUS	GPWUF	CWUF	—	$\overline{\text{TO}}$	$\overline{\text{PD}}$	Z	DC	C	`00-1 1xxx`	`qq-q quuu`(1), (2)
06h	GPIO	—	—	—	—	GP3	GP2	GP1	GP0	`---- xxxx`	`---- uuuu`

Legend: Shaded cells are not used by PORT registers, read as '0', – = unimplemented, read as '0', x = unknown, u = unchanged, q = depends on condition.

Note 1: If Reset was due to wake-up on pin change, then bit 7 = 1. All other Resets will cause bit 7 = 0.

2: If Reset was due to wake-up on comparator change, then bit 6 = 1. All other Resets will cause bit 6 = 0.

5.4 I/O Programming Considerations

5.4.1 BIDIRECTIONAL I/O PORTS

Some instructions operate internally as read followed by write operations. The BCF and BSF instructions, for example, read the entire port into the CPU, execute the bit operation and rewrite the result. Caution must be used when these instructions are applied to a port where one or more pins are used as input/outputs. For example, a BSF operation on bit 2 of GPIO will cause all eight bits of GPIO to be read into the CPU, bit 2 to be set and the GPIO value to be written to the output latches. If another bit of GPIO is used as a bidirectional I/O pin (say bit 0), and it is defined as an input at this time, the input signal present on the pin itself would be read into the CPU and rewritten to the data latch of this particular pin, overwriting the previous content. As long as the pin stays in the Input mode, no problem occurs. However, if bit 0 is switched into Output mode later on, the content of the data latch may now be unknown.

Example 5-1 shows the effect of two sequential Read-Modify-Write instructions (e.g., BCF, BSF, etc.) on an I/O port.

A pin actively outputting a high or a low should not be driven from external devices at the same time in order to change the level on this pin ("wired OR", "wired AND"). The resulting high output currents may damage the chip.

EXAMPLE 5-1: READ-MODIFY-WRITE INSTRUCTIONS ON AN I/O PORT

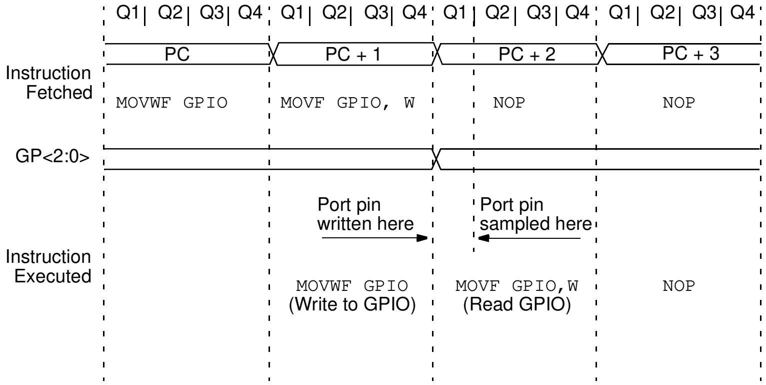
```
;Initial GPIO Settings
;GPIO<3:2> Inputs
;GPIO<1:0> Outputs
;
;                  GPIO latch    GPIO pins
;                  ----------    ----------
 BCF    GPIO,   1 ;---- pp01     ---- pp11
 BCF    GPIO,   0 ;---- pp10     ---- pp11
 MOVLW  007h;
 TRIS   GPIO      ;---- pp10     ---- pp11
;
```

Note 1: The user may have expected the pin values to be ---- pp00. The 2nd BCF caused GP1 to be latched as the pin value (High).

5.4.2 SUCCESSIVE OPERATIONS ON I/O PORTS

The actual write to an I/O port happens at the end of an instruction cycle, whereas for reading, the data must be valid at the beginning of the instruction cycle (Figure 5-2). Therefore, care must be exercised if a write followed by a read operation is carried out on the same I/O port. The sequence of instructions should allow the pin voltage to stabilize (load dependent) before the next instruction causes that file to be read into the CPU. Otherwise, the previous state of that pin may be read into the CPU rather than the new state. When in doubt, it is better to separate these instructions with a NOP or another instruction not accessing this I/O port.

FIGURE 5-2: SUCCESSIVE I/O OPERATION (PIC10F200/202/204/206)



This example shows a write to GPIO followed by a read from GPIO.

Data setup time = (0.25 TCY – TPD)

where: TCY = instruction cycle

TPD = propagation delay

Therefore, at higher clock frequencies, a write followed by a read may be problematic.

6.0 TIMER0 MODULE AND TMR0 REGISTER (PIC10F200/202)

The Timer0 module has the following features:

- 8-bit timer/counter register, TMR0
- Readable and writable
- 8-bit software programmable prescaler
- Internal or external clock select:
  - Edge select for external clock

Figure 6-1 is a simplified block diagram of the Timer0 module.

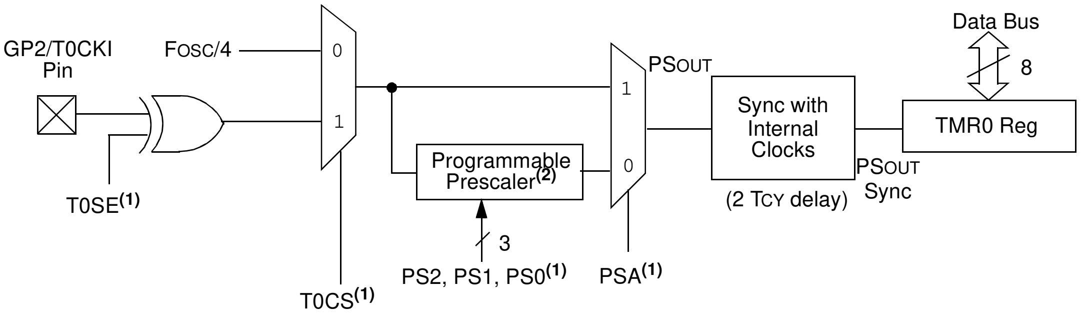
Timer mode is selected by clearing the T0CS bit (OPTION<5>). In Timer mode, the Timer0 module will increment every instruction cycle (without prescaler). If TMR0 register is written, the increment is inhibited for the following two cycles (Figure 6-2 and Figure 6-3). The user can work around this by writing an adjusted value to the TMR0 register.

Counter mode is selected by setting the T0CS bit (OPTION<5>). In this mode, Timer0 will increment either on every rising or falling edge of pin T0CKI. The T0SE bit (OPTION<4>) determines the source edge. Clearing the T0SE bit selects the rising edge. Restrictions on the external clock input are discussed in detail in **Section 6.1 "Using Timer0 with an External Clock (PIC10F200/202)"**.

The prescaler may be used by either the Timer0 module or the Watchdog Timer, but not both. The prescaler assignment is controlled in software by the control bit, PSA (OPTION<3>). Clearing the PSA bit will assign the prescaler to Timer0. The prescaler is not readable or writable. When the prescaler is assigned to the Timer0 module, prescale values of 1:2, 1:4, 1:256 are selectable. **Section 6.2 "Prescaler"** details the operation of the prescaler.

A summary of registers associated with the Timer0 module is found in Table 6-1.

FIGURE 6-1: TIMER0 BLOCK DIAGRAM



Note 1: Bits T0CS, T0SE, PSA, PS2, PS1 and PS0 are located in the OPTION register.

2: The prescaler is shared with the Watchdog Timer (Figure 6-5).

FIGURE 6-2: TIMER0 TIMING: INTERNAL CLOCK/NO PRESCALE

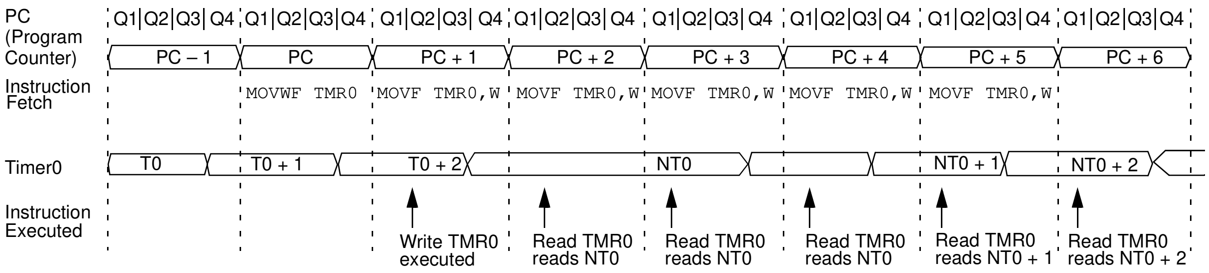


FIGURE 6-3: TIMER0 TIMING: INTERNAL CLOCK/PRESCALE 1:2

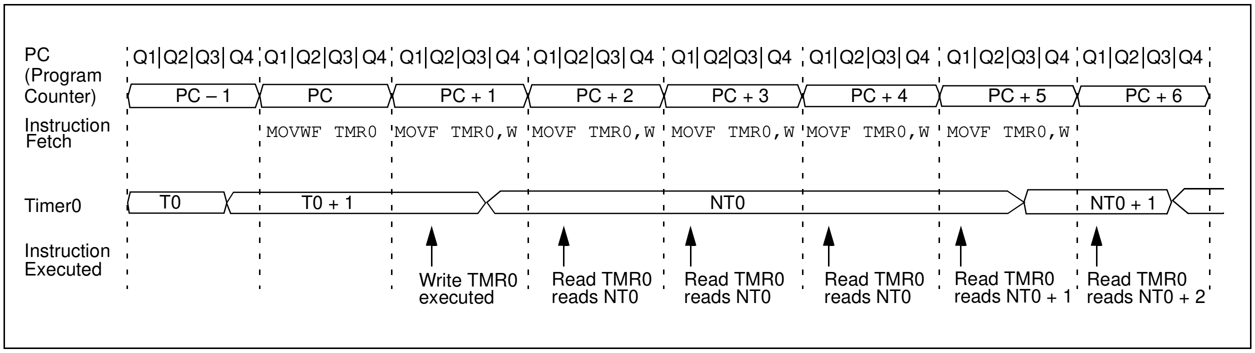


TABLE 6-1: REGISTERS ASSOCIATED WITH TIMER0

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-On Reset	Value on All Other Resets
01h	TMR0	Timer0 – 8-bit Real-Time Clock/Counter								xxxx xxxx	uuuu uuuu
N/A	OPTION	GPWU	GPPU	T0CS	T0SE	PSA	PS2	PS1	PS0	1111 1111	1111 1111
N/A	TRISGPIO(1)	—	—	—	—	I/O Control Register				---- 1111	---- 1111

Legend: Shaded cells not used by Timer0. – = unimplemented, x = unknown, u = unchanged.
Note 1: The TRIS of the T0CKI pin is overridden when T0CS = 1.

6.1 Using Timer0 with an External Clock (PIC10F200/202)

When an external clock input is used for Timer0, it must meet certain requirements. The external clock requirement is due to internal phase clock (TOSC) synchronization. Also, there is a delay in the actual incrementing of Timer0 after synchronization.

6.1.1 EXTERNAL CLOCK SYNCHRONIZATION

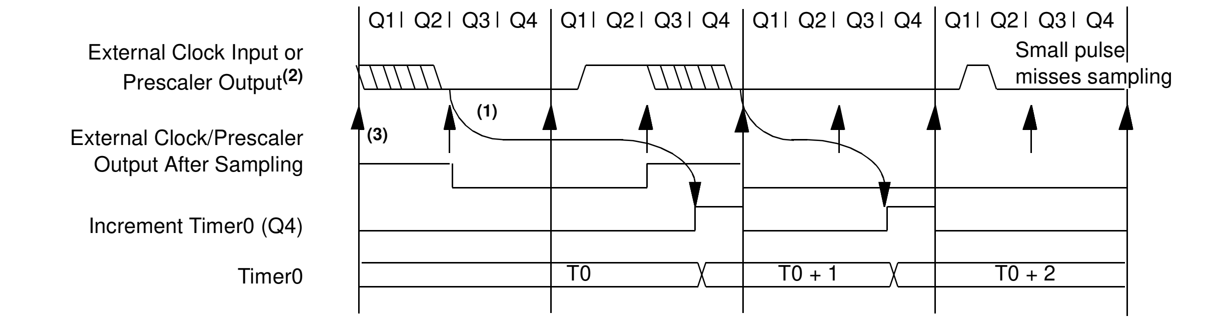
When no prescaler is used, the external clock input is the same as the prescaler output. The synchronization of T0CKI with the internal phase clocks is accomplished by sampling the prescaler output on the Q2 and Q4 cycles of the internal phase clocks (Figure 6-4). Therefore, it is necessary for T0CKI to be high for at least 2 TOSC (and a small RC delay of 2 Tt0H) and low for at least 2 TOSC (and a small RC delay of 2 Tt0H). Refer to the electrical specification of the desired device.

When a prescaler is used, the external clock input is divided by the asynchronous ripple counter-type prescaler, so that the prescaler output is symmetrical. For the external clock to meet the sampling requirement, the ripple counter must be taken into account. Therefore, it is necessary for T0CKI to have a period of at least 4 TOSC (and a small RC delay of 4 Tt0H) divided by the prescaler value. The only requirement on T0CKI high and low time is that they do not violate the minimum pulse width requirement of Tt0H. Refer to parameters 40, 41 and 42 in the electrical specification of the desired device.

6.1.2 TIMER0 INCREMENT DELAY

Since the prescaler output is synchronized with the internal clocks, there is a small delay from the time the external clock edge occurs to the time the Timer0 module is actually incremented. Figure 6-4 shows the delay from the external clock edge to the timer incrementing.

FIGURE 6-4: TIMER0 TIMING WITH EXTERNAL CLOCK


Q1 | Q2 | Q3 | Q4 | Q1 | Q2 | Q3 | Q4 | Q1 | Q2 | Q3 | Q4 | Q1 | Q2 | Q3 | Q4
External Clock Input or Prescaler Output(2)
Small pulse misses sampling
(1)
(3)
External Clock/Prescaler Output After Sampling
Increment Timer0 (Q4)
Timer0
T0
T0 + 1
T0 + 2

Note 1: Delay from clock input change to Timer0 increment is 3 TOSC to 7 TOSC (Duration of Q = TOSC). Therefore, the error in measuring the interval between two edges on Timer0 input = ±4 TOSC max.

2: External clock if no prescaler selected; prescaler output otherwise.

3: The arrows indicate the points in time where sampling occurs.

6.2 Prescaler

An 8-bit counter is available as a prescaler for the Timer0 module or as a postscaler for the Watchdog Timer (WDT), respectively (see **Section 9.6 “Watchdog Timer (WDT)”**). For simplicity, this counter is being referred to as “prescaler” throughout this data sheet.

> **Note:** The prescaler may be used by either the Timer0 module or the WDT, but not both. Thus, a prescaler assignment for the Timer0 module means that there is no prescaler for the WDT and vice versa.

The PSA and PS<2:0> bits (OPTION<3:0>) determine prescaler assignment and prescale ratio.

When assigned to the Timer0 module, all instructions writing to the TMR0 register (e.g., `CLRF 1`, `MOVWF 1`, `BSF 1,x`, etc.) will clear the prescaler. When assigned to WDT, a `CLRWDT` instruction will clear the prescaler along with the WDT. The prescaler is neither readable nor writable. On a Reset, the prescaler contains all ‘`0`’s.

6.2.1 SWITCHING PRESCALER ASSIGNMENT

The prescaler assignment is fully under software control (i.e., it can be changed “on-the-fly” during program execution). To avoid an unintended device Reset, the following instruction sequence (Example 6-1) must be executed when changing the prescaler assignment from Timer0 to the WDT.

EXAMPLE 6-1: CHANGING PRESCALER (TIMER0 → WDT)

```
CLRWDT                ;Clear WDT
CLRF    TMR0          ;Clear TMR0 & Prescaler
MOVLW   ‘00xx1111’b   ;These 3 lines (5, 6, 7)
OPTION                ;are required only if
                      ;desired
CLRWDT                ;PS<2:0> are 000 or 001
MOVLW   ‘00xx1xxx’b   ;Set Postscaler to
OPTION                ;desired WDT rate
```